

Computer algebra for combinatorics

Alin Bostan



MPRI, C-2-22,
February 17, 2022

Give (and prove!) a simple formula for

$$\sum_{k=0}^n (-1)^k \binom{n}{k} \binom{2k}{n}$$

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```
> T:=(-1)^k*binomial(n,k)*binomial(2*k,n):
> first_terms:=seq(add(T, k=0..n), n=0..6):
> guess_rec:=gfun:-listtorec(terms, u(n))[1];
```

$$\{u(n+1) + 2u(n) = 0, \quad u(0) = 1\}$$

```
> rsolve(guess_rec, u(n));
```

$$(-2)^n$$

- ▷ Is this a proof?
- ▷ Can it be turned into a proof?
- ▷ Is this guessing procedure always guaranteed to work?

Give (and prove!) a simple formula for

$$\sum_{k=0}^n (-1)^k \binom{n}{k} \binom{2k}{n}$$



Combinatorial Identities

H. W. Gould

Combinatorial Identities

A STANDARDIZED SET OF TABLES

LISTING 500 BINOMIAL COEFFICIENT SUMMATIONS

"Scientia non habet inimicum nisi ignorantiam"

HENRY W. GOULD
Professor of Mathematics
West Virginia University

Revised Edition
Martinsburg, W. Va. 1972

Give (and prove!) a simple formula for

$$\sum_{k=0}^n (-1)^k \binom{n}{k} \binom{2k}{n}$$

▷ Table look-up:

$$3.64) \quad \sum_{k=0}^n (-1)^k \binom{n}{k} \binom{2k}{j} = (-1)^n \binom{n}{j-n} 2^{2n-j}$$

$$3.150) \quad \sum_{k=0}^n (-1)^k \binom{n}{k} \binom{x+kz}{j} = \begin{cases} 0, & 0 \leq j < n, \\ (-1)^n z^n, & j = n. \end{cases}$$

- ▷ Is this a proof?
- ▷ Can it be turned into a proof?
- ▷ Is this guessing procedure always guaranteed to work?

Give (and prove!) a simple formula for

$$\sum_{k=0}^n (-1)^k \binom{n}{k} \binom{2k}{n}$$

```
> T:=(-1)^k*binomial(n,k)*binomial(2*k,n):  
> sum(T, k=0..n);
```

$$(-2)^n$$

- ▷ Is this a proof?
- ▷ Is it always guaranteed to work?
- ▷ What is behind this proof?

Give (and prove!) a simple formula for $\sum_{k=0}^n (-1)^k \binom{n}{k} \binom{2k}{n}$

```
> T:=(-1)^k*binomial(n,k)*binomial(2*k,n):
> Zpair:=SumTools[Hypergeometric][Zeilberger](T, n, k, Sn):
> tel:=Zpair[1];
```

$$S_n + 2$$

```
> cert:=Zpair[2];
```

$$\frac{(2k - n - 1)(2k - n)(-1)^k \binom{n}{k} \binom{2k}{n}}{(-n + k - 1)(n + 1)}$$

```
> is_zero:=(subs(n=n+1,T) + 2*T) - (subs(k=k+1,cert) - cert):
> simplify(convert(is_zero,GAMMA));
```

$$0$$

Give (and prove!) a simple formula for $\sum_{k=0}^n (-1)^k \binom{n}{k} \binom{2k}{n}$

```
> with(BinomSums):
> T2:=(-1)^k * Binomial2(n,k)*Binomial2(2*k,n):
> S := Sum(t^n*Sum(T2, k=0..infinity), n=0..infinity):
> series(BinomSums[computesum](S, 5), t);
```

$$1 - 2t + 4t^2 - 8t^3 + 16t^4 - 32t^5 + O(t^6)$$

```
> R, ord := BinomSums[sumtores](S, u);
```

$$\frac{1}{2t+1}, [t]$$

Give (and prove!) a simple formula for $\sum_{k=0}^n (-1)^k \binom{n}{k} \binom{2k}{n}$

- ▷ What really happened in the previous slide?
- ▷ The algorithm started from the pre-tabulated formulas

$$\binom{n}{k} = \operatorname{res}_u \frac{(1+u)^n}{u^{k+1}}, \quad \binom{2k}{n} = \operatorname{res}_u \frac{(1+u)^{2k}}{u^{n+1}}$$

- ▷ It then performed the summation $\sum_{n=0}^{\infty} \sum_{k=0}^n (-1)^k \cdot \frac{(1+u)^n}{u^{k+1}} \cdot \frac{(1+u)^{2k}}{u^{n+1}} \cdot t^n$ expressing the GF of the input binomial sum as the residue (w.r.t. u and v) of

$$R := \frac{\frac{1}{v\left(1 - \frac{(1+u)t}{v}\right)} + \frac{(1+v)^2}{uv\left(1 + \frac{(1+v)^2(1+u)t}{uv}\right)}}{v^2 + u + 2v + 1}$$

- ▷ It finally performed a successive pole/residue analysis, proving that

$$\operatorname{res}_{u,v} R = \operatorname{res}_v \frac{v}{t v^3 + 2t v^2 + v^2} = \frac{1}{2t + 1}.$$

Up to $k \leftrightarrow n - k$, the identity is equivalent to
$$\sum_{k=0}^n (-1)^k \binom{n}{k} \binom{2n-2k}{n} = 2^n.$$

▷ Let us re-write binomials as quotients of products of factorials

$$u_n = \sum_{k=0}^n (-1)^k \cdot \frac{n!}{k!(n-k)!} \cdot \frac{(2n-2k)!}{n!(n-2k)!} = \sum_{k=0}^n (-1)^k \cdot \frac{1}{k!} \cdot \frac{1}{(n-k)!} \cdot \frac{(2n-2k)!}{(n-2k)!}$$

and then in terms of “rising factorials” (or, “Pochhammer symbols”)
 $(a)_n = a(a+1) \cdots (a+n-1)$, using the rewriting rules:

$$(n-k)! = \frac{(-1)^k \cdot n!}{(-n)_k} \quad \text{and} \quad (a)_{2k} = 4^k \cdot \left(\frac{a}{2}\right)_k \cdot \left(\frac{a+1}{2}\right)_k$$

▷ We get
$$u_n = \binom{2n}{n} \cdot \sum_{k=0}^n \frac{(-n)_k \cdot (-n)_{2k}}{(1)_k \cdot (-2n)_{2k}} = \binom{2n}{n} \cdot \sum_{k=0}^n \frac{\left(-\frac{n}{2}\right)_k \cdot \left(\frac{-n+1}{2}\right)_k}{(1)_k \cdot \left(\frac{1}{2} - n\right)_k}$$

▷ We conclude using the “Chu-Vandermonde” hypergeometric identity that

$$u_n = \binom{2n}{n} \cdot {}_2F_1\left(-\frac{n}{2}, \frac{1}{2} - \frac{n}{2} \mid 1\right) = 2^n.$$

▷ “We reduced an identity to another identity: what’s the point?”

Up to $k \leftrightarrow n - k$, the identity is equivalent to $\sum_{k=0}^n (-1)^k \binom{n}{k} \binom{2n-2k}{n} = 2^n$.

▷ Consider the polynomial

$$P_n(x) := \frac{1}{n!} \cdot \frac{\partial^n}{\partial x^n} (x^2 - 1)^n$$

▷ By the Leibniz differentiation rule,

$$P_n(x) = \frac{1}{n!} \cdot \sum_{k=0}^n \binom{n}{k} \cdot \frac{\partial^k}{\partial x^k} (x+1)^n \cdot \frac{\partial^{n-k}}{\partial x^{n-k}} (x-1)^n,$$

hence $P_n(x) = \sum_{k=0}^n \binom{n}{k}^2 (x-1)^{n-k} (x+1)^k$, in particular $P_n(1) = 2^n$.

▷ By the binomial theorem, $(x^2 - 1)^n = \sum_{k=0}^n (-1)^k \binom{n}{k} x^{2n-2k}$, hence

$$P_n(x) = \sum_{k=0}^n (-1)^k \binom{n}{k} \binom{2n-2k}{n} x^{n-2k}.$$

▷ In conclusion, $2^n = P_n(1) = \sum_{k=0}^n (-1)^k \binom{n}{k} \binom{2n-2k}{n}$.

We will prove that

$$\sum_{k=0}^n (-1)^k \binom{n}{k} \binom{2n-2k}{n} = 2^n$$

by counting subsets of $\{1, \dots, n\}$ in the following way:

(1) There is an obvious bijection between subsets of $\{1, \dots, n\}$ and subsets with n elements of the set $\{1, \dots, n, \bar{1}, \dots, \bar{n}\}$ which contain either k or \bar{k}

(2) Now to count the latter we can do inclusion/exclusion:

- $\binom{2n}{n}$ counts all n -element sets
- This counts too many, because it counts also subsets which contain both k and \bar{k}
- To delete those, we subtract $\binom{n}{1} \cdot \binom{2n-2}{n-2}$
- But this deletes too many, since it counts those who have k and \bar{k} and ℓ and $\bar{\ell}$ twice
- Hence one adds $\binom{n}{2} \cdot \binom{2n-4}{n-4}$
- And so on.

Summary: ingredients of an experimental mathematics approach

- Trial and error phase
 - Guess the answer Hermite-Padé approximants
 - Look up for possible generalizations bibliographic searches
 - Reasoning and proving phase – understand what's “inside the box”
 - Use built-in routines computer algebra software
 - Use a specific summation approach Zeilberger's algorithm
 - Use an alternative / better one residues
 - Bonus phase – attacking from different angles
 - Hypergeometric approach Chu-Vandermonde identity
 - Direct algebraic approach Legendre polynomials
 - Combinatorial approach Bijection (most human creativity demanding)
- ▷ What is your favorite proof?
- ▷ Why? (Criteria: length/beauty/trickiness/naturalness)

Enumerative Combinatorics: science of counting

Area of mathematics primarily concerned with counting discrete objects.

▷ Main outcome: theorems

Computer Algebra: effective mathematics

Area of computer science primarily concerned with the algorithmic manipulation of algebraic objects.

▷ Main outcome: algorithms

Computer Algebra for **Enumerative Combinatorics**

Today: **Algorithms** for proving **Theorems** on **Lattice Paths Combinatorics**.

An (innocent looking) combinatorial question

Let $\mathcal{S} = \{\uparrow, \leftarrow, \searrow\}$. An \mathcal{S} -walk is a path in \mathbb{Z}^2 using only steps from \mathcal{S} . Show that, for any integer n , the following quantities are equal:

(i) number a_n of n -steps \mathcal{S} -walks confined to the upper half plane $\mathbb{Z} \times \mathbb{N}$ that start and finish at the origin $(0,0)$ (*excursions*);

(ii) number b_n of n -steps \mathcal{S} -walks confined to the quarter plane \mathbb{N}^2 that start at the origin $(0,0)$ and finish on the diagonal of \mathbb{N}^2 (*diagonal walks*).

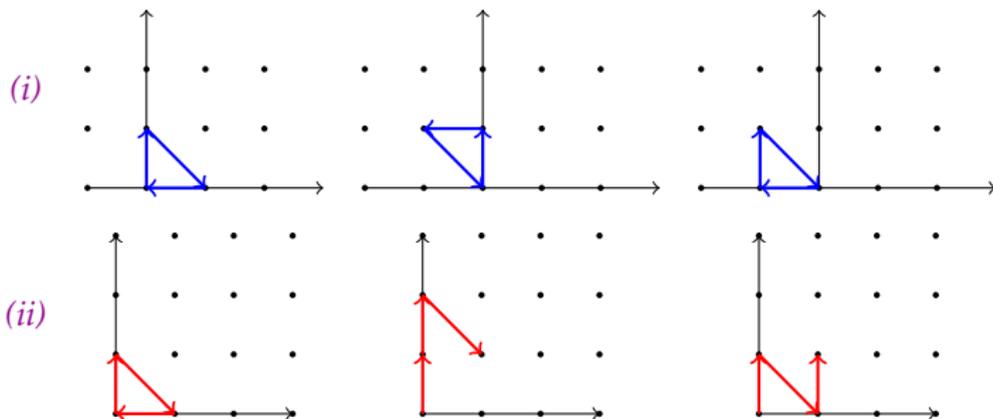
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For instance, for $n = 3$, this common value is $a_3 = b_3 = 3$:



Teaser 1: This “exercise” is non-trivial

Teaser 2: It can be solved using **Experimental Math** and **Computer Algebra**

Teaser 3: ...by two robust and efficient algorithmic techniques,
Guess-and-Prove and **Creative Telescoping**

Why count walks?

Many objects can be encoded by (confined) walks:

- probability theory (voting, games of chance, branching processes, ...)
- discrete mathematics (permutations, trees, words, urns, ...)
- statistical physics (Ising model, ...)
- operations research (queueing theory, ...)

7TH INTERNATIONAL CONFERENCE ON
LATTICE PATH COMBINATORICS AND APPLICATIONS



Siena, Italy July 4-7, 2010

HOME	TOPICS to be covered include (but are not limited to) :	
Photo	Lattice path enumeration	Random walks
Program	Plane Partitions	Non parametric statistical inference
Proceedings	Young tableaux	Discrete distributions and urn models
Submission	q-calculus	Queueing theory
Important dates	Orthogonal polynomials	Analysis of algorithms
Participants		Graph Theory and Applications
General Information		Self-dual codes and unimodular lattices
		Bijections between paths and other combinatoric structures

Counting walks is an old topic: the ballot problem [Bertrand, 1887]

Suppose that candidates A and B are running in an election. If a votes are cast for A and b votes are cast for B , where $a > b$, then the probability that A stays ahead of B throughout the counting of the ballots is $(a - b)/(a + b)$.

Lattice path reformulation: find the number of paths in \mathbb{Z}^2 with a upsteps ↗ and b downsteps ↘ that start at the origin and never touch the x -axis



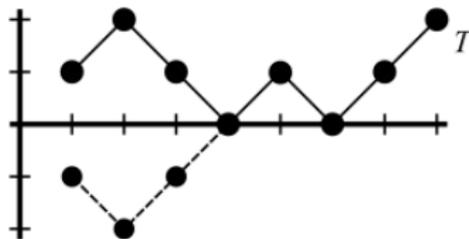
- ▷ Without the constraint, the number of such paths is $\binom{a+b}{a}$
→ a Guess-and-Prove proof in a few slides

Counting walks is an old topic: the ballot problem [Bertrand, 1887]

Suppose that candidates A and B are running in an election. If a votes are cast for A and b votes are cast for B , where $a > b$, then the probability that A stays ahead of B throughout the counting of the ballots is $(a - b)/(a + b)$.

Lattice path reformulation: find the number of paths in \mathbb{Z}^2 with $a - 1$ upsteps \nearrow and b downsteps \searrow that start at $(1, 1)$ and never touch the x -axis

Reflection principle [Aebly, 1923]: paths in \mathbb{Z}^2 from $(1, 1)$ to $T(a + b, a - b)$ that do touch the x -axis are in bijection with paths in \mathbb{Z}^2 from $(1, -1)$ to T



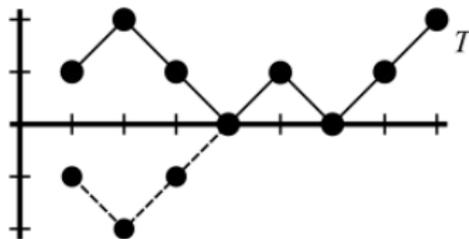
$$\text{Answer: } \underbrace{\binom{a+b-1}{a-1}}_{\text{paths in } \mathbb{Z}^2 \text{ from } (1, 1) \text{ to } T} - \underbrace{\binom{a+b-1}{b-1}}_{\text{paths in } \mathbb{Z}^2 \text{ from } (1, -1) \text{ to } T}$$

Counting walks is an old topic: the ballot problem [Bertrand, 1887]

Suppose that candidates A and B are running in an election. If a votes are cast for A and b votes are cast for B , where $a > b$, then the probability that A stays ahead of B throughout the counting of the ballots is $(a - b)/(a + b)$.

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Reflection principle [Aebly, 1923]: *paths in \mathbb{Z}^2 from $(1, 1)$ to $T(a + b, a - b)$ that do touch the x -axis are in bijection with paths in \mathbb{Z}^2 from $(1, -1)$ to T*



Answer: $\underbrace{(\text{paths in } \mathbb{Z}^2 \text{ from } (1, 1) \text{ to } T) - (\text{paths in } \mathbb{Z}^2 \text{ from } (1, -1) \text{ to } T)}$

$$\binom{a+b-1}{a-1} - \binom{a+b-1}{b-1} = \frac{a-b}{a+b} \binom{a+b}{a}$$

Lot of recent activity; many recent contributors:

Arquès, Bacher, Banderier, Beaton, Bernardi, Biane, Bostan, Bousquet-Mélou, Buchacher, Budd, Chyzak, Cori, Courtiel, Denisov, Dreyfus, Du, Duchon, Dulucq, Duraj, Fayolle, Fisher, Flajolet, Fusy, Garbit, Gessel, Gouyou-Beauchamps, Guttmann, Guy, Hardouin, van Hoeij, Hou, Iasnogorodski, Johnson, Kauers, Kenyon, Koutschan, Krattenthaler, Kreweras, Kurkova, Lecouvey, Malyshev, Melczer, Miller, Mishna, Niederhausen, Owczarek, Pech, Petkovšek, Prellberg, Raschel, Rechnitzer, Roques, Sagan, Salvy, Sheffield, Singer, Tarrago, Trotignon, Verron, Viennot, Wachtel, Wallner, Wang, Wilf, D. Wilson, M. Wilson, Xu, Yatchak, Yeats, Zeilberger, ...

etc.

...but it is still a very hot topic

Lot of recent activity; many recent contributors:

Arquès, Bacher, Banderier, Beaton, Bernardi, Biane, Bostan, Bousquet-Mélou, Buchacher, Budd, Chyzak, Cori, Courtiel, Denisov, Dreyfus, Du, Duchon, Dulucq, Duraj, Fayolle, Fisher, Flajolet, Fusy, Garbit, Gessel, Gouyou-Beauchamps, Guttmann, Guy, Hardouin, van Hoeij, Hou, Iasnogorodski, Johnson, Kauers, Kenyon, Koutschan, Krattenthaler, Kreweras, Kurkova, Lecouvey, Malyshev, Melczer, Miller, Mishna, Niederhausen, Owczarek, Pech, Petkovšek, Prellberg, Raschel, Rechnitzer, Roques, Sagan, Salvy, Sheffield, Singer, Tarrago, Trotignon, Verron, Viennot, Wachtel, Wallner, Wang, Wilf, D. Wilson, M. Wilson, Xu, Yatchak, Yeats, Zeilberger, ...

etc.

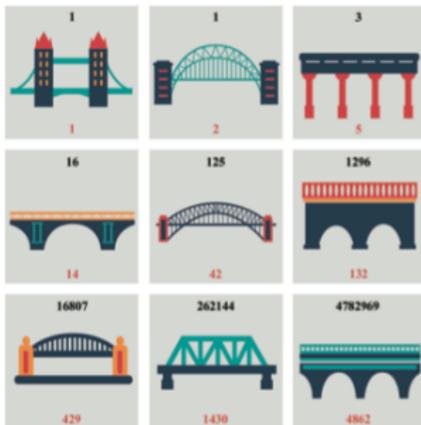
~~Specific question
Ad hoc solution~~



Systematic approach

DISCRETE MATHEMATICS AND ITS APPLICATIONS

HANDBOOK OF ENUMERATIVE COMBINATORICS



Edited by
Miklós Bóna

 **CRC Press**
Taylor & Francis Group
A CHAPMAN & HALL BOOK

Chapter 10

Lattice Path Enumeration

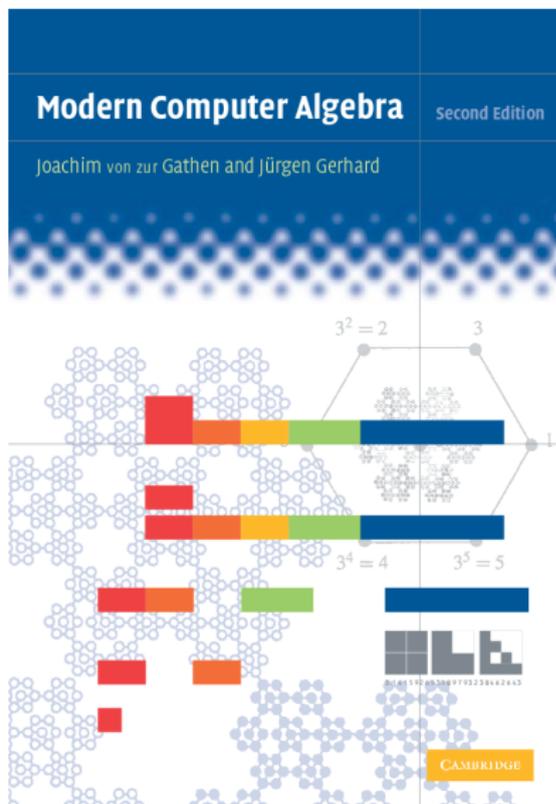
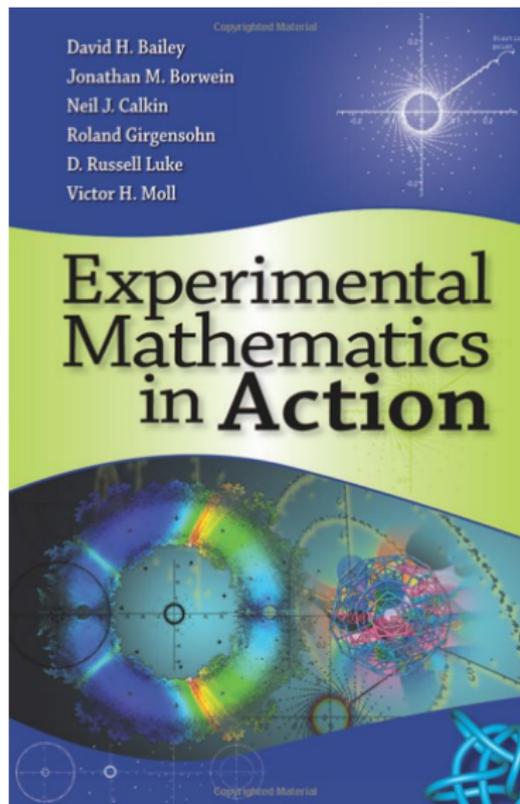
Christian Krattenthaler

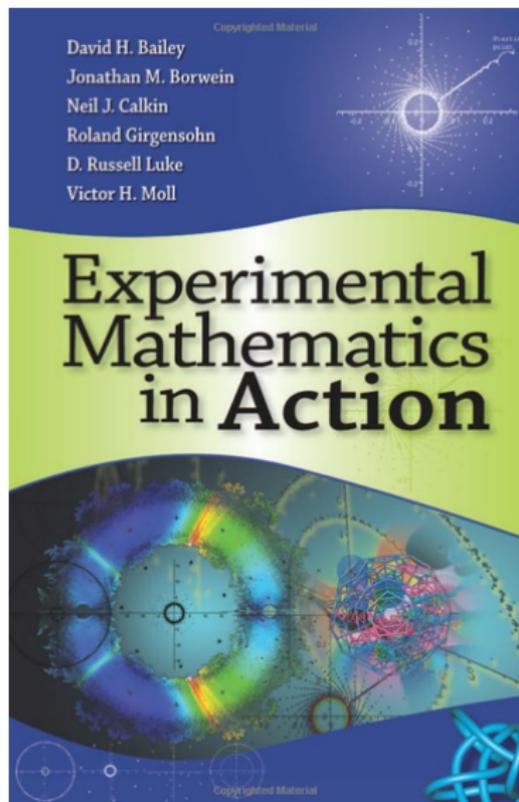
Universität Wien

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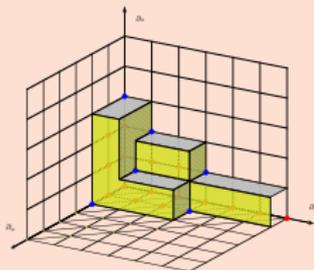
Our approach: Experimental Mathematics using Computer Algebra





Algorithmes Efficaces en Calcul Formel

Alin BOSTAN
Frédéric CHYZAK
Marc GIUSTI
Romain LEBRETON
Grégoire LECERF
Bruno SALVY
Éric SCHOST



- ▷ Walks in \mathbb{N}^2 starting at $(0,0)$ and using steps in a fixed subset \mathcal{S} of

$$\{\swarrow, \leftarrow, \nearrow, \uparrow, \searrow, \rightarrow, \swarrow, \downarrow\}.$$

- ▷ Counting sequence: $q_{\mathcal{S}}(n)$ = number of \mathcal{S} -walks of length n

- ▷ Length generating function:

$$Q_{\mathcal{S}}(t) = \sum_{n=0}^{\infty} q_{\mathcal{S}}(n)t^n \in \mathbb{Z}[[t]]$$

Lattice walks with small steps in the quarter plane

- ▷ Walks in \mathbb{N}^2 starting at $(0,0)$ and using steps in a fixed subset \mathcal{S} of

$$\{\swarrow, \leftarrow, \nearrow, \uparrow, \searrow, \rightarrow, \downarrow\}.$$

- ▷ Refinement: $q_{\mathcal{S}}(i, j; n)$ = number of \mathcal{S} -walks of length n ending at (i, j)

- ▷ Full generating function (with “catalytic ” variables x, y):

$$Q_{\mathcal{S}}(x, y; t) = \sum_{i, j, n=0}^{\infty} q_{\mathcal{S}}(i, j; n) x^i y^j t^n \in \mathbb{Z}[[x, y, t]]$$

- ▷ Actually: $Q_{\mathcal{S}}(x, y; t) \in \mathbb{Z}[x, y][[t]]$ and $Q_{\mathcal{S}}(1, 1; t) = Q_{\mathcal{S}}(t)$

Entire books dedicated to small-steps walks in the quarter plane!

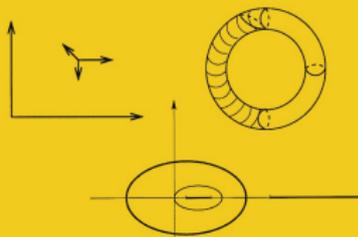
Applications of Mathematics
Stochastic Modelling and Applied Probability

40

Guy Fayolle
Roudolf Iasnogorodski
Vadim Malyshev

Random Walks in the Quarter-Plane

Algebraic Methods,
Boundary Value Problems
and Applications



Probability Theory and Stochastic Modelling 40

Guy Fayolle
Roudolf Iasnogorodski
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Random Walks in the Quarter Plane

Algebraic Methods, Boundary Value
Problems, Applications to Queueing
Systems and Analytic Combinatorics

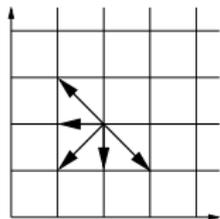
Second Edition



Among the 2^8 step sets $\mathcal{S} \subseteq \{-1, 0, 1\}^2 \setminus \{(0, 0)\}$, some are:

Small-steps models of interest

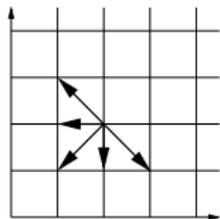
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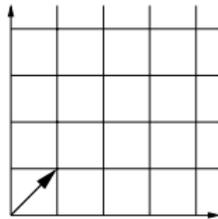
trivial,

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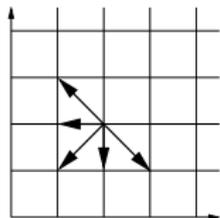
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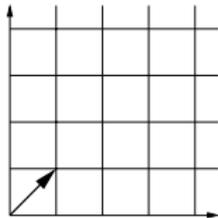
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Small-steps models of interest

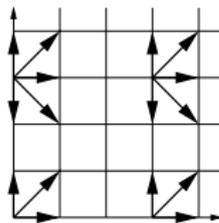
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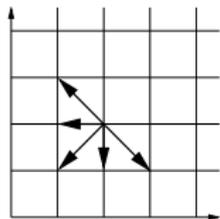
simple,



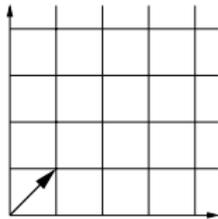
intrinsic to the
half plane,

Small-steps models of interest

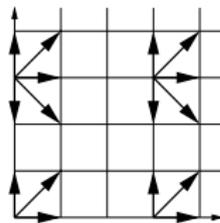
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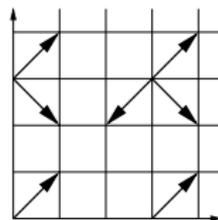
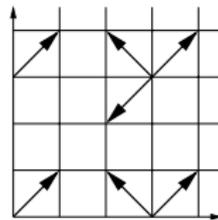
trivial,



simple,



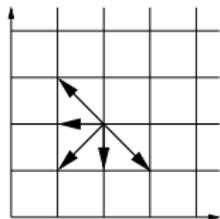
intrinsic to the
half plane,



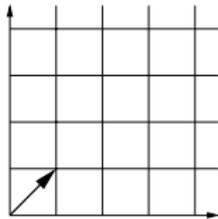
symmetrical.

Small-steps models of interest

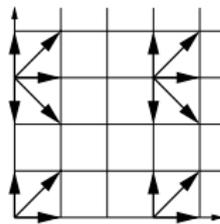
Among the 2^8 step sets $\mathcal{S} \subseteq \{-1, 0, 1\}^2 \setminus \{(0, 0)\}$, some are:



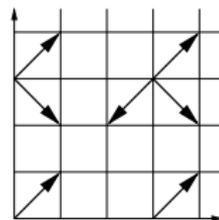
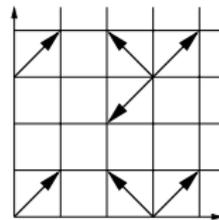
trivial,



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One is left with [79 interesting distinct models](#).

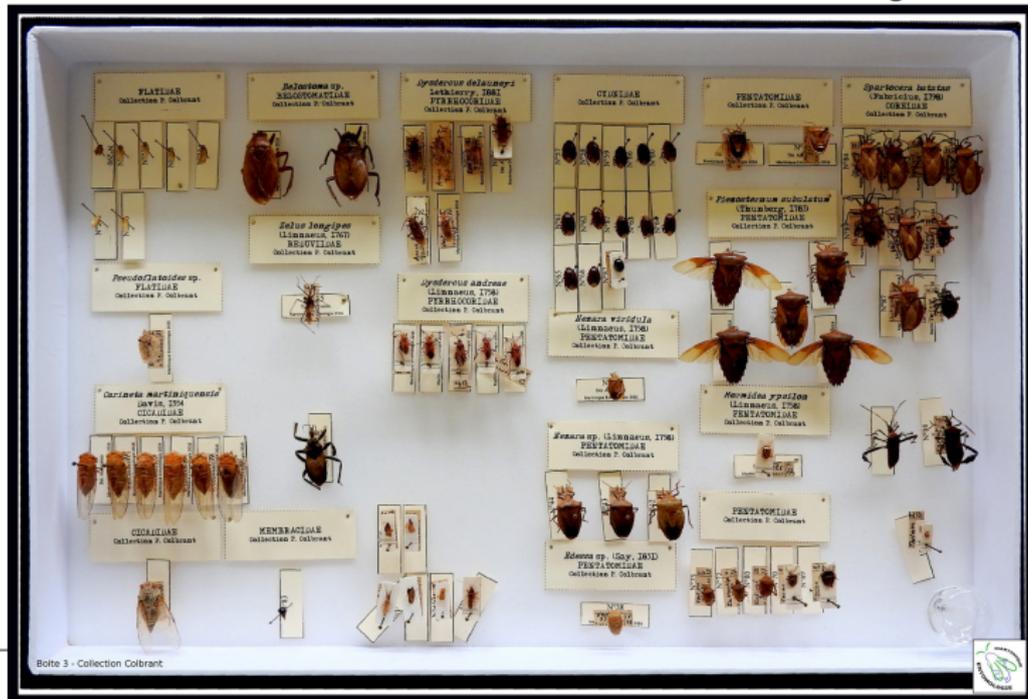
The 79 small-steps quadrant models



Task: classify their generating functions!

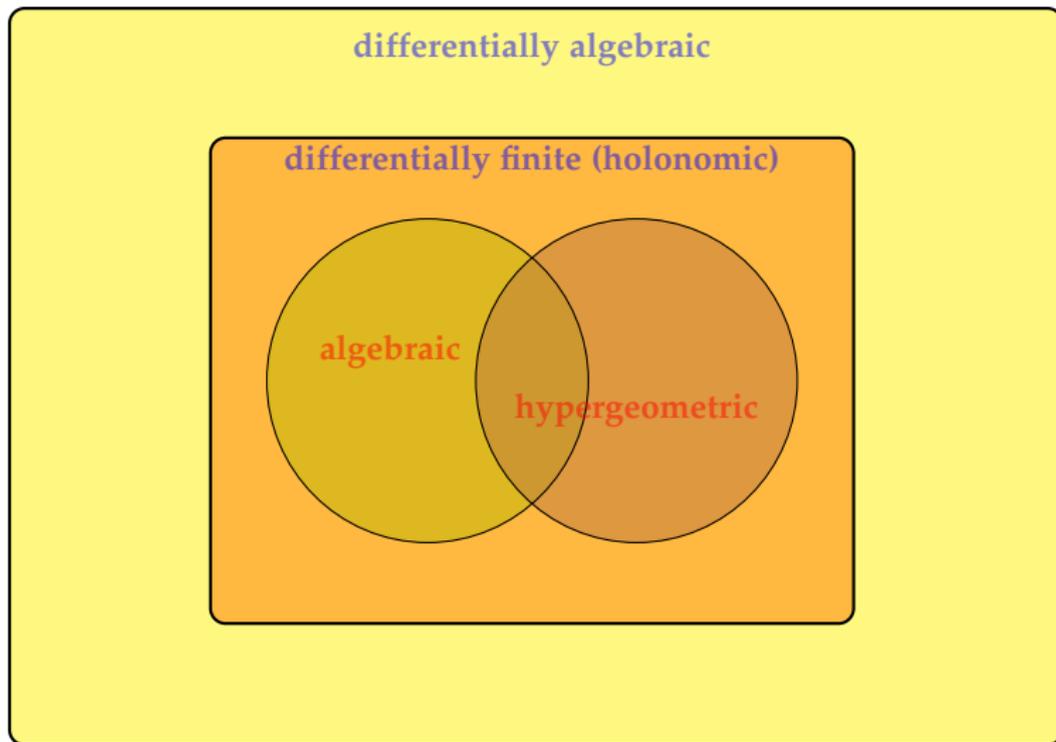


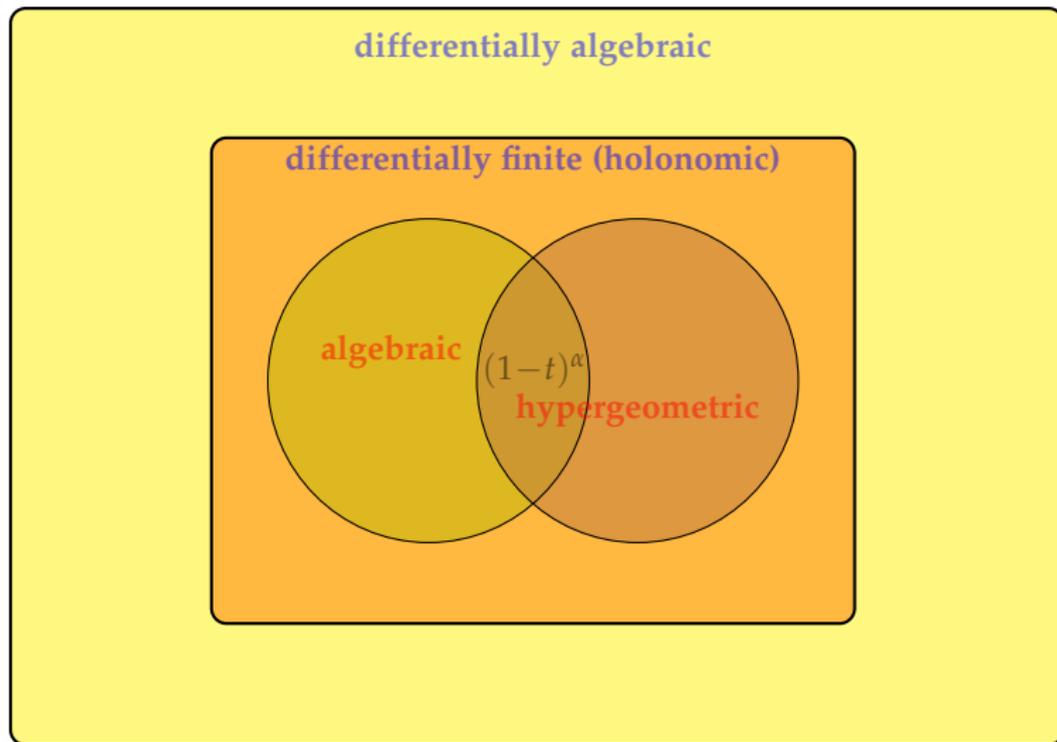
Non-singular

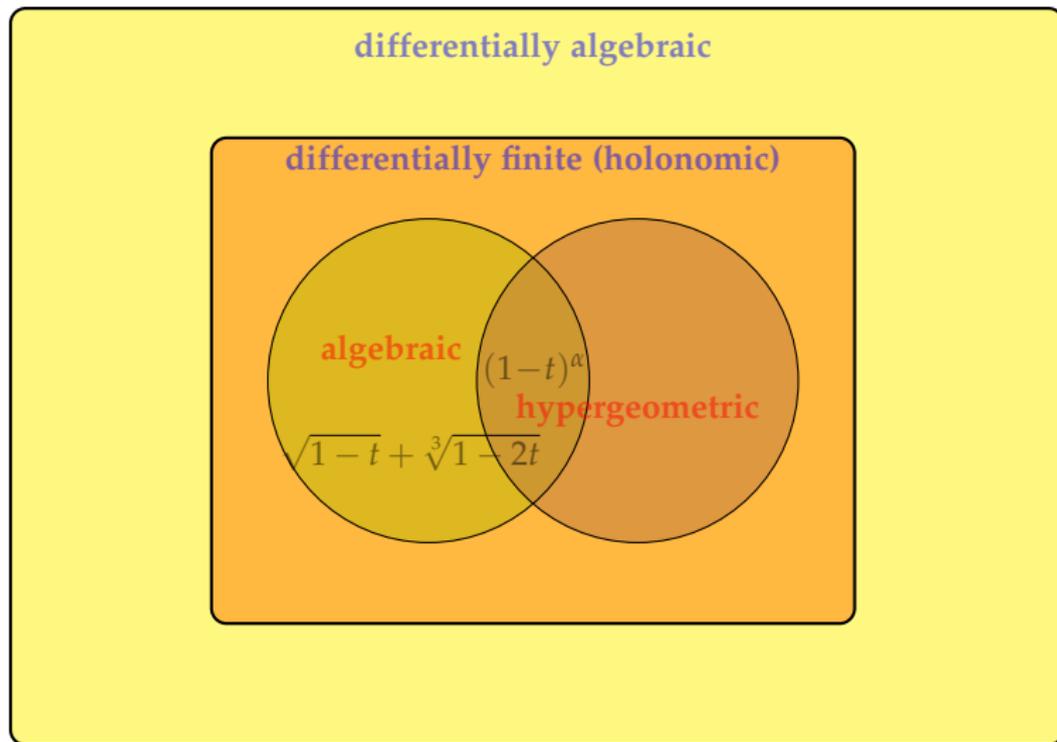


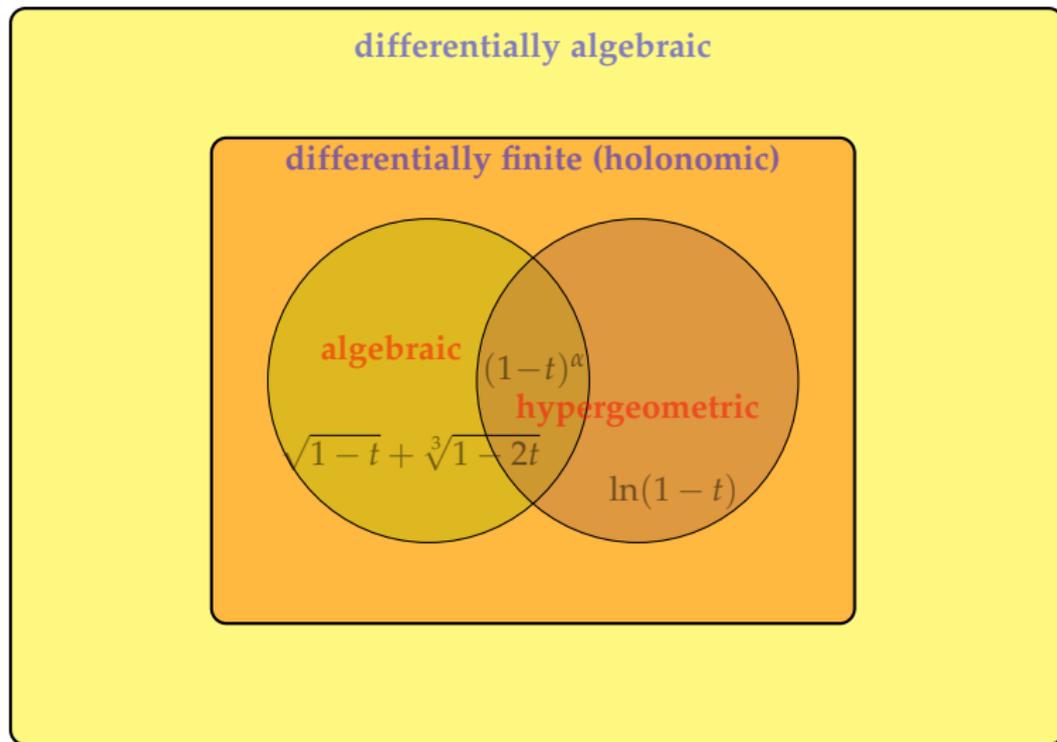
Singular

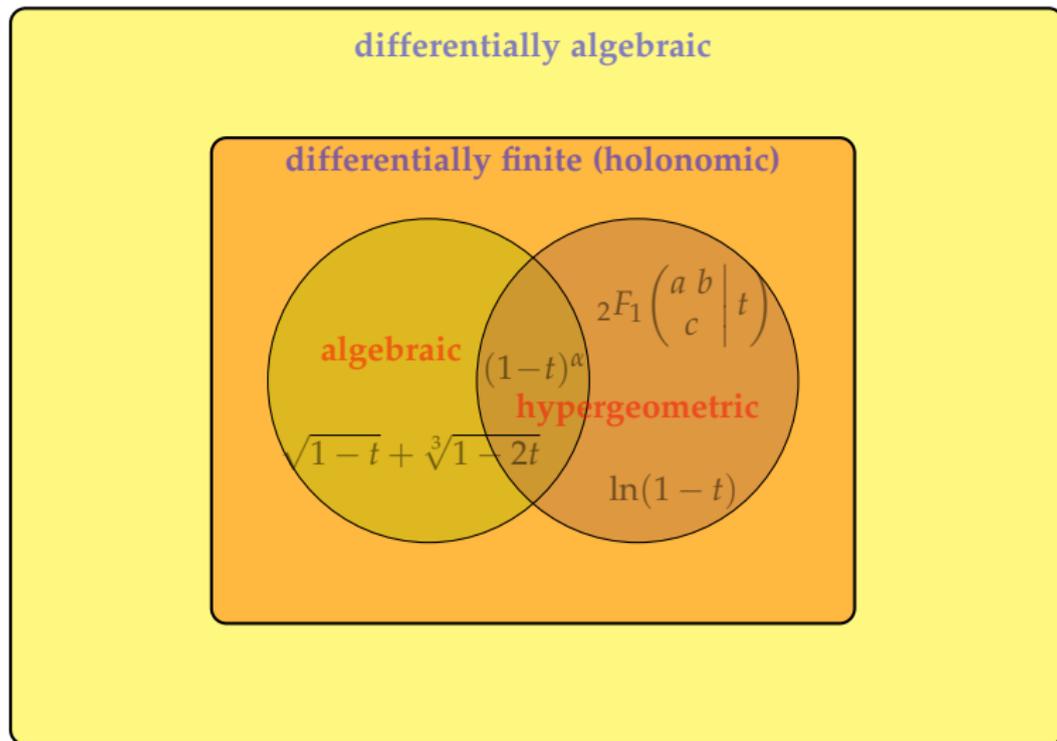




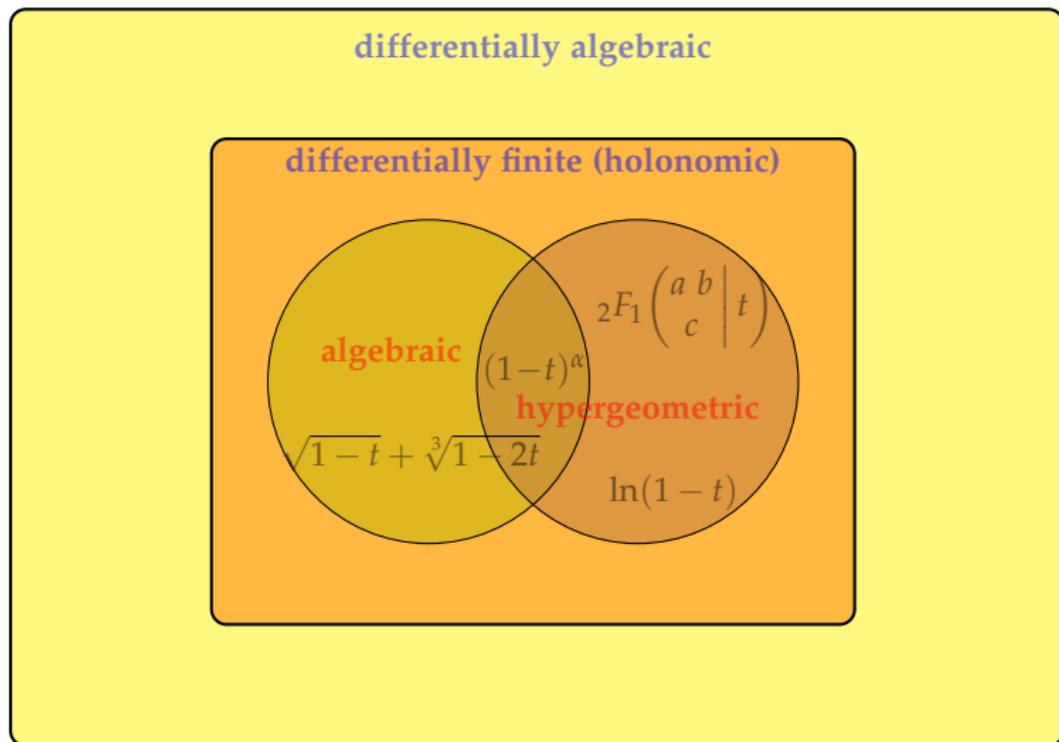




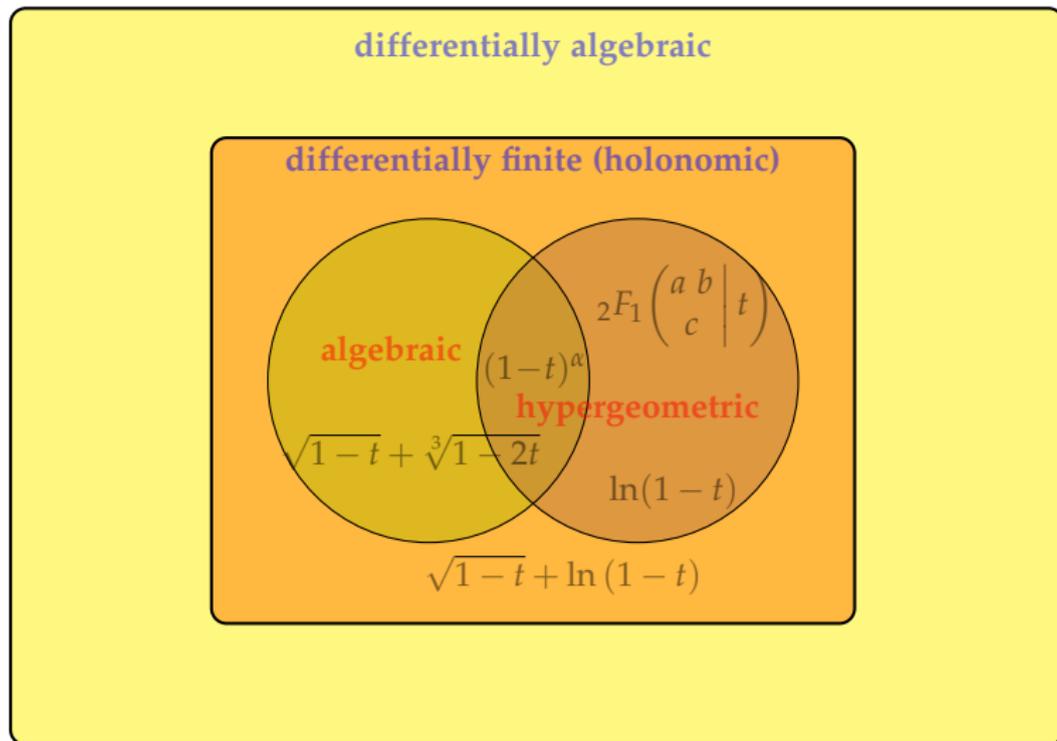




$${}_2F_1\left(\begin{matrix} a & b \\ c \end{matrix} \middle| t\right) := \sum_{n=0}^{\infty} \frac{(a)_n (b)_n}{(c)_n} \frac{t^n}{n!}, \quad \text{where } (a)_n = a(a+1) \cdots (a+n-1).$$

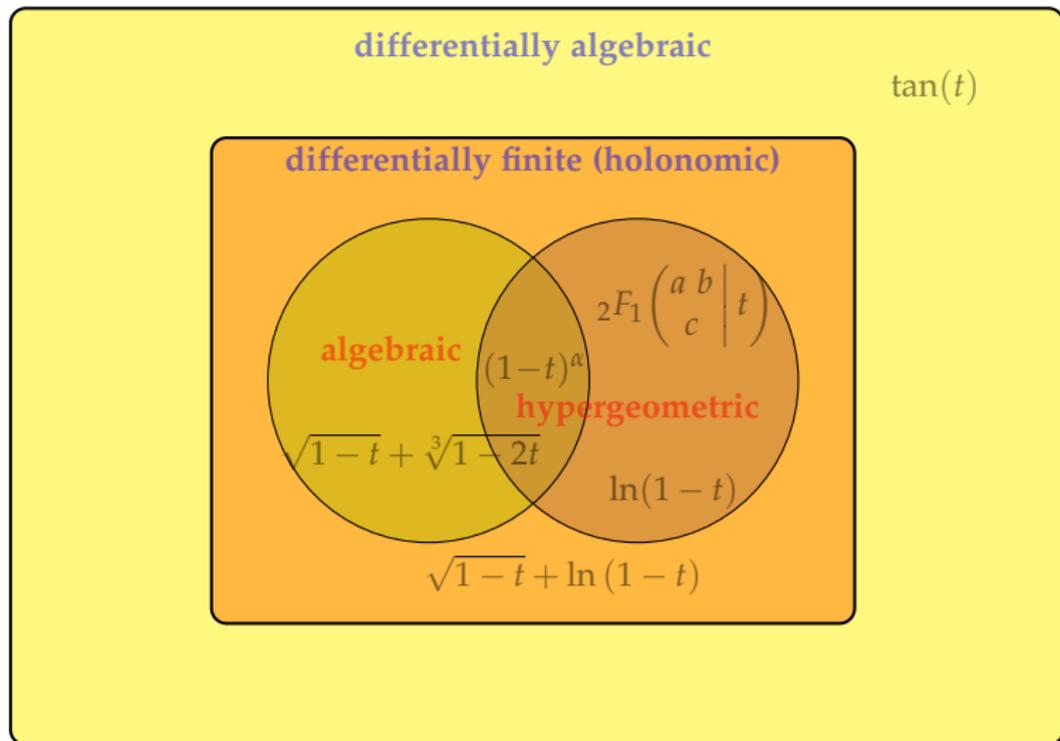


E.g., $(1-t)^\alpha = {}_2F_1\left(\begin{matrix} -\alpha & 1 \\ 1 \end{matrix} \middle| t\right)$, $\ln(1-t) = -t \cdot {}_2F_1\left(\begin{matrix} 1 & 1 \\ 2 \end{matrix} \middle| t\right) = -\sum_{n=1}^{\infty} \frac{t^n}{n}$



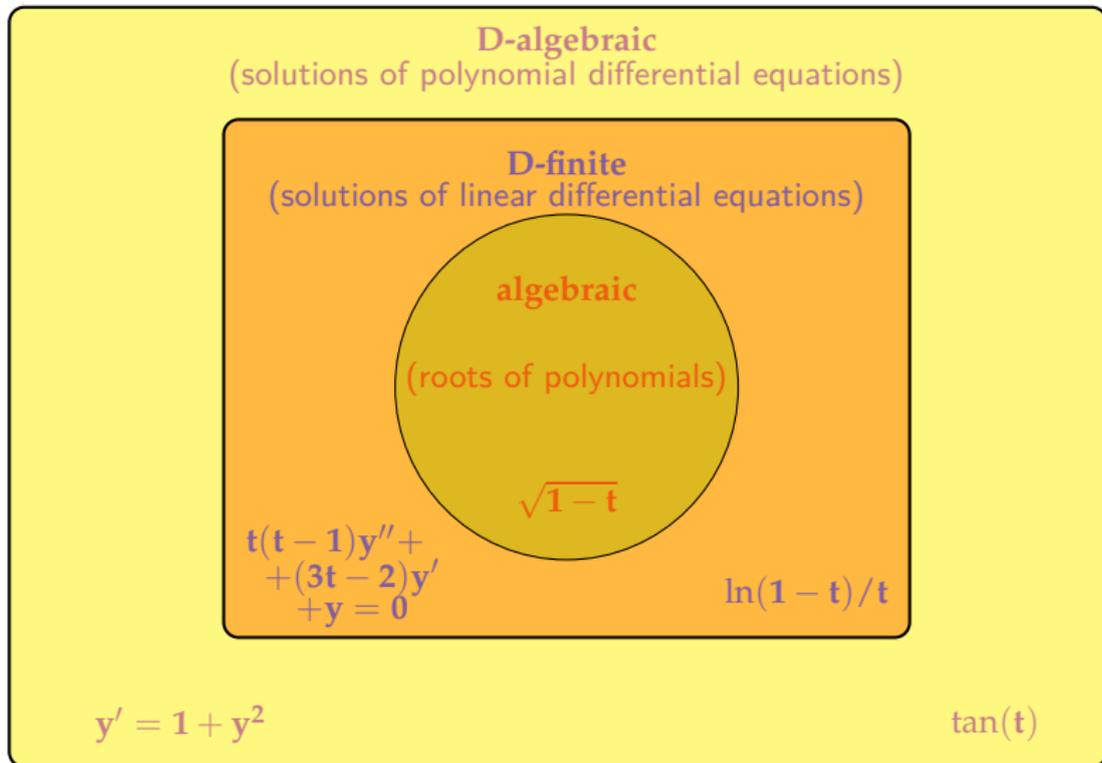
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Classification criterion: properties of generating functions



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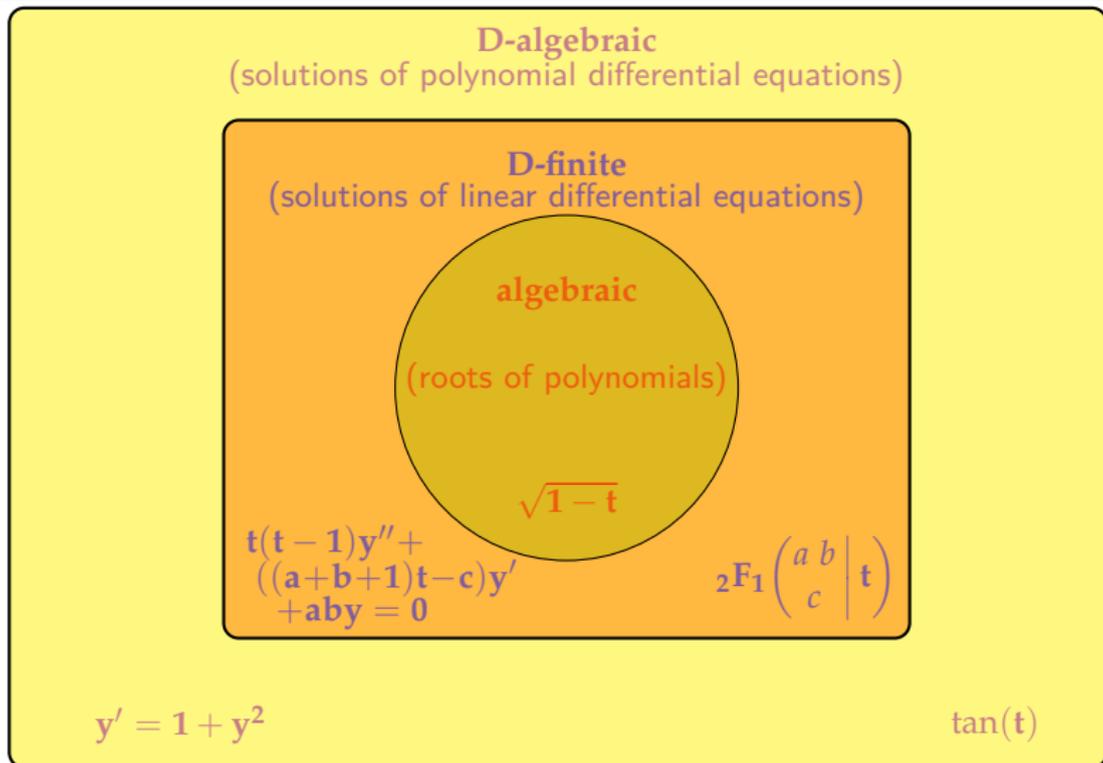
Classification criterion: properties of generating functions



D-transcendental

$$\Gamma(t) = \int_0^{\infty} x^{t-1} e^{-x} dx$$

Classification criterion: properties of generating functions



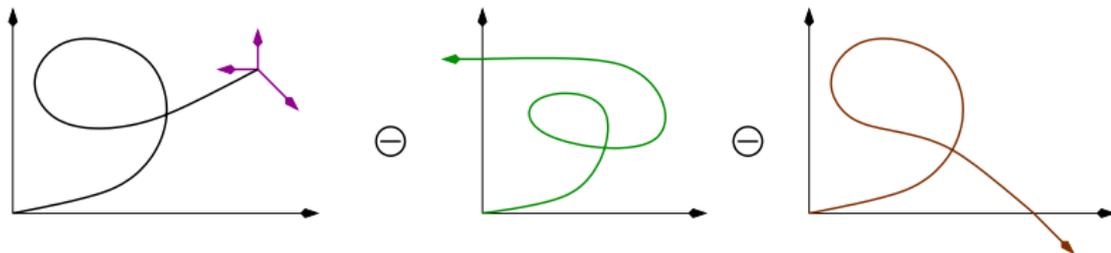
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Algebraic reformulation of main task: solving a functional equation

Generating function: $Q(x, y) \equiv Q(x, y; t) = \sum_{i, j, n=0}^{\infty} q(i, j; n) x^i y^j t^n \in \mathbb{Z}[[x, y, t]]$

Recursive construction yields the *kernel equation*

$$Q(x, y) = 1 + t \left(y + \frac{1}{x} + x \frac{1}{y} \right) Q(x, y) - t \frac{1}{x} Q(0, y) - tx \frac{1}{y} Q(x, 0)$$

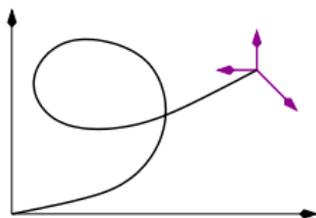


Algebraic reformulation of main task: solving a functional equation

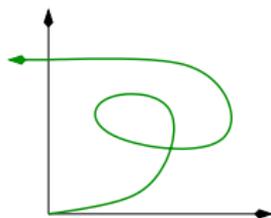
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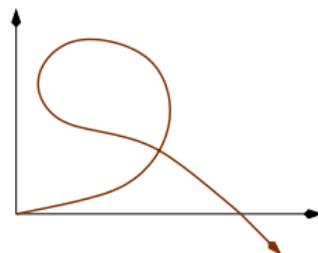
$$\left(1 - t \left(y + \frac{1}{x} + x \frac{1}{y}\right)\right) xyQ(x, y) = xy - tyQ(0, y) - tx^2Q(x, 0)$$



⊖



⊖

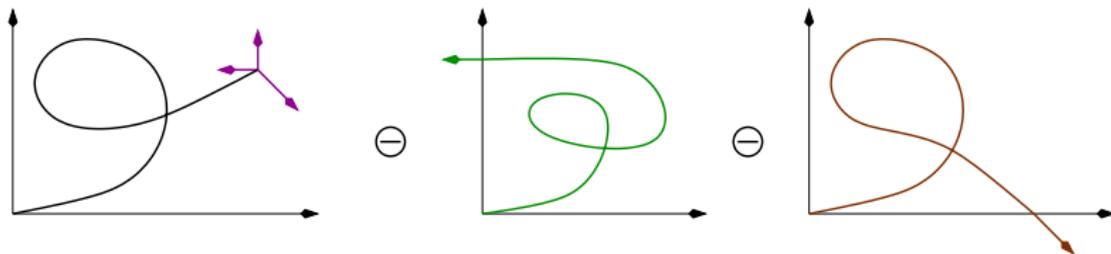


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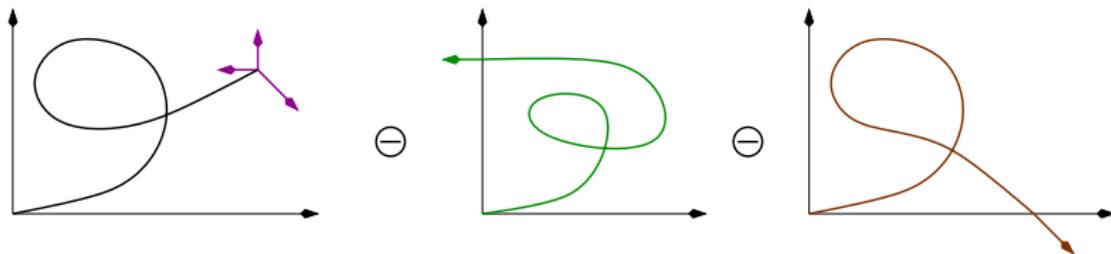
New task: Solve this functional equation!

Algebraic reformulation of main task: solving a functional equation

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$$\left(1 - t \left(y + \frac{1}{x} + x \frac{1}{y}\right)\right) xyQ(x, y) = xy - tyQ(0, y) - tx^2Q(x, 0)$$



New task: For the other models – solve 78 similar equations!

“Special” models of walks in the quarter plane

Dyck: 

Motzkin: 

Pólya: 

Kreweras: 

Gessel: 

Gouyou-Beauchamps: 

King walks: 

Tandem walks: 

▷ Kernel equation:

$$(y - tx(1 + y^2)) \cdot Q(x, y) = y - tx \cdot Q(x, 0)$$

▷ Kernel method [Knuth, 1968]:

- let $y_0 \in \mathbb{Q}[t][[x]]$ be the power series root of $K = y - tx(1 + y^2)$

$$y_0 = \frac{1 - \sqrt{1 - 4t^2x^2}}{2tx} = tx + t^3x^3 + 2t^5x^5 + \dots \in \mathbb{Q}[t][[x]]$$

- plug $y = y_0$ in the kernel equation $\implies Q(x, 0) = \frac{y_0}{tx}$
- conclude algebraicity:

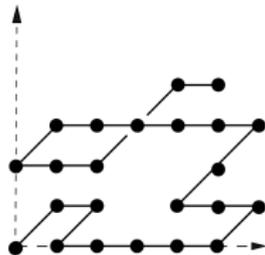
$$Q(x, y) = \frac{y - y_0}{K} = \frac{\sqrt{1 - 4t^2x^2} + 2txy - 1}{2tx(y - tx(1 + y^2))}$$

▷ Same method proves *algebraicity* for all models intrinsic to the half plane

- $g(i, j; n)$ = number of n -steps $\{\nearrow, \swarrow, \leftarrow, \rightarrow\}$ -walks in \mathbb{N}^2 from $(0, 0)$ to (i, j)

Question: What is the nature of the generating function

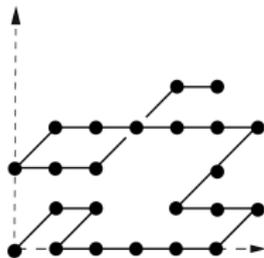
$$G(x, y; t) = \sum_{i, j, n=0}^{\infty} g(i, j; n) x^i y^j t^n ?$$



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Theorem [B., Kauers, 2010]

$G(x, y; t)$ is an algebraic function[†].

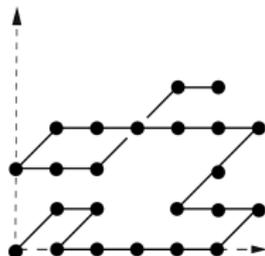
▷ computer-driven discovery/proof via *algorithmic Guess-and-Prove*

[†] Minimal polynomial $P(G(x, y; t); x, y, t) = 0$ has $> 10^{11}$ terms; ≈ 30 Gb (6 DVDs!)

- $g(n)$ = number of n -steps $\{\nearrow, \nwarrow, \leftarrow, \rightarrow\}$ -walks in \mathbb{N}^2

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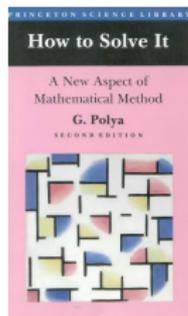
$$G(t) = \sum_{n=0}^{\infty} g(n) t^n ?$$



Corollary [B., Kauers, 2010] (former conjecture of Gessel's)

$$(3n + 1) g(2n) = (12n + 2) g(2n - 1) \text{ and } (n + 1) g(2n + 1) = (4n + 2) g(2n)$$

▷ computer-driven discovery/proof via *algorithmic Guess-and-Prove*



Guessing and Proving

George Pólya

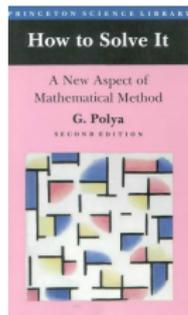


What is “scientific method”? Philosophers and non-philosophers have discussed this question and have not yet finished discussing it. Yet as a first introduction it can be described in three syllables:

Guess and test.

Mathematicians too follow this advice in their research although they sometimes refuse to confess it. They have, however, something which the other scientists cannot really have. For mathematicians the advice is

First guess, then prove.



Guessing and Proving

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Guess-and-Prove: a toy example

Question: Find $B_{i,j} :=$ the number of $\{\rightarrow, \uparrow\}$ -walks in \mathbb{N}^2 from $(0,0)$ to (i,j)

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- ① There are 2 ways to get to (i,j) , either from $(i-1,j)$, or from $(i,j-1)$:

$$B_{i,j} = B_{i-1,j} + B_{i,j-1}$$

- ② There is only one way to get to a point on an axis: $B_{i,0} = B_{0,j} = 1$

▷ These two rules completely determine all the numbers $B_{i,j}$

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(I) Generate data:

⋮							
1	7	28	84	210	462	924	
1	6	21	56	126	252	462	
1	5	15	35	70	126	210	
1	4	10	20	35	56	84	
1	3	6	10	15	21	28	
1	2	3	4	5	6	7	
1	1	1	1	1	1	1	...

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(II) Guess:

→ ...

→ $\frac{(i+1)(i+2)}{2}$

→ $i+1$

→ 1

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(II) Guess:

$$B_{i,j} \stackrel{?}{=} \frac{(i+j)!}{i!j!}$$

Guess-and-Prove: a toy example

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(III) Prove: If

$C_{i,j} \stackrel{\text{def}}{=} \frac{(i+j)!}{i!j!}$, then

$$\frac{C_{i-1,j}}{C_{i,j}} + \frac{C_{i,j-1}}{C_{i,j}} = \frac{i}{i+j} + \frac{j}{i+j} = 1$$

and $C_{i,0} = C_{0,j} = 1$.

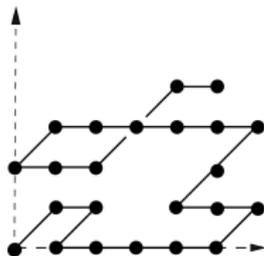
Thus $B_{i,j} = C_{i,j}$

Guess-and-Prove for Gessel walks

- $g(i, j; n)$ = number of n -steps $\{\nearrow, \swarrow, \leftarrow, \rightarrow\}$ -walks in \mathbb{N}^2 from $(0, 0)$ to (i, j)

Question: What is the nature of the generating function

$$G(x, y; t) = \sum_{i, j, n=0}^{\infty} g(i, j; n) x^i y^j t^n ?$$



Answer: [B., Kauers, 2010] $G(x, y; t)$ is an algebraic function[†].

Approach: → very general and robust!

- ① **Generate data:** compute G to precision t^{1200} (≈ 1.5 billion coeffs!)
- ② **Guess:** conjecture polynomial equations for $G(x, 0; t)$ and $G(0, y; t)$ (degree 24 each, coeffs. of degree (46, 56), with 80-bits digits coeffs.)
- ③ **Prove:** multivariate resultants of (very big) polynomials (30 pages each)

[†] Minimal polynomial $P(G(x, y; t); x, y, t) = 0$ has $> 10^{11}$ terms; ≈ 30 Gb (6 DVDs!)

A typical Guess-and-Prove algorithmic proof

Theorem ["Gessel excursions are algebraic"]

$$g(t) := G(0,0; \sqrt{t}) = \sum_{n=0}^{\infty} \frac{(5/6)_n (1/2)_n}{(5/3)_n (2)_n} (16t)^n \text{ is algebraic.}$$

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Proof: First **guess** a polynomial $P(t, T)$ in $\mathbb{Q}[t, T]$, then **prove** that P admits the power series $g(t) = \sum_{n=0}^{\infty} g_n t^n$ as a root.

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- 2 **Implicit function theorem:** $\exists!$ root $r(t) \in \mathbb{Q}[[t]]$ of P .

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- ② **Implicit function theorem:** $\exists!$ root $r(t) \in \mathbb{Q}[[t]]$ of P .
- ③ $r(t) = \sum_{n=0}^{\infty} r_n t^n$ **being algebraic, it is D-finite**, and so (r_n) is **P-recursive**:

$$(n+2)(3n+5)r_{n+1} - 4(6n+5)(2n+1)r_n = 0, \quad r_0 = 1$$

$$\implies \text{solution } r_n = \frac{(5/6)_n (1/2)_n}{(5/3)_n (2)_n} 16^n = g_n, \text{ thus } g(t) = r(t) \text{ is algebraic.}$$

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```
> P:=gfun:-listtoalgeq([seq(pochhammer(5/6,n)*pochhammer(1/2,n)/
  pochhammer(5/3,n)/pochhammer(2,n)*16^n, n=0..100)], g(t)):
> gfun:-diffeqtoec(gfun:-algeqtodiffeq(P[1], g(t)), g(t), r(n));
```

A typical Guess-and-Prove algorithmic proof

Theorem ["Gessel excursions are algebraic"]

$$g(t) := G(0,0;\sqrt{t}) = \sum_{n=0}^{\infty} \frac{(5/6)_n(1/2)_n}{(5/3)_n(2)_n} (16t)^n \text{ is algebraic.}$$

Proof: First **guess** a polynomial $P(t, T)$ in $\mathbb{Q}[t, T]$, then **prove** that P admits the power series $g(t) = \sum_{n=0}^{\infty} g_n t^n$ as a root.

- 1 Find P such that $P(t, g(t)) = 0 \pmod{t^{100}}$ by **(structured) linear algebra**.
- 2 **Implicit function theorem:** $\exists!$ root $r(t) \in \mathbb{Q}[[t]]$ of P .
- 3 $r(t) = \sum_{n=0}^{\infty} r_n t^n$ **being algebraic, it is D-finite**, and so (r_n) is **P-recursive**:

$$(n+2)(3n+5)r_{n+1} - 4(6n+5)(2n+1)r_n = 0, \quad r_0 = 1$$

$$\implies \text{solution } r_n = \frac{(5/6)_n(1/2)_n}{(5/3)_n(2)_n} 16^n = g_n, \text{ thus } g(t) = r(t) \text{ is algebraic.}$$

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> P:=gfun:-listtoalgeq([seq(pochhammer(5/6,n)*pochhammer(1/2,n)/
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```

▷ Steps 1 & 3 rely on **polynomial linear algebra** (*Hermite-Padé approximants*).

Sur la généralisation des fractions continues algébriques.

(Par M. CH. HERMITE, membre de l'Institut, à Paris.)

[Extrait d'une lettre à M. Pincherle (*).]

. Le problème que j'ai en vue est le suivant: Etant donné n séries S_1, S_2, \dots, S_n procédant suivant les puissances d'une variable x , déterminer les polynômes X_1, X_2, \dots, X_n des degrés $\mu_1, \mu_2, \dots, \mu_n$ de manière à avoir

$$S_1 X_1 + S_2 X_2 + \dots + S_n X_n = S x^{\mu_1 + \mu_2 + \dots + \mu_n + n - 1},$$

où S est une série de même nature que S_1, S_2 , etc. La question ainsi posée est entièrement déterminée, et une remarque de calcul intégral en donne la complète solution dans le cas particulier où les séries sont de simples exponentielles. C'est ce que je vais montrer, je me proposerai ensuite de faire sortir, en vue du cas général, les enseignements que contient cette solution.

*Sur la généralisation des fractions continues algébriques;***PAR M. H. PADÉ,**Docteur ès Sciences mathématiques,
Professeur au lycée de Lille.

INTRODUCTION.

M. Hermite s'est, dans un travail récemment paru (1), occupé de la généralisation des fractions continues algébriques. La question est de déterminer les polynômes X_1, X_2, \dots, X_n , de degrés $\mu_1, \mu_2, \dots, \mu_n$, qui satisfont à l'équation

$$S_1 X_1 + S_2 X_2 + \dots + S_n X_n = S x^{\mu_1 + \mu_2 + \dots + \mu_n + n - 1},$$

S_1, S_2, \dots, S_n étant des séries entières données, et S une série également entière. Ou plutôt, il s'agit d'obtenir un algorithme qui permette le calcul de proche en proche de ces systèmes de n polynômes, et qui

Definition: A **Hermite-Padé approximant** of type $\mathbf{d} = (d_1, \dots, d_n) \in \mathbb{N}^n$ for $\mathbf{F} = (f_1, \dots, f_n) \in \mathbb{K}[[x]]^n$ is a $\mathbf{P} = (P_1, \dots, P_n) \in \mathbb{K}[x]^n \setminus \{0\}$ such that:

- (1) $P_1 f_1 + \dots + P_n f_n = O(x^\sigma)$ with $\sigma = \sum_i (d_i + 1) - 1$,
- (2) $\deg(P_i) \leq d_i$ for all i .

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- ▷ Very useful concept in number theory (irrationality/transcendence):
 - [Hermite, 1873]: e is transcendent; [Lindemann, 1882]: π is transc.
 - [Apéry, 1978; Beukers, 1981]: $\zeta(3) = \sum_{n \geq 1} \frac{1}{n^3}$ is irrational;
[Rivoal, 2000]: there are infinitely many k such that $\zeta(2k+1) \notin \mathbb{Q}$.
- ▷ Very useful tool in computer algebra
 - **algebraic approximants** when $f_\ell = A^{\ell-1}$ for a given $A \in \mathbb{K}[[x]]$
 - **differential approximants** when $f_\ell = A^{(\ell-1)}$ for a given $A \in \mathbb{K}[[x]]$

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▷ Many fast algorithms: [Beckermann, Labahn, 1994], [Giorgi, Jeannerod, Villard, 2003], [B., Jeannerod, Schost, 2007, 2008], [Zhou, Labahn, 2012], [Jeannerod, Neiger, Villard, 2020], [Rosenkilde, Storjohann, 2016, 2021], etc.

▷ `gfun` (Maple), `Guess.m` (Mathematica), `ore_algebra` (SageMath), etc.

Algorithmic classification of models with D-finite $Q_{\mathcal{S}}(t) := Q_{\mathcal{S}}(1, 1; t)$

	OEIS	\mathcal{S}	Pol size	LDE size	Rec size		OEIS	\mathcal{S}	Pol size	LDE size	Rec size
1	A005566		—	(3, 4)	(2, 2)	13	A151275		—	(5, 24)	(9, 18)
2	A018224		—	(3, 5)	(2, 3)	14	A151314		—	(5, 24)	(9, 18)
3	A151312		—	(3, 8)	(4, 5)	15	A151255		—	(4, 16)	(6, 8)
4	A151331		—	(3, 6)	(3, 4)	16	A151287		—	(5, 19)	(7, 11)
5	A151266		—	(5, 16)	(7, 10)	17	A001006		(2, 2)	(2, 3)	(2, 1)
6	A151307		—	(5, 20)	(8, 15)	18	A129400		(2, 2)	(2, 3)	(2, 1)
7	A151291		—	(5, 15)	(6, 10)	19	A005558		—	(3, 5)	(2, 3)
8	A151326		—	(5, 18)	(7, 14)						
9	A151302		—	(5, 24)	(9, 18)	20	A151265		(6, 8)	(4, 9)	(6, 4)
10	A151329		—	(5, 24)	(9, 18)	21	A151278		(6, 8)	(4, 12)	(7, 4)
11	A151261		—	(4, 15)	(5, 8)	22	A151323		(4, 4)	(2, 3)	(2, 1)
12	A151297		—	(5, 18)	(7, 11)	23	A060900		(8, 9)	(3, 5)	(2, 3)

Equation sizes = (order, degree)

▷ Computerized discovery: enumeration + guessing [B., Kauers, 2009]

Algorithmic classification of models with D-finite $Q_{\mathcal{S}}(t) := Q_{\mathcal{S}}(1, 1; t)$

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Equation sizes = (order, degree)

- ▷ Computerized discovery: enumeration + guessing [B., Kauers, 2009]
- ▷ 1–22: DF confirmed by human proofs in [Bousquet-Mélou, Mishna, 2010]
- ▷ 23: DF confirmed by a human proof in [B., Kurkova, Raschel, 2017]
- ▷ All: explicit eqs. proved via CA [B., Chyzak, van Hoeij, Kauers, Pech, 2017]

Algorithmic classification of models with D-finite $Q_{\mathcal{S}}(t) := Q_{\mathcal{S}}(1, 1; t)$

	OEIS	\mathcal{S}	algebraic?	asymptotics		OEIS	\mathcal{S}	algebraic?	asymptotics
1	A005566		N	$\frac{4}{\pi} \frac{4^n}{n}$	13	A151275		N	$\frac{12\sqrt{30}}{\pi} \frac{(2\sqrt{6})^n}{n^2}$
2	A018224		N	$\frac{2}{\pi} \frac{4^n}{n}$	14	A151314		N	$\frac{\sqrt{6}\lambda\mu C^{5/2}}{5\pi} \frac{(2C)^n}{n^2}$
3	A151312		N	$\frac{\sqrt{6}}{\pi} \frac{6^n}{n}$	15	A151255		N	$\frac{24\sqrt{2}}{\pi} \frac{(2\sqrt{2})^n}{n^2}$
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5	A151266		N	$\frac{1}{2} \sqrt{\frac{3}{\pi}} \frac{3^n}{n^{1/2}}$	17	A001006		Y	$\frac{3}{2} \sqrt{\frac{3}{\pi}} \frac{3^n}{n^{3/2}}$
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9	A151302		N	$\frac{1}{3} \sqrt{\frac{5}{2\pi}} \frac{5^n}{n^{1/2}}$	20	A151265		Y	$\frac{2\sqrt{2}}{\Gamma(1/4)} \frac{3^n}{n^{3/4}}$
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► Computerized discovery: convergence acceleration + LLL [B., Kauers, '09]

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- ▷ Computerized discovery: convergence acceleration + LLL [B., Kauers, '09]
- ▷ Asympt. confirmed by human proofs via ACSV in [Melczer, Wilson, 2016]
- ▷ Transcendence proofs via CA [B., Chyzak, van Hoeij, Kauers, Pech, 2017]

Theorem [B., Chyzak, van Hoeij, Kauers, Pech, 2017]

Let \mathcal{S} be one of the models 1–19. Then

- $Q_{\mathcal{S}}(x, y; t)$ is expressible using (integrals of) ${}_2F_1$ expressions.
- $Q_{\mathcal{S}}(x, y; t)$ is transcendental.

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Example (King walks in the quarter plane, [A151331](#))

$$Q_{\begin{smallmatrix} \uparrow \downarrow \\ \leftarrow \rightarrow \end{smallmatrix}}(t) = \frac{1}{t} \int_0^t \frac{1}{(1+4x)^3} \cdot {}_2F_1\left(\frac{3}{2} \mid \frac{3}{2} \mid \frac{16x(1+x)}{(1+4x)^2}\right) dx$$

$$= 1 + 3t + 18t^2 + 105t^3 + 684t^4 + 4550t^5 + 31340t^6 + 219555t^7 + \dots$$

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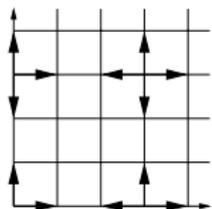
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- ▷ Computer-driven discovery and proof; no human proof yet.
- ▷ Proof uses: (1) **kernel method** and (2) **creative telescoping**
+ (3) **ODE factoring** and (4) **ODE solving**.

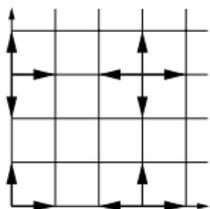


The kernel $K(x, y; t) := 1 - t \cdot \sum_{(i,j) \in \mathcal{S}} x^i y^j = 1 - t \left(x + \frac{1}{x} + y + \frac{1}{y} \right)$ is left invariant under the change of (x, y) into the elements of

$$\mathcal{G}_{\mathcal{S}} := \left\{ (x, y), \left(\frac{1}{x}, y \right), \left(\frac{1}{x}, \frac{1}{y} \right), \left(x, \frac{1}{y} \right) \right\}$$

Kernel equation:

$$\begin{aligned} K(x, y; t)xyQ(x, y; t) &= xy - txQ(x, 0; t) - tyQ(0, y; t) \\ -K(x, y; t)\frac{1}{x}yQ\left(\frac{1}{x}, y; t\right) &= -\frac{1}{x}y + t\frac{1}{x}Q\left(\frac{1}{x}, 0; t\right) + tyQ(0, y; t) \\ K(x, y; t)\frac{1}{x}\frac{1}{y}Q\left(\frac{1}{x}, \frac{1}{y}; t\right) &= \frac{1}{x}\frac{1}{y} - t\frac{1}{x}Q\left(\frac{1}{x}, 0; t\right) - t\frac{1}{y}Q\left(0, \frac{1}{y}; t\right) \end{aligned}$$

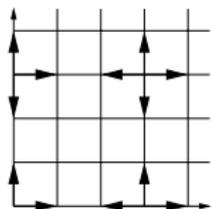


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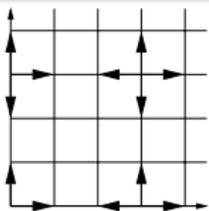
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Summing up and taking positive parts yields:

$$xyQ(x, y; t) = [x > y] \frac{xy - \frac{1}{x}y + \frac{1}{x}\frac{1}{y} - x\frac{1}{y}}{K(x, y; t)}$$



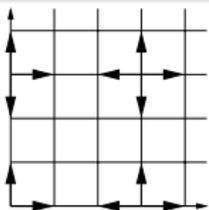
The kernel $K(x, y; t) := 1 - t \cdot \sum_{(i,j) \in \mathcal{S}} x^i y^j = 1 - t \left(x + \frac{1}{x} + y + \frac{1}{y} \right)$ is left invariant under the change of (x, y) into the elements of

$$\mathcal{G}_{\mathcal{S}} := \left\{ (x, y), \left(\frac{1}{x}, y \right), \left(\frac{1}{x}, \frac{1}{y} \right), \left(x, \frac{1}{y} \right) \right\}$$

Kernel equation:

$$\begin{aligned} K(x, y; t)xyQ(x, y; t) &= xy - txQ(x, 0; t) - tyQ(0, y; t) \\ -K(x, y; t)\frac{1}{x}yQ\left(\frac{1}{x}, y; t\right) &= -\frac{1}{x}y + t\frac{1}{x}Q\left(\frac{1}{x}, 0; t\right) + tyQ(0, y; t) \\ K(x, y; t)\frac{1}{x}\frac{1}{y}Q\left(\frac{1}{x}, \frac{1}{y}; t\right) &= \frac{1}{x}\frac{1}{y} - t\frac{1}{x}Q\left(\frac{1}{x}, 0; t\right) - t\frac{1}{y}Q\left(0, \frac{1}{y}; t\right) \\ -K(x, y; t)x\frac{1}{y}Q\left(x, \frac{1}{y}; t\right) &= -x\frac{1}{y} + txQ(x, 0; t) + t\frac{1}{y}Q\left(0, \frac{1}{y}; t\right) \end{aligned}$$

$$\text{GF} = \text{PosPart} \left(\frac{\text{OS}}{\text{Ker}} \right) \text{ is D-finite [Lipshitz, 1988]}$$



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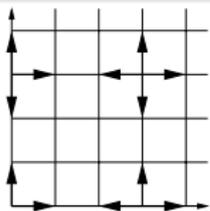
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$$\text{GF} = \text{PosPart} \left(\frac{\text{OS}}{\text{Ker}} \right) \text{ is D-finite [Lipshitz, 1988]}$$

▷ Argument works if $\text{OS} \neq 0$: algebraic version of the reflection principle



The kernel $K(x, y; t) := 1 - t \cdot \sum_{(i,j) \in \mathcal{S}} x^i y^j = 1 - t \left(x + \frac{1}{x} + y + \frac{1}{y} \right)$ is left invariant under the change of (x, y) into the elements of

$$\mathcal{G}_{\mathcal{S}} := \left\{ (x, y), \left(\frac{1}{x}, y \right), \left(\frac{1}{x}, \frac{1}{y} \right), \left(x, \frac{1}{y} \right) \right\}$$

Kernel equation:

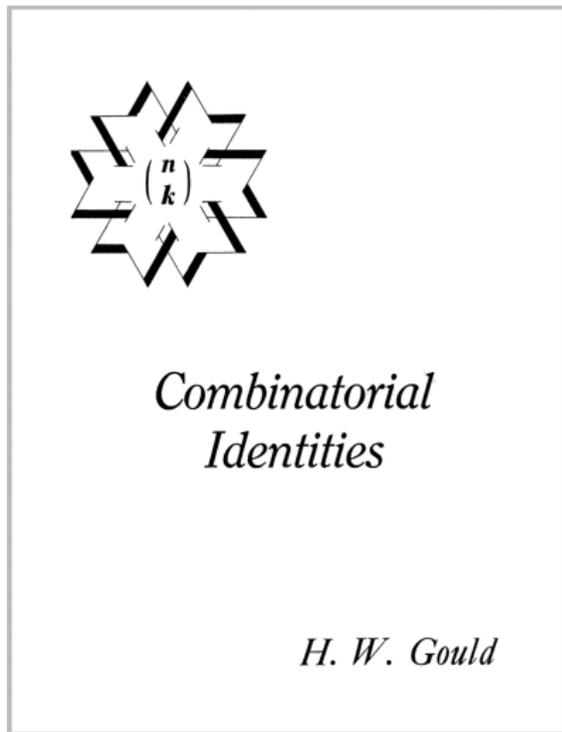
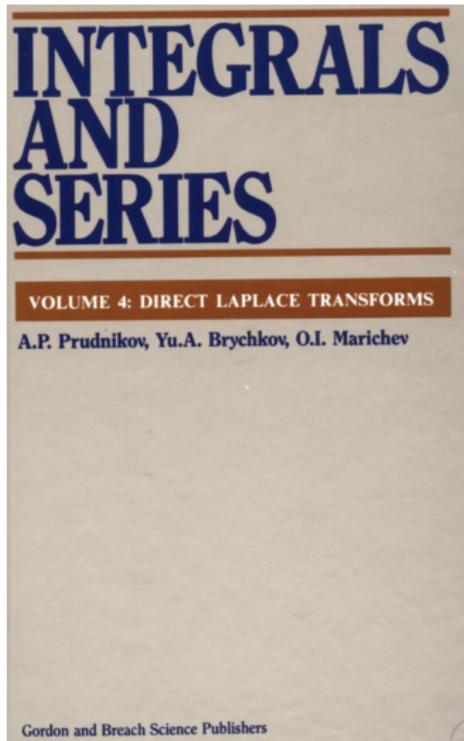
$$\begin{aligned} K(x, y; t)xyQ(x, y; t) &= xy - txQ(x, 0; t) - tyQ(0, y; t) \\ -K(x, y; t)\frac{1}{x}yQ\left(\frac{1}{x}, y; t\right) &= -\frac{1}{x}y + t\frac{1}{x}Q\left(\frac{1}{x}, 0; t\right) + tyQ(0, y; t) \\ K(x, y; t)\frac{1}{x}\frac{1}{y}Q\left(\frac{1}{x}, \frac{1}{y}; t\right) &= \frac{1}{x}\frac{1}{y} - t\frac{1}{x}Q\left(\frac{1}{x}, 0; t\right) - t\frac{1}{y}Q(0, \frac{1}{y}; t) \\ -K(x, y; t)x\frac{1}{y}Q\left(x, \frac{1}{y}; t\right) &= -x\frac{1}{y} + txQ(x, 0; t) + t\frac{1}{y}Q(0, \frac{1}{y}; t) \end{aligned}$$

$$\text{GF} = \text{PosPart} \left(\frac{\text{OS}}{\text{Ker}} \right) \text{ is D-finite [Lipshitz, 1988]}$$

▷ **Creative Telescoping** finds a differential equation for $\text{GF} = \int \text{RatFrac}$

(2) Creative Telescoping

“An algorithmic toolbox for multiple sums and integrals with parameters”



(2) Creative Telescoping

“An algorithmic toolbox for multiple sums and integrals with parameters”

DOUBLE INTEGRALS

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$$6. \int_0^{\infty} \int_0^{\infty} x^{2\alpha+1} y^{\alpha} - y^{\alpha} - \alpha x^{\alpha} \operatorname{ber}_{\alpha}(axy) \, dx \, dy = \left(\frac{ap}{2}\right)^{\alpha} \frac{4 \operatorname{Re}(\cos(3\pi\alpha/4)) - a^2 \sin(3\pi\alpha/4)}{16p^2 q^2 + a^4} \quad (\operatorname{Re} \alpha > -1)$$

$$7. \int_0^{\infty} \int_0^{\infty} x^{2\alpha+1} y^{\alpha} - \alpha x^{\alpha} - \alpha y^{\alpha} \operatorname{bei}_{\alpha}(axy) \, dx \, dy = (-1)^{\alpha} \sqrt{\frac{\pi}{2}} \frac{a^{\alpha\alpha}}{2^{2\alpha+1} p^{2\alpha+1} q^{2\alpha+1}} e^{-a^2/(2p^2 q^2)} D_{2\alpha+1} \left(\frac{a^2}{4p\sqrt{2}q} \right)$$

$$8. \int_0^{\infty} \int_0^{\infty} x^{2\alpha+1} y^{\alpha} - \alpha x^{\alpha} - \alpha y^{\alpha} \operatorname{ber}_{\alpha\alpha}(axy) \, dx \, dy = (-1)^{\alpha} \sqrt{\frac{\pi}{2}} \frac{a^{2\alpha}}{2^{2\alpha+1} p^{2\alpha+1} q^{2\alpha+1}} e^{-a^2/(2p^2 q^2)} D_{2\alpha} \left(\frac{a^2}{4p\sqrt{2}q} \right)$$

$$9. \int_0^{\infty} \int_0^{\infty} xy^{\alpha} - \alpha x^{\alpha} - \alpha y^{\alpha} \operatorname{bei}(axy) \, dx \, dy = \frac{\sqrt{\pi}}{2^{\alpha} p^{\alpha} q^{\alpha/2}} e^{-a^2/(2p^2 q^2)}$$

$$10. \int_0^{\infty} \int_0^{\infty} xy^{\alpha} - \alpha x^{\alpha} - \alpha y^{\alpha} \operatorname{ber}(axy) \, dx \, dy = \frac{\sqrt{\pi}}{8p\sqrt{q}} e^{-a^2/(2p^2 q^2)}$$

$$11. \int_0^{\infty} \int_0^{\infty} xy^{\alpha} - \alpha x^{\alpha} - \alpha y^{\alpha} [\operatorname{ber}^2(axy) + \operatorname{bei}^2(axy)] \, dx \, dy = \frac{1}{8p^2} e^{-a^2/(16p^2 q^2)}$$

$$12. \int_0^{\infty} \int_0^{\infty} xy^{\alpha} - \alpha x^{\alpha} - \alpha y^{\alpha} \operatorname{ber}(axy) \, dx \, dy = \frac{\pi}{16\sqrt{p}q} J_0 \left(\frac{a^2}{8\sqrt{p}q} \right)$$

$$13. \int_0^{\infty} \int_0^{\infty} xy K_{\alpha}(ax) K_{\beta}(by) (\operatorname{ber} x \operatorname{ber} y - \operatorname{bei} x \operatorname{bei} y) \, dx \, dy = \frac{ab-1}{(a^2+1)(b^2+1)} \quad (\operatorname{Re} a^{\alpha}, \operatorname{Re} b^{\beta} > 1/2)$$

3.2.12. Integrals containing orthogonal polynomials.

$$1. \int_0^1 \int_0^1 \frac{y^{\alpha\beta+1}}{(1-y^{\alpha})^{\beta+1}} \sin^{2\alpha} x P_n^{(\alpha, \beta)}(2) [\cos a \cos b + y^{\alpha} \sin a \sin b]^2 \, dx \, dy = \frac{\sqrt{\pi}}{2} \Gamma \left[\frac{\alpha-\beta}{\alpha+1}, \beta+1/2 \right] \frac{P_n^{(\alpha, \beta)}(\cos 2a) P_n^{(\alpha, \beta)}(\cos 2b)}{P_n^{(\alpha, \beta)}(1)} \quad [\operatorname{Re} \alpha > \beta > -1/2]$$

$$2. \int_{-1}^1 \int_{-1}^1 (1-x^2)^{\alpha-1} (1-y^2)^{\beta-1} C_n^{\alpha+\beta}(ax+by) \, dx \, dy = \pi \Gamma \left[\frac{\alpha}{\alpha+1/2}, \frac{\beta}{\beta+1/2} \right] \frac{C_n^{\alpha+\beta}(0) P_n^{(\alpha-1/2, \beta-1/2)}(A) P_n^{(\alpha-1/2, \beta-1/2)}(B)}{P_n^{(\alpha-1/2, \beta-1/2)}(-1) P_n^{(\alpha-1/2, \beta-1/2)}(1)} \quad [2\alpha = \sqrt{1-A}(1-B), 2\beta = \sqrt{1+A}(1+B); \operatorname{Re} \alpha, \operatorname{Re} \beta > 0]$$

$$3. \int_0^1 \int_0^1 \frac{P_n(x) P_n(y)}{\sqrt{1-x^2} \sqrt{1-y^2}} \, dx \, dy = \begin{cases} (-1)^{n/2} 2^n \frac{(n-1)!}{n! (2n+1)} & [n = n \text{ even}] \\ 0 & [\text{otherwise}] \end{cases}$$

$$7. \sum_{k=0}^{2n} (-1)^k \binom{2n}{k} \binom{2k}{k} \frac{1}{\binom{n+k}{k} 2^{2k}} = \frac{\binom{6n}{2n}}{\binom{2n}{n}} 2^{-4n}$$

$$8. \sum_{k=0}^{2n} (-1)^k \binom{2n}{k} \binom{2k}{k} \frac{2k+1}{\binom{n+k}{k} (n+k+1)} = 1$$

$$9. \sum_{k=0}^n (-1)^k \binom{n}{k} \binom{x+k}{k} \frac{2^{2k}}{\binom{2k}{k} (2k+1)} = \frac{2^{2n}}{\binom{2n}{n} (2n+1)} \binom{n-x-\frac{1}{2}}{n}$$

$$10. \sum_{k=0}^n (-1)^k \binom{n}{k} \binom{n+k}{k} \frac{2^{2k}}{\binom{2k}{k} (2k+1)} = \frac{(-1)^n}{2n+1}$$

$$11. \sum_{k=0}^n (-1)^k \binom{n}{k} \binom{x+k}{k} \frac{2^{2k}}{\binom{2k}{k} (x+k)} = (-1)^n \frac{\binom{2x}{2n}}{x \binom{x}{n}}$$

$$12. \sum_{k=0}^n (-1)^k \binom{n}{k} \binom{n+k}{k} \frac{2^{2k}}{\binom{2k}{k} (n+k)} = \frac{(-1)^n}{n}, (n \geq 1)$$

$$13. \sum_{k=0}^n \binom{n}{k} \binom{-n-1}{n-k} \frac{1}{\binom{k+r}{k}} = (-1)^n r \sum_{k=0}^n \binom{n}{k}^2 \frac{1}{k+r}$$

(2) Creative Telescoping

“An algorithmic toolbox for multiple sums and integrals with parameters”

Example [Apéry 1978]: $A_n = \sum_{k=0}^n \binom{n}{k}^2 \binom{n+k}{k}^2$ satisfies the recurrence

$$(n+1)^3 A_{n+1} + n^3 A_{n-1} = (2n+1)(17n^2 + 17n + 5)A_n.$$

▷ Key fact used to prove that $\zeta(3) := \sum_{n \geq 1} \frac{1}{n^3} \approx 1.202056903 \dots$ is irrational.

1. Journées Arithmétiques de Marseille-Luminy, June 1978

The board of programme changes informed us that R. Apéry (Caen) would speak Thursday, 14.00 “Sur l’irrationalité de $\zeta(3)$.” Though there had been earlier rumours of his claiming a proof, scepticism was general. The lecture tended to strengthen this view to rank disbelief. Those who listened casually, or who were afflicted with being non-Francophone, appeared to hear only a sequence of unlikely assertions.

7. ICM '78, Helsinki, August 1978

Neither Cohen nor I had been able to prove (5) or (5) in the intervening 2 months. After a few days of fruitless effort the specific problem was mentioned to Don Zagier (Bonn), and with irritating speed he showed that indeed the sequence $\{b'_n\}$ satisfies the recurrence (4). This more or less broke the dam and (5) and (5) were quickly conquered. Henri Cohen addressed a very well-attended meeting at 17.00 on Friday, August 18 in the language of the majority, proving (5) and explaining how this implied the

[Van der Poorten, 1979: “A proof that Euler missed”]

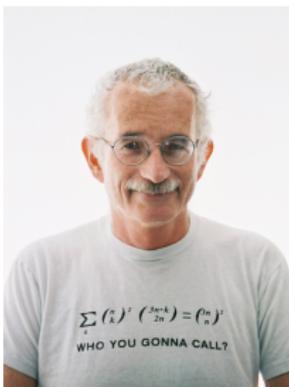
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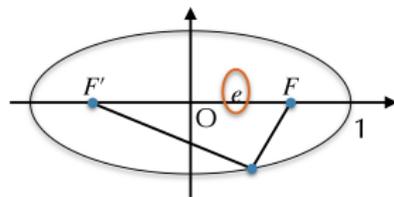
[Zeilberger, 1990: “The method of creative telescoping”]

(2) Creative Telescoping

“An algorithmic toolbox for multiple sums and integrals with parameters”

Example [Euler, 1733]: Perimeter of an ellipse of eccentricity e , semi-major axis 1

$$p(e) = 4 \int_0^1 \sqrt{\frac{1 - e^2 u^2}{1 - u^2}} du = 4 \oint \frac{dx dy}{1 - \frac{1 - e^2 x^2}{(1 - x^2)y^2}}$$



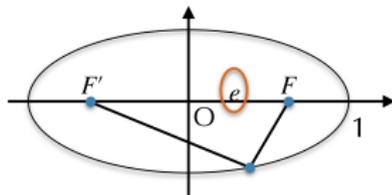
Principle: Find algorithmically

(2) Creative Telescoping

“An algorithmic toolbox for multiple sums and integrals with parameters”

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Principle: Find algorithmically

$$\begin{aligned} & \left((e - e^3) \partial_e^2 + (1 - e^2) \partial_e + e \right) \left(\frac{1}{1 - \frac{1-e^2x^2}{(1-x^2)y^2}} \right) = \\ & \partial_x \left(\frac{e(1+x-x^2-x^3)y^2(2x-3+y^2+x^2(3e^2-y^2-2))}{(y^2+x^2(e^2-y^2)-1)^2} \right) \\ & \quad + \partial_y \left(\frac{2e(e^2-1)x(1+x^3)y^3}{(y^2+x^2(e^2-y^2)-1)^2} \right) \end{aligned}$$

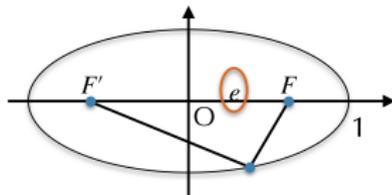
▷ Conclusion: $(e - e^3) \cdot p''(e) + (1 - e^2) \cdot p'(e) + e \cdot p(e) = 0.$

(2) Creative Telescoping

“An algorithmic toolbox for multiple sums and integrals with parameters”

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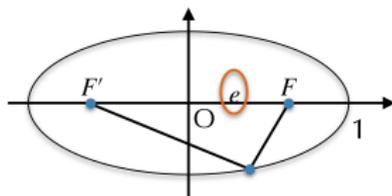
▷ Conclusion: $p(e) = \frac{\pi}{2} \cdot {}_2F_1 \left(-\frac{1}{2}, \frac{1}{2} \mid e^2 \right) = 2\pi - \frac{\pi}{2}e^2 - \frac{3\pi}{32}e^4 - \dots$

(2) Creative Telescoping

“An algorithmic toolbox for multiple sums and integrals with parameters”

Example [Euler, 1733]: Perimeter of an ellipse of eccentricity e , semi-major axis 1

$$p(e) = 4 \int_0^1 \sqrt{\frac{1-e^2u^2}{1-u^2}} du = 4 \not\int \frac{dx dy}{1 - \frac{1-e^2x^2}{(1-x^2)y^2}}$$



Principle: Find algorithmically

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▷ Drawback: Size(certificates) \gg Size(telescopers).

(2) Creative Telescoping: several generations of algorithms

- 1G, *elimination-based*: [Fasenmyer, 1947], [Lipshitz, 1988], [Zeilberger, 1990], [Takayama, 1990], [Wilf, Zeilberger, 1990], [Chyzak, Salvy, 2000]
 - 2G, *linear diff/rec rational solving*: [Zeilberger, 1990], [Zeilberger, 1991], [Almkvist, Zeilberger, 1990], [Chyzak, 2000], [Koutschan, 2010]
 - 3G, combines $1G + 2G + \text{linear algebra}$: [Apagodu, Zeilberger, 2005], [Koutschan 2010], [Chen, Kauers 2012], [Chen, Kauers, Koutschan 2014]
- ▷ Advantages:
- 1G–3G: very general algorithms;
 - 2G/3G algorithms are able to solve non-trivial applications.
- ▷ Drawbacks:
- 1G: slow;
 - 2G: bad or unknown complexity;
 - 1G and 3G: non-minimality of telescopers;
 - 1G–3G: all compute (big) certificates.

(2) Creative Telescoping: several generations of algorithms

4G: roots in [Ostrogradsky, 1845], [Hermite, 1872] and [Picard, 1902]

- univariate:

- rational f : [B., Chen, Chyzak, Li, 2010]
- hyperexponential f : [B., Chen, Chyzak, Li, Xin, 2013]
- hypergeometric Σ : [Chen, Huang, Kauers, Li, 2015], [Huang, 2016]
- mixed $f + \Sigma$: [B., Dumont, Salvy, 2016]
- algebraic f : [Chen, Kauers, Koutschan, 2016]
- D-finite Fuchsian f : [Chen, van Hoeij, Kauers, Koutschan, 2018]
- D-finite f : [B., Chyzak, Lairez, Salvy, 2018], [van der Hoeven, 2018]

- multiple:

- rational bivariate \mathcal{H} : [Chen, Kauers, Singer, 2012], [Chen, Du, Kauers, 2021]
- rational: [B., Lairez, Salvy, 2013], [Lairez 2016]
- binomial sums: [B., Lairez, Salvy, 2017]

▷ Advantages:

- good complexity;
- minimality of telescopers;
- do not need to compute certificates;
- fast in practice.

▷ Drawback: not (yet) as general as 1G–3G algorithms.

(2) 4G Creative Telescoping

Algorithm for the integration of rational functions [B., Lairez, Salvy, 2013]

- **Input:** $R(e, \mathbf{x})$ a rational function in e and $\mathbf{x} = x_1, \dots, x_n$.
- **Output:** A linear ODE $T(e, \partial_e)y = 0$ satisfied by $y(e) = \int R(e, \mathbf{x})dx$.
- **Complexity:** $\mathcal{O}(D^{8n+2})$, where $D = \deg R$.
- **Output size:** T has order $\leq D^n$ in ∂_e and degree $\leq D^{3n+2}$ in e .

- ▷ Roots in [B., Chen, Chyzak, Li, 2010] ($n = 1$).
- ▷ Relies on **generalized Hermite reduction** and **polynomial linear algebra**.
- ▷ Avoids the (costly) computation of **certificates**, of size $\Omega(D^{n^2/2})$.
- ▷ Previous algorithms: complexity (at least) doubly exponential in n .
- ▷ Highly non-trivial extension by [Lairez, 2016]: very efficient in practice.

Models 1–19: explicit expressions and transcendence

Theorem [B., Chyzak, van Hoeij, Kauers, Pech, 2017]

Let \mathcal{S} be one of the models 1–19. Then

- $Q_{\mathcal{S}}(t)$ is expressible using (integrals of) ${}_2F_1$ expressions.
- $Q_{\mathcal{S}}(t)$ is transcendental, except for $\mathcal{S} = \begin{smallmatrix} \uparrow \\ \downarrow \end{smallmatrix}$ and $\mathcal{S} = \begin{smallmatrix} \uparrow \downarrow \\ \downarrow \uparrow \end{smallmatrix}$.

Example (King walks in the quarter plane, A151331)

$$Q_{\begin{smallmatrix} \uparrow \downarrow \\ \downarrow \uparrow \end{smallmatrix}}(t) = \frac{1}{t} \int_0^t \frac{1}{(1+4x)^3} \cdot {}_2F_1\left(\frac{3}{2} \quad \frac{3}{2} \mid \frac{16x(1+x)}{(1+4x)^2}\right) dx$$
$$= 1 + 3t + 18t^2 + 105t^3 + 684t^4 + 4550t^5 + 31340t^6 + 219555t^7 + \dots$$

- ▷ Computer-driven discovery and proof; no human proof yet.
- ▷ Proof uses: (1) **kernel method** and (2) **creative telescoping**
+ (3) **ODE factoring** and (4) **ODE solving**.

Theorem (Apéry's power series is transcendental)

$$f(t) = \sum_n A_n t^n, \quad \text{where } A_n = \sum_{k=0}^n \binom{n}{k}^2 \binom{n+k}{k}^2, \quad \text{is transcendental.}$$

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Proof:

① Creative telescoping:

$$(n+1)^3 A_{n+1} + n^3 A_{n-1} = (2n+1)(17n^2 + 17n + 5)A_n, \quad A_0 = 1, A_1 = 5$$

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② **Conversion** from recurrence to differential equation $L(f) = 0$, where

$$L = (t^4 - 34t^3 + t^2)\partial_t^3 + (6t^3 - 153t^2 + 3t)\partial_t^2 + (7t^2 - 112t + 1)\partial_t + t - 5$$

Theorem (Apéry's power series is transcendental)

$$f(t) = \sum_n A_n t^n, \quad \text{where } A_n = \sum_{k=0}^n \binom{n}{k}^2 \binom{n+k}{k}^2, \quad \text{is transcendental.}$$

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$$\left\{ 1 + 5t + O(t^2), \ln(t) + (5\ln(t) + 12)t + O(t^2), \ln(t)^2 + (5\ln(t)^2 + 24\ln(t))t + O(t^2) \right\}$$

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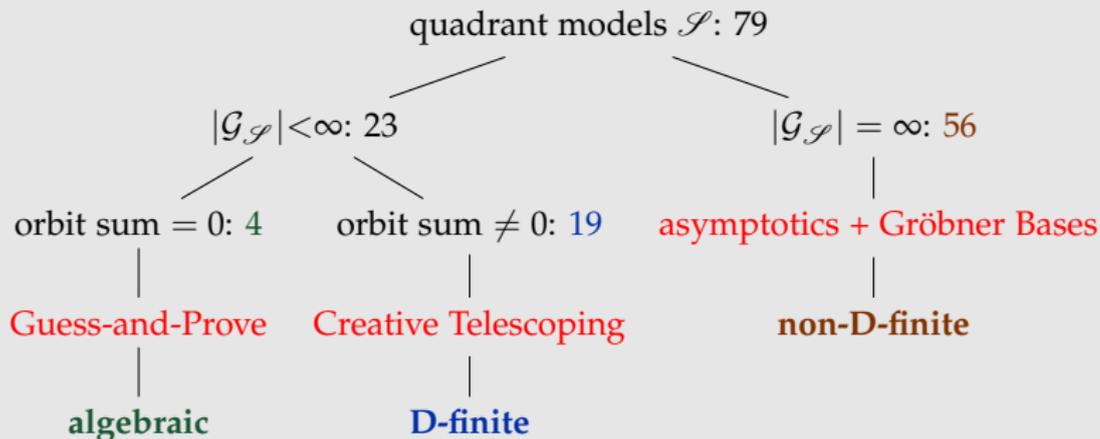
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⑤ **Conclusion:** f is transcendental[†]

[†] f algebraic would imply a full basis of algebraic solutions for L_f^{\min} [Tannery, 1875].

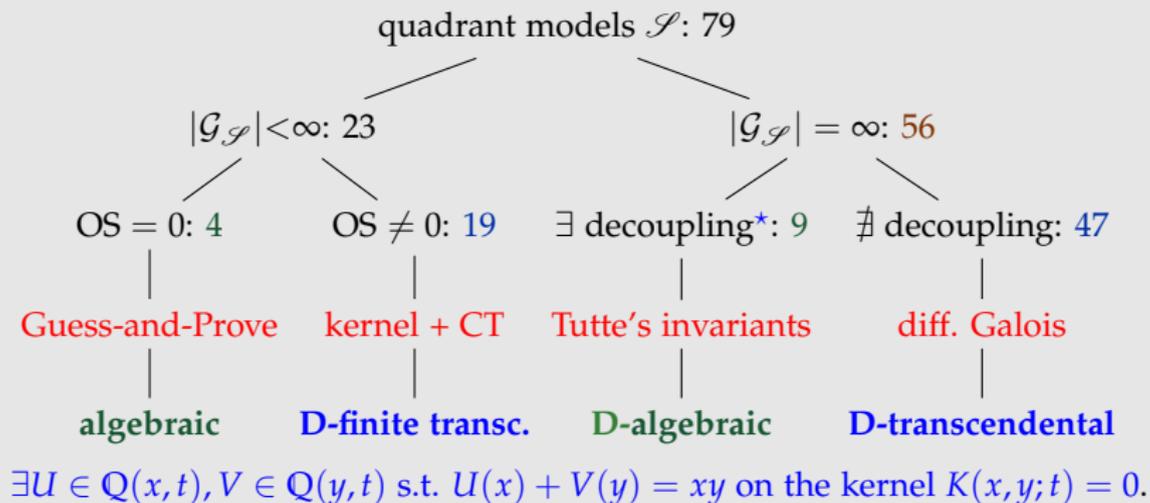
Summary: classification of walks with small steps in \mathbb{N}^2

$Q_{\mathcal{S}}$ is D-finite \iff a certain group $\mathcal{G}_{\mathcal{S}}$ is finite (!)



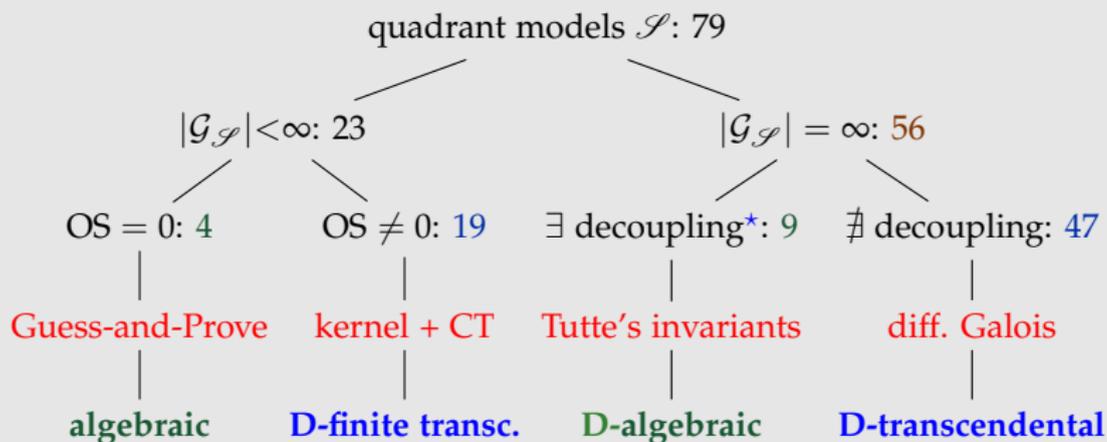
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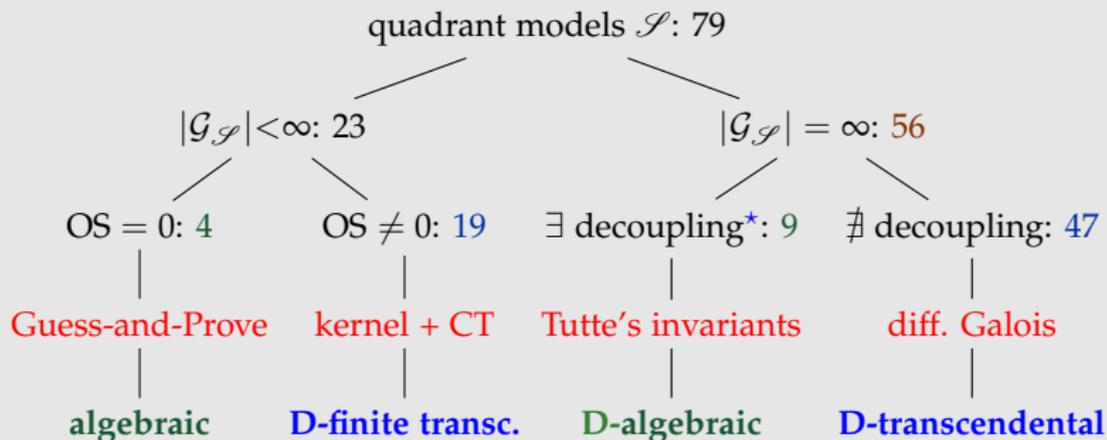
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▷ Many contributors (2010–2021): Bernardi, B., Bousquet-Mélou, Chyzak, Dreyfus, Hardouin, van Hoeij, Kauers, Kurkova, Mishna, Pech, Raschel, Roques, Salvy, Singer

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▷ Proofs use **various tools**: algebra, complex analysis, probability theory, differential Galois theory, computer algebra, etc.



Enumerative Combinatorics and Computer Algebra enrich one another



Classification of $Q(x, y; t)$ **fully completed** for 2D small-steps walks



Robust algorithmic methods, based on efficient algorithms:

- **Guess-and-Prove**
- **Creative Telescoping**



Brute-force and/or use of naive algorithms = **hopeless**.

E.g. size of algebraic equations for $G(x, y; t) \approx 30\text{Gb}$.



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Lack of “purely human” proofs for some results.



Many beautiful open questions for 2D walk models with **repeated** or **large** steps, and in different cones, and in **dimension** > 2 .

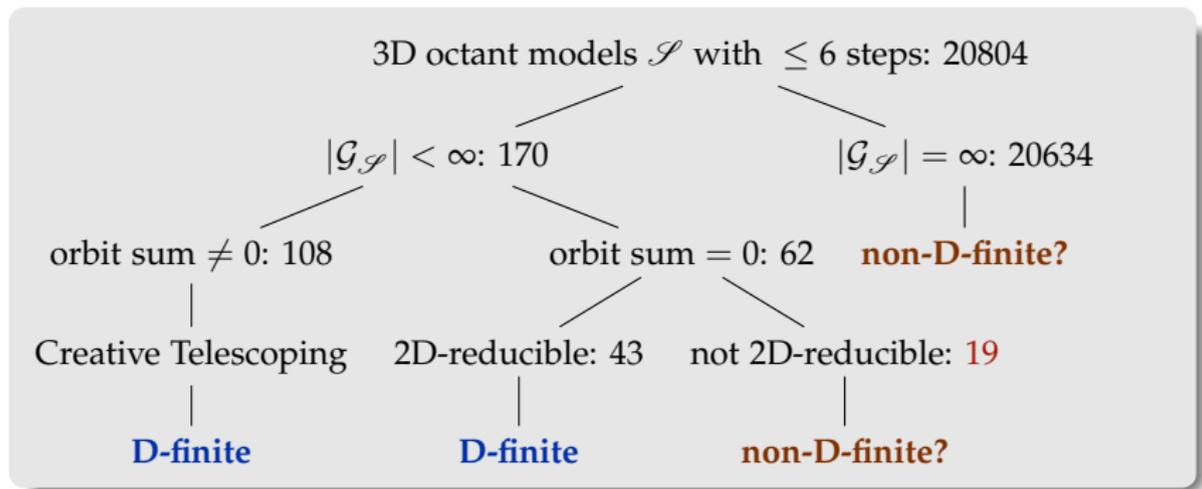
Thanks for your attention!

- Automatic classification of restricted lattice walks, with M. Kauers. *Proceedings of FPSAC*, 2009.
- The complete generating function for Gessel walks is algebraic, with M. Kauers. *Proceedings of the American Mathematical Society*, 2010.
- Explicit formula for the generating series of diagonal 3D Rook paths, with F. Chyzak, M. van Hoeij and L. Pech. *Séminaire Lotharingien de Combinatoire*, 2011.
- Non-D-finite excursions in the quarter plane, with K. Raschel and B. Salvy. *Journal of Combinatorial Theory A*, 2014.
- On 3-dimensional lattice walks confined to the positive octant, with M. Bousquet-Mélou, M. Kauers and S. Melczer. *Annals of Comb.*, 2016.
- A human proof of Gessel's lattice path conjecture, with I. Kurkova, K. Raschel, *Transactions of the American Mathematical Society*, 2017.
- Hypergeometric expressions for generating functions of walks with small steps in the quarter plane, with F. Chyzak, M. van Hoeij, M. Kauers and L. Pech, *European Journal of Combinatorics*, 2017.
- Counting walks with large steps in an orthant, with M. Bousquet-Mélou and S. Melczer, *Journal of the European Mathematical Society*, 2021.
- Computer Algebra in the Service of Enumerative Combinatorics, *Proceedings of ISSAC'21*, 2021.

Bonus

Beyond dimension 2: walks with small-steps in \mathbb{N}^3

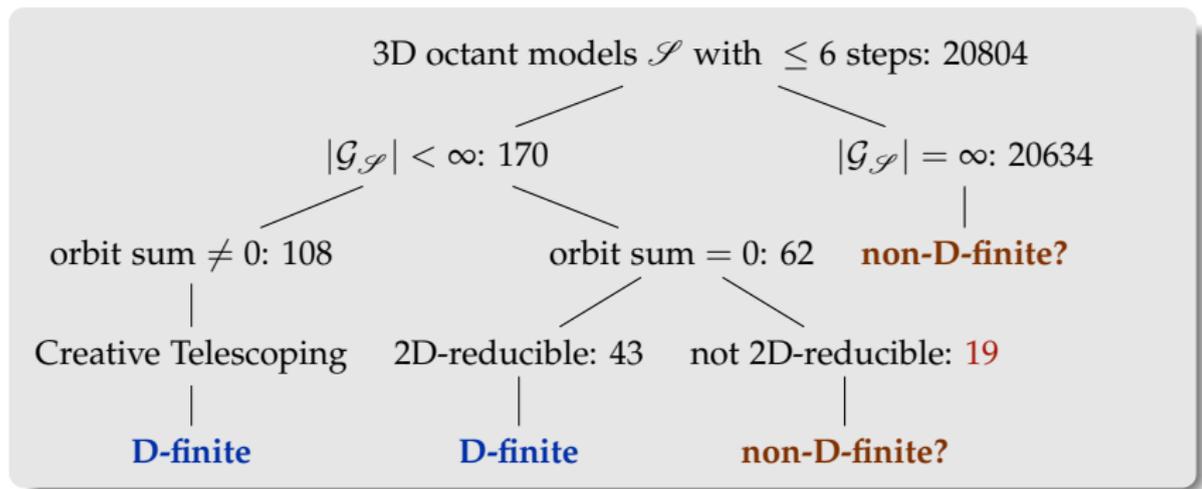
▷ $2^{3^3-1} \approx 67$ million models, of which ≈ 11 million inherently 3D



[B., Bousquet-Mélou, Kauers, Melczer, 2016] + [Du, Hou, Wang, 2017];
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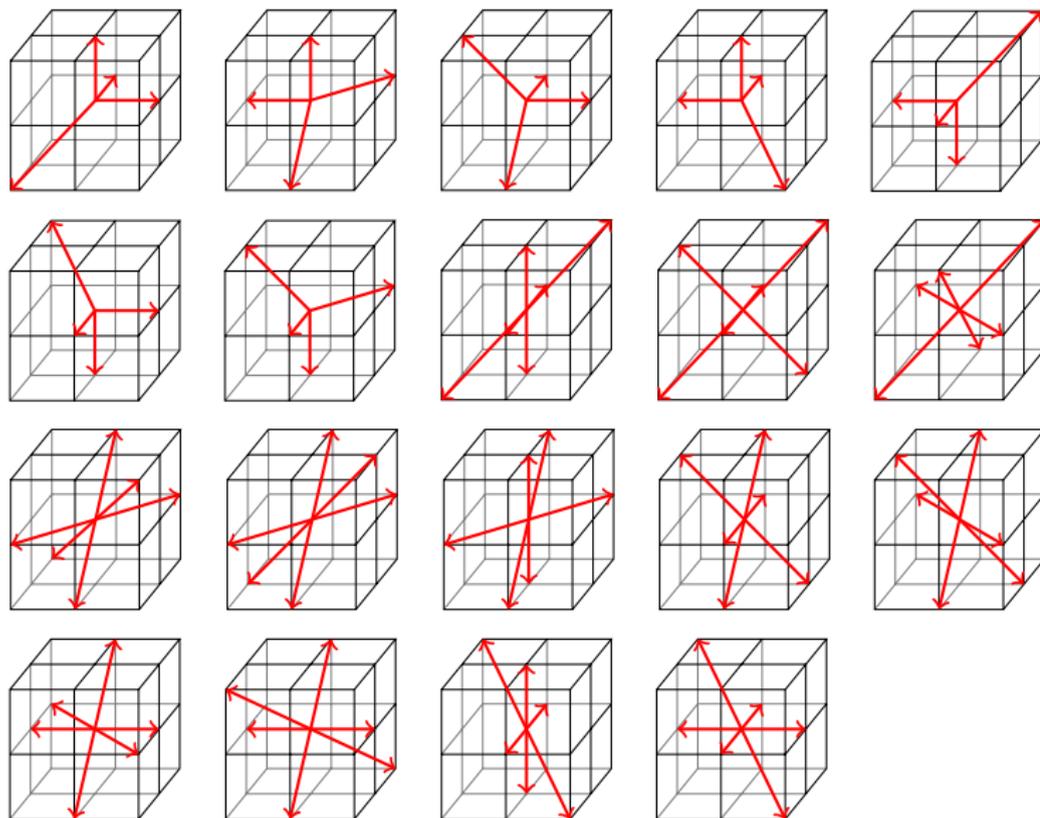


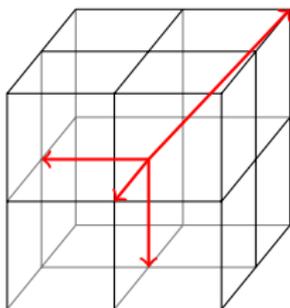
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Question: differential finiteness \iff finiteness of the group?

Answer: probably no

19 mysterious 3D-models: finite $\mathcal{G}_{\mathcal{S}}$ and possibly non-D-finite $Q_{\mathcal{S}}$



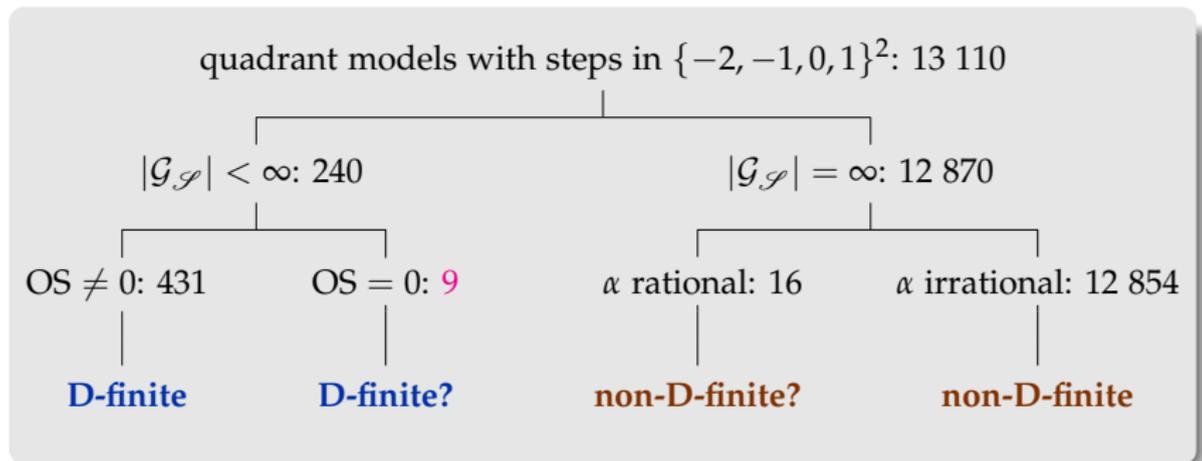


Numerical computations [Dahne, Salvy, 2020] suggest:

$$k_{4n} = C \cdot 256^n / n^\alpha, \text{ for } \alpha = 3.3257570041744 \dots \notin \mathbb{Q},$$

so excursions are very probably non-D-finite

Beyond small steps: Walks in \mathbb{N}^2 with large steps



[B., Bousquet-Mélou, Melczer, 2021]

Question: differential finiteness \iff finiteness of the group?

Answer: ?

Two challenging models with large steps

Conjecture 1 [B., Bousquet-Mélou, Melczer, 2021]

For the model  the excursions generating function $Q(0,0;t^{1/2})$ equals

$$\frac{1}{3t} - \frac{1}{6t} \cdot \left(\frac{1-12t}{(1+36t)^{1/3}} \cdot {}_2F_1 \left(\frac{1}{6}, \frac{2}{3} \mid \frac{108t(1+4t)^2}{(1+36t)^2} \right) + \sqrt{1-12t} \cdot {}_2F_1 \left(-\frac{1}{6}, \frac{2}{3} \mid \frac{108t(1+4t)^2}{(1-12t)^2} \right) \right).$$

Conjecture 2 [B., Bousquet-Mélou, Melczer, 2021]

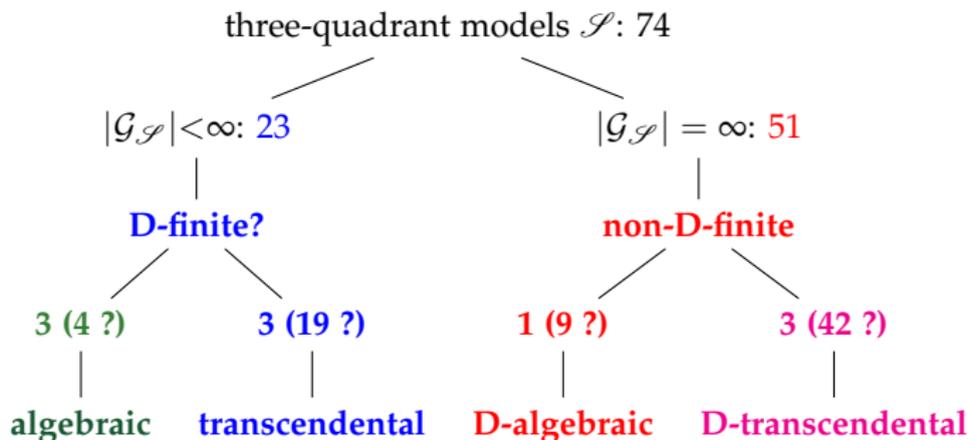
For the model  the excursions generating function $Q(0,0;t)$ equals

$$\frac{(1-24U+120U^2-144U^3)(1-4U)}{(1-3U)(1-2U)^{3/2}(1-6U)^{9/2}},$$

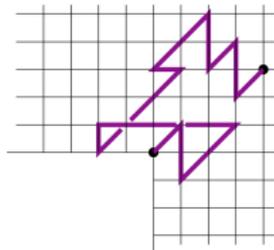
where $U = t^4 + 53t^8 + 4363t^{12} + \dots$ is the unique series in $\mathbb{Q}[[t]]$ satisfying

$$U(1-2U)^3(1-3U)^3(1-6U)^9 = t^4(1-4U)^4.$$

Beyond the first quadrant: three-quadrant walks with small steps



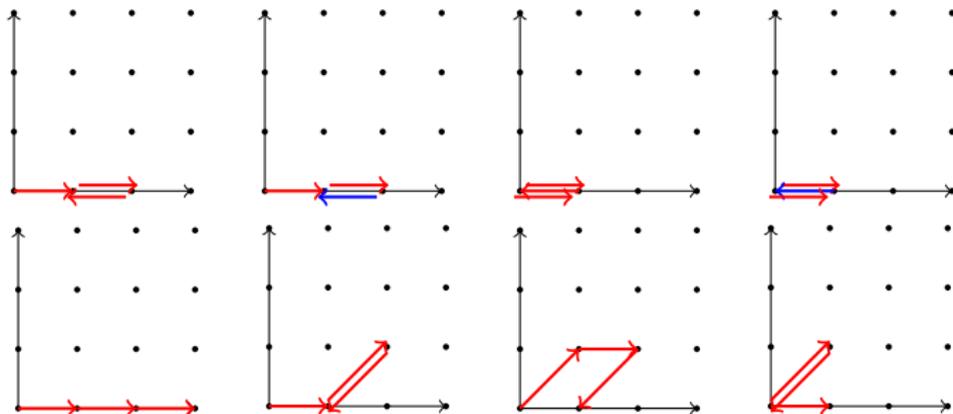
- ▷ Partial classification due to
[Bousquet-Mélou, 2016], [Raschel, Trotignon, 2019],
[Mustapha, 2019], [Dreyfus, Trotignon, 2020],
[Bousquet-Mélou, Wallner, 2021], [Bousquet-Mélou, 2021]



A difficult quadrant model with repeated steps

Theorem [B., Bousquet-Mélou, Kauers, Melczer, 2016]

Let $a_n = \# \left\{ \begin{array}{c} \bullet \\ \bullet \end{array} \begin{array}{c} \bullet \\ \bullet \end{array} \right\} - \text{walks of length } n \text{ in } \mathbb{N}^2 \text{ from } (0,0) \text{ to } (\star,0) \left. \vphantom{\left\{ \right.} \right\}$. Then $f(t) = \sum_n a_n t^n = 1 + t + 4t^2 + 8t^3 + 39t^4 + 98t^5 + \dots$ is transcendental.



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Proof:

- ① Discover and certify a differential equation L for $f(t)$ of order 11 and degree 73 high-tech Guess-and-Prove
- ② If $\text{ord}(L_f^{\min}) \leq 10$, then $\text{deg}_t(L_f^{\min}) \leq 580$ apparent singularities
- ③ Rule out this possibility differential Hermite-Padé approximants
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- ▷ General minimization algorithm and application to transcendence
[B., Rivoal, Salvy, 2021]

Solution of the “exercise”

- The kernel equation reads (with $K(x, y) = 1 - t(y + \bar{x} + x\bar{y})$):

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$$(27t^4 - t)A''(t) + (108t^3 - 4)A'(t) + 54t^2A(t) = 0.$$

> Zeilberger(1/x * sqrt((t-x)^2 - 4*t^2*x^3)/(2*t^2*x^2), t, x, Dt);

The group of the model $\{\uparrow, \leftarrow, \searrow\}$

Step set $\mathcal{S} = \{(-1,0), (0,1), (1,-1)\}$, with **characteristic polynomial**

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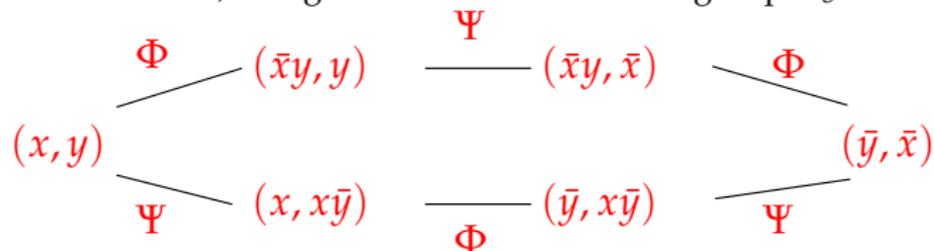
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Φ and Ψ are involutions, and generate a finite dihedral group D_3 of order 6:



- Orbit equation:

$$\begin{aligned} &xyQ(x, y) - \bar{x}y^2Q(\bar{x}y, y) + \bar{x}^2yQ(\bar{x}y, \bar{x}) \\ &\quad - \bar{x}\bar{y}Q(\bar{y}, \bar{x}) + x\bar{y}^2Q(\bar{y}, x\bar{y}) - x^2\bar{y}Q(x, x\bar{y}) = \\ &\quad \frac{xy - \bar{x}y^2 + \bar{x}^2y - \bar{x}\bar{y} + x\bar{y}^2 - x^2\bar{y}}{1 - t(y + \bar{x} + x\bar{y})} \end{aligned}$$

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- **Corollary** [Bousquet-Mélou & Mishna, 2010]:

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$$B(t) = [z^0]Q(z, \bar{z}) = [u^{-1}v^{-1}z^{-1}] \frac{\bar{u}\bar{v} - u\bar{v}^2 + u^2\bar{v} - uv + \bar{u}v^2 - \bar{u}^2v}{z(1 - zu)(1 - v\bar{z})(1 - t(\bar{v} + u + \bar{u}v))}$$

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- **Creative Telescoping** gives a differential equation for $B(t)$:

$$(27t^4 - t)B''(t) + (108t^3 - 4)B'(t) + 54t^2B(t) = 0.$$

We have proved that $A(t)$ and $B(t)$ are both solutions of

$$(27t^4 - t)y''(t) + (108t^3 - 4)y'(t) + 54t^2y(t) = 0.$$

Solving this equation proves:

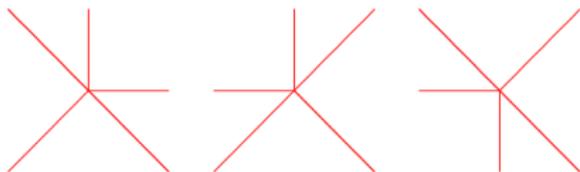
$$A(t) = B(t) = {}_2F_1\left(\begin{matrix} 1/3 & 2/3 \\ 2 \end{matrix} \middle| 27t^3\right) = \sum_{n=0}^{\infty} \frac{(3n)!}{n!^3} \frac{t^{3n}}{n+1}.$$

Thus the two sequences are equal to

$$a_{3n} = b_{3n} = \frac{(3n)!}{n!^2 \cdot (n+1)!}, \quad \text{and} \quad a_m = b_m = 0 \quad \text{if } 3 \text{ does not divide } m.$$

Example with infinite group: the scarecrows

[B., Raschel, Salvy, 2014]: $Q_{\mathcal{S}}(0,0;t)$ is not D-finite for the models



▷ For the 1st and the 3rd, the excursions sequence $[t^n] Q_{\mathcal{S}}(0,0;t)$

$1, 0, 0, 2, 4, 8, 28, 108, 372, \dots$

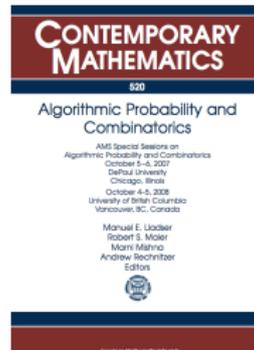
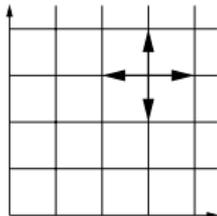
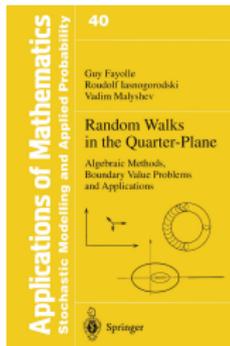
is $\sim K \cdot 5^n \cdot n^{-\alpha}$, with $\alpha = 1 + \pi / \arccos(1/4) = 3.383396\dots$

[Denisov, Wachtel, 2015]

▷ The **irrationality** of α prevents $Q_{\mathcal{S}}(0,0;t)$ from being D-finite.

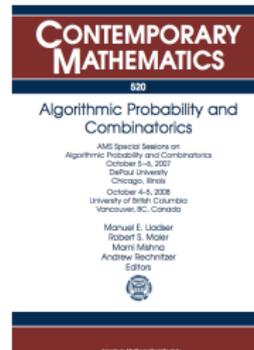
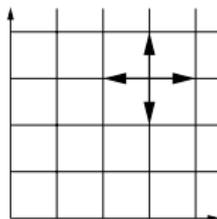
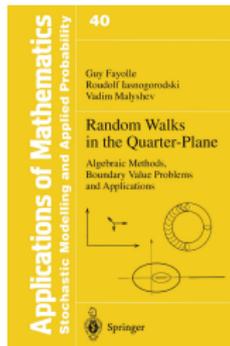
[Katz, 1970; Chudnovsky, 1985; André, 1989]

The group of a model: the simple walk case



The characteristic polynomial $\chi_{\mathcal{L}} := x + \frac{1}{x} + y + \frac{1}{y}$

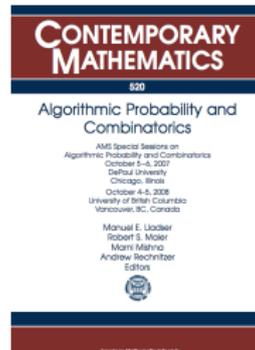
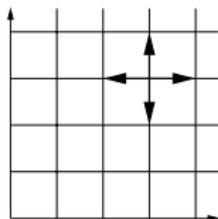
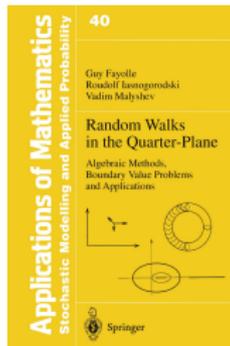
The group of a model: the simple walk case



The characteristic polynomial $\chi_{\mathcal{L}} := x + \frac{1}{x} + y + \frac{1}{y}$ is left invariant under

$$\psi(x, y) = \left(x, \frac{1}{y}\right), \quad \phi(x, y) = \left(\frac{1}{x}, y\right),$$

The group of a model: the simple walk case



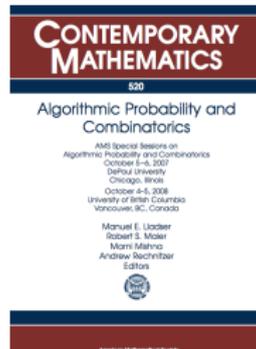
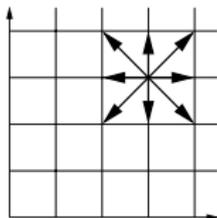
The characteristic polynomial $\chi_{\mathcal{L}} := x + \frac{1}{x} + y + \frac{1}{y}$ is left invariant under

$$\psi(x, y) = \left(x, \frac{1}{y}\right), \quad \phi(x, y) = \left(\frac{1}{x}, y\right),$$

and thus under any element of the group

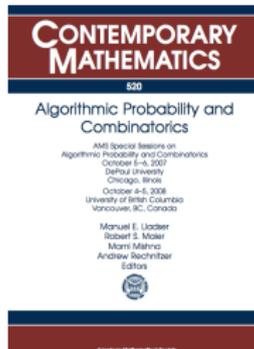
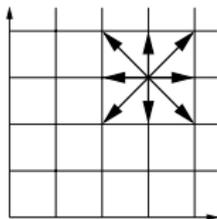
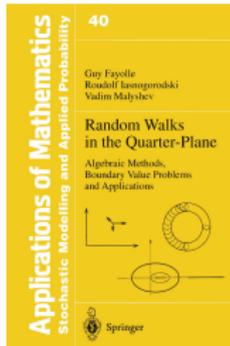
$$\langle \psi, \phi \rangle = \left\{ (x, y), \left(x, \frac{1}{y}\right), \left(\frac{1}{x}, \frac{1}{y}\right), \left(\frac{1}{x}, y\right) \right\}.$$

The group of a model



The generating polynomial $\chi_{\mathcal{S}} := \sum_{(i,j) \in \mathcal{S}} x^i y^j = \sum_{i=-1}^1 B_i(y) x^i = \sum_{j=-1}^1 A_j(x) y^j$

The group of a model



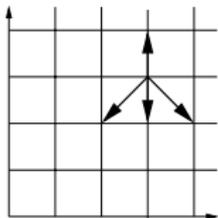
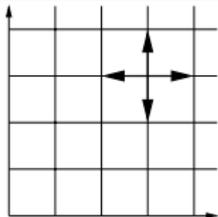
The generating polynomial $\chi_{\mathcal{S}} := \sum_{(i,j) \in \mathcal{S}} x^i y^j = \sum_{i=-1}^1 B_i(y) x^i = \sum_{j=-1}^1 A_j(x) y^j$
is left invariant under the birational involutions

$$\psi(x, y) = \left(x, \frac{A_{-1}(x)}{A_{+1}(x)} \frac{1}{y} \right), \quad \phi(x, y) = \left(\frac{B_{-1}(y)}{B_{+1}(y)} \frac{1}{x}, y \right),$$

and thus under any element of the (dihedral) group

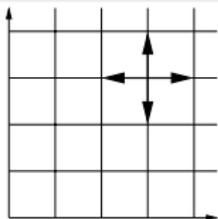
$$\mathcal{G}_{\mathcal{S}} := \langle \psi, \phi \rangle.$$

Examples of groups

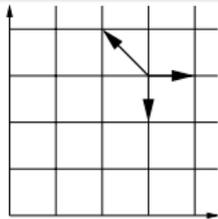


Order 4,

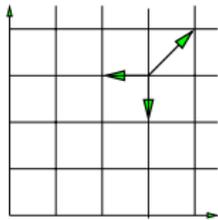
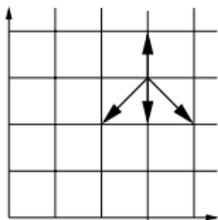
Examples of groups



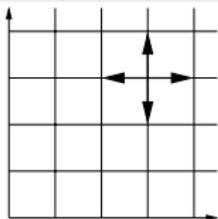
Order 4,



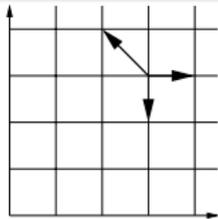
order 6,



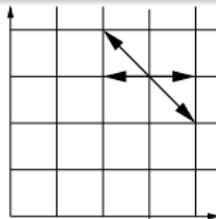
Examples of groups



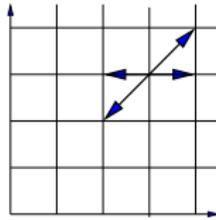
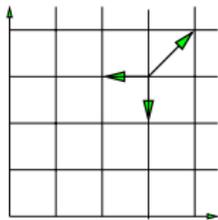
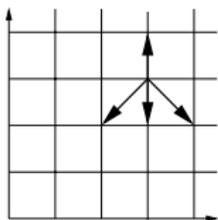
Order 4,



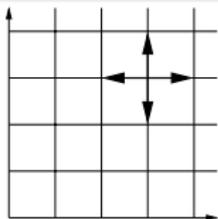
order 6,



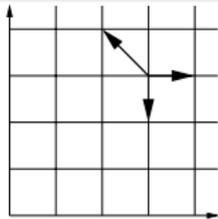
order 8,



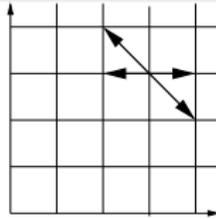
Examples of groups



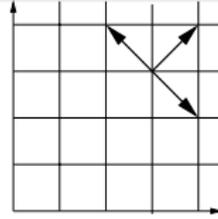
Order 4,



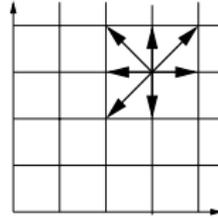
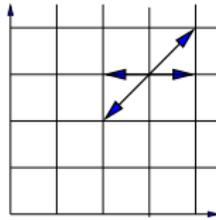
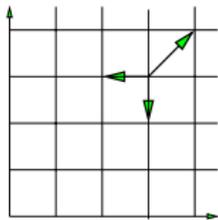
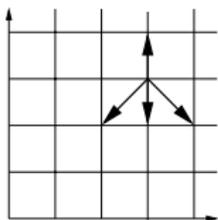
order 6,



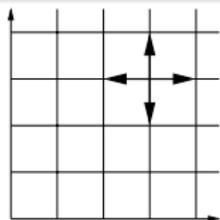
order 8,



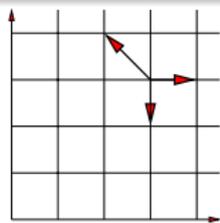
order ∞ .



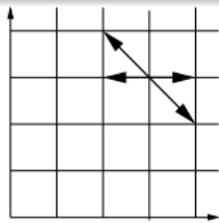
Examples of groups



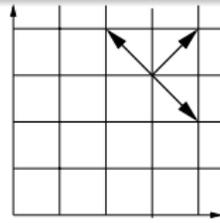
Order 4,



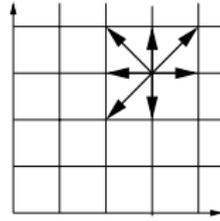
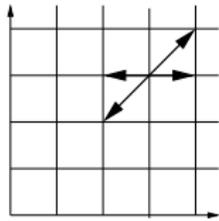
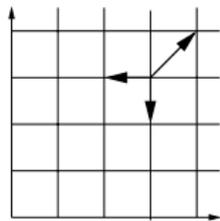
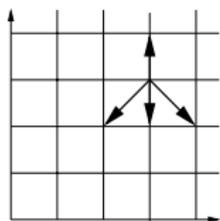
order 6,



order 8,



order ∞ .



$$\begin{array}{ccccc}
 & \Phi & \begin{pmatrix} y \\ x \end{pmatrix}, y & \Psi & \begin{pmatrix} y \\ x \end{pmatrix}, \frac{1}{x} & \Phi & & \\
 \begin{pmatrix} x, y \end{pmatrix} & \diagdown & & \text{---} & & \diagdown & & \begin{pmatrix} 1 \\ y \end{pmatrix}, \frac{1}{x} \\
 & \Psi & \begin{pmatrix} x, \frac{x}{y} \end{pmatrix} & \Phi & \begin{pmatrix} 1 \\ y \end{pmatrix}, \frac{x}{y} & \Psi & & \\
 & & & & & & &
 \end{array}$$