vers Mars	eille	R & S		S. C. M.	The floor
		Dorbsergelier	DA Lui 143	nigny	
Le Carelle Le	s Baumettes	DAC			
	Algeb	ora and A	Analys	is (APR CO
学行にすって		28	85m		MONTPHEET
	24	Philippe Du	mas	1-17-0	
1-1A	Cc 21	des Escourtines Inria Sacla	y	Col de Su 215m	giton
Col de Sormiou 181m	Col des Baun 200m Ph:	ilippe.dumas@in	nria.fr	18 50	Candelle
	19185271		Same Sh	Morgiou	And the second second

Journées Aléa 2016 7-11 mars 2016 Centre International de Rencontres Mathématiques

Cap Redon

Méditerranée

Calanque de Sormiou

Calanque de Morgiou

Calangue de La Triperie

Presqu'Île de la Triperie

Cap Wolgibu

All is available at http://specfun.inria.fr/dumas/Research/DAC/



Part I

Algebra

Overview of Part I

What is a DAC recurrence?

Algebraic machinery

Linear operators Basic functional properties

Definition of DAC recurrences

Comparison of types Generating functions Some links Mahler and sections All links Types equivalence Anatoli Karatsuba Frank Gray Rational sequence Linear representation Divide and conquer algorithms Combinatorics on words Number theory Moritz Stern, Achille Brocot



EURO-DA

European Design Automati



Meaning DAC Design Automation Conference DAC Digital-to-Analog Converter DAC Development Assistance Committee (OEC DAC Discretionary Access Control DAC District Advisory Council DAC Data Access Component DAC Downhill Assist Contr DAC Department of Arts a Mfrica) Dei Amoris DAC Divide and Conquer

Karatsuba's polynomial multiplication

$$a = a_0(x) + x^k a_1(x), \qquad b = b_0(x) + x^k b_1(x)$$

$$ab = a_0 \times b_0 + x^k ((a_0 + a_1) \times (b_0 + b_1) - a_0 \times b_0 - a_1 \times b_1) + x^{2k} a_1 \times b_1$$

Karatsuba's polynomial multiplication



floor and ceil type

What is a DAC recurrence? Gray code as usual binary code

n	bin(n)	$\operatorname{gray}(n)$	u_n
0	00	00	0
1	01	01	1
2	10	11	3
3	11	10	2
4	100	110	6
5	101	111	7
6	110	101	5
$\overline{7}$	111	100	4



Gray code as usual binary code

$$u_{4n} = 2u_{2n},$$

$$u_{4n+1} = -4u_n + 3u_{2n} + u_{2n+1},$$

$$u_{4n+2} = -4u_n + u_{2n} + 3u_{2n+1},$$

$$u_{4n+3} = 2u_{2n+1},$$
 with $u_0 = 0.$

by case type

From floor and ceil type to by case type: obvious!

$$u_n = 2u_{\lceil \frac{n}{2} \rceil} + u_{\lfloor \frac{n}{2} \rfloor} + 4(n-1), \qquad n \ge 2, \qquad \text{with } u_0 = 0, \, u_1 = 1,$$

$$u_{2n} = 3u_n + 8n - 4,$$
 with $u_0 = 0$
 $u_{2n+1} = 2u_{n+1} + u_n + 8n,$ with $u_1 = 1$

But from by case type to floor and ceil type?

$$u_{4n} = 2u_{2n},$$

$$u_{4n+1} = -4u_n + 3u_{2n} + u_{2n+1},$$

$$u_{4n+2} = -4u_n + u_{2n} + 3u_{2n+1},$$

$$u_{4n+3} = 2u_{2n+1},$$
 with $u_0 = 0.$

Algebraic machinery



$$u(x) = \sum_{n \ge 0} u_n x^n$$

 $u(x)$ formal series in $\mathbb{K}[[x]]$
$$(u_n)$$
 sequence in $\mathbb{K}^{\mathbb{N}}$

Both are exactly the same object.

$$u(x) = \sum_{n \ge 0} u_n x^n$$

radix $b \ge 2$ Mahler operator $Mu(x) = u(x^b)$

$$u(x) = \sum_{n \ge 0} u_n x^n$$

radix $b \ge 2$ Mahler operator $Mu(x) = u(x^b)$

$$0 \le r < b$$
 section operator $T_{b,r}u(x) = \sum_{k\ge 0} u_{bk+r}x^k$

$$u(x) = \sum_{n \ge 0} u_n x^n$$

radix $b \ge 2$ Mahler operator $Mu(x) = u(x^b)$

$$0 \le r < b$$
 section operator $T_{b,r}u(x) = \sum_{k\ge 0} u_{bk+r}x^k$

forward shift
$$Su(x) = \sum_{n \ge 0} u_{n+1} x^n$$

$$u(x) = \sum_{n \ge 0} u_n x^n$$

radix $b \ge 2$ Mahler operator $Mu(x) = u(x^b)$ $0 \le r < b$ section operator $T_{b,r}u(x) = \sum_{k\ge 0} u_{bk+r}x^k$

forward shift
$$Su(x) = \sum_{n \ge 0} u_{n+1} x^n$$

backward shift
$$xu(x) = \sum_{n \ge 0} u_{n-1}x^n$$

Algebraic machinery: Linear operators radix b = 2, by far the most usual case

Mahler operator

$$Mu(x) = u_0 \qquad \qquad + u_1 x^2 \qquad \qquad + u_2 x^4$$

. . . .

Algebraic machinery: Linear operators radix b = 2, by far the most usual case

Mahler operator

$$Mu(x) = u_0 \qquad \qquad + u_1 x^2 \qquad \qquad + u_2 x^4$$

section operator

 $T_{2,0}u(x) = u_0 + u_2x + u_4x^2 + u_6x^3 + \cdots$ even part

 $T_{2,1}u(x) = u_1 + u_3x + u_5x^2 + u_7x^3 + \cdots$ odd part

Algebraic machinery: Linear operators radix b = 2, by far the most usual case

Mahler operator

$$Mu(x) = u_0 \qquad \qquad + u_1 x^2 \qquad \qquad + u_2 x^4$$

section operator

 $T_{2,0}u(x) = u_0 + u_2x + u_4x^2 + u_6x^3 + \cdots$ even part

$$T_{2,1}u(x) = u_1 + u_3x + u_5x^2 + u_7x^3 + \cdots$$
 odd part

forward shift

$$Su(x) = u_1 + u_2x + u_3x^2 + u_4x^3 + \cdots$$

backward shift

$$xu(x) = u_0 x + u_1 x^2 + u_2 x^3 + u_3 x^4 + \cdots$$

$$T_{b,0}M = 1, \quad T_{b,r}M = 0 \qquad 1 \le r < b \qquad \text{obvious}$$

 $Mx = x^b M \qquad \text{obvious}$
 $ST_{b,r} = T_{b,r}S^b \qquad \text{the same, but...}$

$$ST_{b,r} = T_{b,r}S^b$$
, the same, but...

$$T_{b,r}u(x) = u_r + u_{b+r}x + u_{2b+r}x^2 + u_{3b+r}x^3 + \cdots$$

$$S^{b}u(x) = u_{b} + u_{b+1}x + u_{b+2}x^{2} + u_{b+3}x^{3} + \cdots$$

$$ST_{b,r} = T_{b,r}S^b$$
, the same, but...

$$T_{b,r}u(x) = u_r + u_{b+r}x + u_{2b+r}x^2 + u_{3b+r}x^3 + \cdots$$

$$ST_{b,r}u(x) = u_{b+r} + u_{2b+r}x + u_{3b+r}x^2 + u_{4b+r}x^3 + \cdots$$

$$T_{b,r}S^{b}u(x) = u_{b+r} + u_{2b+r}x^{2} + u_{3b+r}x^{3} + u_{4b+r}x^{3} + \cdots$$
$$S^{b}u(x) = u_{b} + u_{b+1}x + u_{b+2}x^{2} + u_{b+3}x^{3} + \cdots$$

$$ST_{b,r} = T_{b,r}S^b$$
, the same, but...

Proposition

The sections of a rational function are rational functions.

Proof

$$f \in \mathbb{K}(x), S^*f \in \mathcal{F}$$
 with dim $\mathcal{F} < \infty$,
 $g = T_{b,r}f, S^kg = T_{b,r}S^{bk}f \in T_{b,r}\mathcal{F}$ with dim $T_{b,r}\mathcal{F} < \infty$

motto : a subspace left stable by the operator(s)

$$T_{b,r}(fMg) = (T_{b,r}f)g$$
$$\sum_{0 \le r < b} x^r M T_{b,r} = 1$$

$$T_{b,r}(f(x)g(x^b)) = (T_{b,r}f(x))g(x)$$
$$\sum_{0 \le r < b} x^r T_{b,r}f(x^b) = f(x)$$

useful for products

It is possible to rebuild a function from its sections.

Example

$$T_{2,0}\frac{1+3x}{x^3(1+2x)} = \frac{1}{x(1-4x)}, \qquad T_{2,1}\frac{1+3x}{x^3(1+2x)} = \frac{1-6x}{x^2(1-4x)},$$
$$1 \times \frac{1}{x^2(1-4x^2)} + x \times \frac{1-6x^2}{x^4(1-4x^2)} = \frac{1+3x}{x^3(1+2x)} \qquad b = 2$$
$$1 \times T_{2,0}f(x^2) + xT_{2,1}f(x^2) = f(x)$$

 \wedge



Definition A (linear) Mahler equation is an equation

$$\ell_0(x)u(x) + \ell_1(x)u(x^b) + \dots + \ell_d(x)u(x^{b^d}) = v(x)$$

where $\ell_0(x)$, $\ell_1(x)$, ..., $\ell_d(x)$ and v(x) are polynomials in $\mathbb{K}[x]$.

$$L(x, M) = \ell_0(x) + \ell_1(x)M + \dots + \ell_d(x)M^d, \qquad L(x, M)u(x) = v(x)$$

motto : a subspace left stable by the operator(s)

Definition

A divide-and-conquer recurrence is the translation in terms of sequence of a Mahler equation.

Definition $u_{\nu} = 0$ if $\nu \notin \mathbb{N}_{\geq 0}$

Example

$$\begin{aligned} &(x+2x^2)u(x)-(1+x)u(x^2)+u(x^4)=0, & b=2\\ &u_{m-1}+2u_{m-2} & -u_{\frac{m}{2}}-u_{\frac{m-1}{2}} & +u_{\frac{m}{4}}=0, & m\geq 0\\ &u_9+2u_8 & -u_5-u_{\frac{9}{2}} & +u_{\frac{5}{2}}=0, & m=10\\ &u_{10}+2u_9 & -u_{\frac{11}{2}}-u_5 & +u_{\frac{11}{4}}=0, & m=11\\ &u_{11}+2u_{10} & -u_6-u_{\frac{11}{2}} & +u_3=0, & m=12\\ &u_{12}+2u_{11} & -u_{\frac{13}{2}}-u_6 & +u_{\frac{13}{4}}=0, & m=13 \end{aligned}$$

 \triangle

Example

$$(x+2x^2)u(x) - (1+x)u(x^2) + u(x^4) = 0, \qquad b = 2$$

$$u_{m-1} + 2u_{m-2} \quad -u_{\frac{m}{2}} - u_{\frac{m-1}{2}} \quad +u_{\frac{m}{4}} = 0, \qquad m \ge 0$$

$$u_9 + 2u_8 \qquad -u_5 \qquad = 0, \qquad m = 10$$

$$u_{10} + 2u_9 \qquad -u_5 \qquad = 0, \qquad m = 11$$

$$u_{11} + 2u_{10} \qquad -u_6 \qquad \qquad +u_3 = 0, \qquad m = 12$$

$$u_{12} + 2u_{11} - u_6 = 0, \qquad m = 13$$

 \wedge

fractional type

Comparison of types

Three types for the same thing, that's a lot!



 \blacktriangleright reference type = fractional type

$$t_m = u_{\frac{m-s}{b^k}}$$
 $t(x) = x^s u(x^{b^k})$

▶ floor and ceil type

$$t_m = u_{\lfloor \frac{n+s}{b} \rfloor}$$
 $t(x) = x^{-s}(1+x+\cdots+x^{b-1})u(x^b)$

$$-x^{-s}(1+x+\dots+x^{b-1})\sum_{n=0}^{q-1}u_nx^{bn}$$
$$-x^{-r}\sum_{i=0}^{r-1}x^iu_q$$

 $s = bq + r, \, |r| < b, \, \operatorname{sgn}(r) = \operatorname{sgn}(s)$

symmetrical Euclidean division

corrective term = 0 for $-\infty < s \le 0$

ceil ad libitum

$$\left\lceil \frac{n}{b} \right\rceil = \left\lfloor \frac{n+b-1}{b} \right\rfloor$$

► by case type

$$t_m = u_{bk+s}$$
 $t(x) = x^{-q} T_{b,r} u(x) - x^{-q} \sum_{j=0}^{r} u_{bj+r} x^j$
 $s = bq + r, \ 0 \le r < b$

natural Euclidean division

a-1

corrective term = 0 for $-\infty < s < b$

neglecting details:

$$t_m = u_{\frac{m-s}{b^k}} \qquad t(x) = \qquad x^s u(x^{b^k})$$
$$t_m = u_{\lfloor \frac{n+s}{b} \rfloor} \qquad t(x) = \qquad x^{-s}(1+x+\dots+x^{b-1})u(x^b)$$
$$t_m = u_{bk+s} \qquad t(x) = \qquad x^{-q}T_{b,r}u(x)$$

fractional type	<i>A</i> ahler	operator
floor and ceil type	Iahler	operator
by case type se	ection c	perators
floor and ceil type recurrence

Mahler equation _____ fractional type recurrence

system for sections

by case type recurrence



system for sections

by case type recurrence







Comparison of types: Mahler and sections

Theorem

If u is a formal series which is a solution of a non trivial Mahler equation, then, under the action of the section operators, it generates a finite dimensional $\mathbb{K}(x)$ -space. Conversely, if the iterated sections of a formal series u remain in a finite dimensional $\mathbb{K}(x)$ -space, then u is a solution a non trivial Mahler equation.

variation on

Gilles Christol, Teturo Kamae, Michel Mendès France, and Gérard Rauzy. Suites algébriques, automates et substitutions.

Bull. Soc. Math. France, 108(4):401-419, 1980.

motto : a subspace left stable by the operator(s)

Comparison of types: All links



strongly connected graph

Comparison of types: Types equivalence

Theorem

For a sequence (u_n) with support in $\mathbb{N}_{\geq 0}$ and for its generating function u(x), with a given integer $b \geq 2$,

- ▶ a fractional type recurrence,
- ▶ a floor and ceil type recurrence,
- ▶ a by case type equation,
- ▶ a Mahler equation,
- ▶ a system about the sections,

all have the same expressiveness.



$$u_n = 2u_{\lceil \frac{n}{2} \rceil} + u_{\lfloor \frac{n}{2} \rfloor} + 4(n-1) \qquad n \ge 2,$$

with $u(0) = 0, u(1) = 1$

$$u_n = 2u_{\lceil \frac{n}{2} \rceil} + u_{\lfloor \frac{n}{2} \rfloor} + 4(n-1) \qquad n \ge 2,$$

with $u(0) = 0, u(1) = 1$

$$xu(x) - (1+x)(2+x)u(x^2) = -x^2 + 4\frac{x}{(1-x)^2}$$

$$u_n = 2u_{\lceil \frac{n}{2} \rceil} + u_{\lfloor \frac{n}{2} \rfloor} + 4(n-1) \qquad n \ge 2,$$

with $u(0) = 0, u(1) = 1$

$$xu(x) - (1+x)(2+x)u(x^2) = -x^2 + 4\frac{x^3}{(1-x)^2}$$

$$u_{m-1} - \left(2u_{\frac{m}{2}} + 3u_{\frac{m-1}{2}} + u_{\frac{m-2}{2}}\right) = 4(m-1)$$

$$u_n = 2u_{\lceil \frac{n}{2} \rceil} + u_{\lfloor \frac{n}{2} \rfloor} + 4(n-1) \qquad n \ge 2,$$

with $u(0) = 0, u(1) = 1$
$$xu(x) - (1+x)(2+x)u(x^2) = -x^2 + 4\frac{x^3}{(1-x)^2}$$

$$u_{m-1} - (2u_{\frac{m}{2}} + 3u_{\frac{m-1}{2}} + u_{\frac{m-2}{2}}) = 4(m-1)$$

$$u_{2k-1} = 2u_k + u_{k-1} + 8k - 4, \qquad k \ge 2,$$

$$u_{2k} = 3u_k + 8k, \qquad k \ge 1$$

$$u(x) = \frac{(1+x)(2+x)}{x}u(x^2) - x + 4\frac{x^2}{(1-x)^2}$$
$$T_{2,0}u(x) = 3u(x) + \frac{4x+4x^2}{(1-x)^2} \qquad T_{2,1}u(x) = \frac{2+x}{x}u(x) - \frac{1-10x+x^2}{(1-x)^2}$$



March 17, 1953

F. GRAY

2,632,058

PULSE CODE COMMUNICATION

Filed Nov. 13, 1947

4 Sheets-Sheet 1



51/151

$$u_{4n} = 2u_{2n}, u_0 = 0$$

$$u_{4n+1} = -4u_n + 3u_{2n} + u_{2n+1}$$

$$u_{4n+2} = -4u_n + u_{2n} + 3u_{2n+1}$$

$$u_{4n+3} = 2u_{2n+1}$$

$$u_{4n} = 2u_{2n}, u_0 = 0$$

$$u_{4n+1} = -4u_n + 3u_{2n} + u_{2n+1}$$

$$u_{4n+2} = -4u_n + u_{2n} + 3u_{2n+1}$$

$$u_{4n+3} = 2u_{2n+1}$$

$$T_{2,0}u(x) = 2T_{2,0}u(x)$$

$$T_{4,1}u(x) = -4u(x) + 3T_{2,0}u(x) + T_{2,1}u(x)$$

$$T_{4,2}u(x) = -4u(x) + T_{2,0}u(x) + 3T_{2,1}u(x)$$

$$T_{4,3}u(x) = 2T_{2,1}u(x)$$

$$u_{4n} = 2u_{2n}, u_0 = 0$$

$$u_{4n+1} = -4u_n + 3u_{2n} + u_{2n+1}$$

$$u_{4n+2} = -4u_n + u_{2n} + 3u_{2n+1}$$

$$u_{4n+3} = 2u_{2n+1}$$

$$\begin{split} T_{2,0}u(x) &= 2T_{2,0}u(x) \\ T_{4,1}u(x) &= -4u(x) + 3T_{2,0}u(x) + T_{2,1}u(x) \\ T_{4,2}u(x) &= -4u(x) + T_{2,0}u(x) + 3T_{2,1}u(x) \\ T_{4,3}u(x) &= 2T_{2,1}u(x) \end{split}$$

$$v_1(x) = u(x), v_2(x) = T_{2,0}u(x), v_3(x) = T_{2,1}u(x)$$

$$\begin{split} T_{2,0}u(x) &= 2T_{2,0}u(x) \\ T_{4,1}u(x) &= -4u(x) + 3T_{2,0}u(x) + T_{2,1}u(x) \\ T_{4,2}u(x) &= -4u(x) + T_{2,0}u(x) + 3T_{2,1}u(x) \\ T_{4,3}u(x) &= 2T_{2,1}u(x) \end{split}$$

$$v_1(x) = u(x), v_2(x) = T_{2,0}u(x), v_3(x) = T_{2,1}u(x)$$

$$T_{2,0}v_1(x) = v_2(x)$$

$$T_{2,0}v_2(x) = 2v_2(x)$$

$$T_{2,0}v_3(x) = -4v_1(x) + 3v_2(x) + v_3(x)$$

$$T_{2,1}v_1(x) = v_3(x)$$

$$T_{2,1}v_2(x) = -4v_1(x) + v_2(x) + 3v_3(x)$$

$$T_{2,1}v_3(x) = 2v_3(x)$$

$$v_1(x) = u(x), v_2(x) = T_{2,0}u(x), v_3(x) = T_{2,1}u(x)$$

$$T_{2,0}v_1(x) = v_2(x)$$

$$T_{2,0}v_2(x) = 2v_2(x)$$

$$T_{2,0}v_3(x) = -4v_1(x) + 3v_2(x) + v_3(x)$$

$$T_{2,1}v_1(x) = v_3(x)$$

$$T_{2,1}v_2(x) = -4v_1(x) + v_2(x) + 3v_3(x)$$

$$T_{2,1}v_3(x) = 2v_3(x)$$

$$A_0 = \begin{bmatrix} 0 & 0 & -4 \\ 1 & 2 & 3 \\ 0 & 0 & 1 \end{bmatrix}, A_1 = \begin{bmatrix} 0 & -4 & 0 \\ 0 & 1 & 0 \\ 1 & 3 & 2 \end{bmatrix}, C = \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}$$

.

$$v_1(x) = u(x), v_2(x) = T_{2,0}u(x), v_3(x) = T_{2,1}u(x)$$

$$A_0 = \begin{bmatrix} 0 & 0 & -4 \\ 1 & 2 & 3 \\ 0 & 0 & 1 \end{bmatrix}, A_1 = \begin{bmatrix} 0 & -4 & 0 \\ 0 & 1 & 0 \\ 1 & 3 & 2 \end{bmatrix}, C = \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}$$

$$\begin{bmatrix} v_{1}(x) & v_{2}(x) & v_{3}(x) \end{bmatrix} \qquad v(x) = T_{2,0}v(x^{2}) + xT_{2,1}v(x^{2})$$
$$= \begin{bmatrix} v_{1}(x^{2}) & v_{2}(x^{2}) & v_{3}(x^{2}) \end{bmatrix} \begin{bmatrix} 0 & 0 & -4 \\ 1 & 2 & 3 \\ 0 & 0 & 1 \end{bmatrix}$$
$$+ x \begin{bmatrix} v_{1}(x^{2}) & v_{2}(x^{2}) & v_{3}(x^{2}) \end{bmatrix} \begin{bmatrix} 0 & -4 & 0 \\ 0 & 1 & 0 \\ 1 & 3 & 2 \end{bmatrix}$$

•

$$\begin{bmatrix} v_{1}(x) & v_{2}(x) & v_{3}(x) \end{bmatrix} \qquad v(x) = T_{2,0}v(x^{2}) + xT_{2,1}v(x^{2})$$
$$= \begin{bmatrix} v_{1}(x^{2}) & v_{2}(x^{2}) & v_{3}(x^{2}) \end{bmatrix} \begin{bmatrix} 0 & 0 & -4 \\ 1 & 2 & 3 \\ 0 & 0 & 1 \end{bmatrix}$$
$$+ x \begin{bmatrix} v_{1}(x^{2}) & v_{2}(x^{2}) & v_{3}(x^{2}) \end{bmatrix} \begin{bmatrix} 0 & -4 & 0 \\ 0 & 1 & 0 \\ 1 & 3 & 2 \end{bmatrix}$$

$$V(x) = \begin{bmatrix} v_1(x) & v_2(x) & v_3(x) \end{bmatrix}, \qquad A(x) = A_0 + xA_1,$$

 $V(x) = V(x^2)A(x), \qquad u(x) = V(x)C$

 $V(x) = V(x^2)A(x), \qquad u(x) = V(x)C$

$$u(x) = V(x)C$$

$$u(x) = V(x^{2})A(x)C$$

$$u(x) = V(x^{4})A(x^{2})A(x)C$$

$$u(x) = V(x^{8})A(x^{4})A(x^{2})A(x)C$$

$$u(x^{8}) = V(x^{8})C$$

$$u(x^{4}) = V(x^{8})A(x^{4})C$$

$$u(x^{2}) = V(x^{8})A(x^{4})A(x^{2})C$$

$$u(x) = V(x^{8})A(x^{4})A(x^{2})A(x)C$$

$$\begin{split} u(x^8) &= V(x^8)C\\ u(x^4) &= V(x^8)A(x^4)C\\ u(x^2) &= V(x^8)A(x^4)A(x^2)C\\ u(x) &= V(x^8)A(x^4)A(x^2)A(x)C \end{split}$$

 $4~{\rm column}~{\rm vectors}$ in dimension 3

$$A(x) = \begin{bmatrix} 0 & -4x & -4\\ 1 & 2+x & 3\\ x & 3x & 1+2x \end{bmatrix}$$
$$\Gamma(x) = \begin{bmatrix} 1 & 0 & -4x^2 - 4x^4 & -4x - 12x^2 - 8x^3 - 8x^4 - 12x^5 - 4x^6\\ 0 & 1 & 2+3x^2 + x^4 & 4+5x + 7x^2 + 6x^3 + 2x^4 + 3x^5 + x^6\\ 0 & x^4 & x^2 + 3x^4 + 2x^6 & x + 3x^2 + 2x^3 + 6x^4 + 7x^5 + 5x^6 + 4x^7 \end{bmatrix}$$
$$K(x) = \begin{bmatrix} 0\\ 2\frac{(1+x)(1+x^4)}{x}\\ -\frac{(1+x)(1+2x+2x^3+x^4)}{x(1+x^2)}\\ 1 \end{bmatrix}$$

$$K(x) = \begin{bmatrix} 0 \\ 2\frac{(1+x)(1+x^4)}{x} \\ -\frac{(1+x)(1+2x+2x^3+x^4)}{x(1+x^2)} \\ 1 \end{bmatrix}$$

 $x(1+x^2)u(x) - (1+x)(1+2x+2x^3+x^4)u(x^2) + 2(1+x)(1+x^2)(1+x^4)u(x^4) = 0$

to summarize: from

$$\begin{split} T_{2,0}u(x) &= 2T_{2,0}u(x) \\ T_{4,1}u(x) &= -4u(x) + 3T_{2,0}u(x) + T_{2,1}u(x) \\ T_{4,2}u(x) &= -4u(x) + T_{2,0}u(x) + 3T_{2,1}u(x) \\ T_{4,3}u(x) &= 2T_{2,1}u(x) \end{split}$$

t	1	
	t	,
~	-	-

$$\begin{aligned} x(1+x^2)u(x) \\ &-(1+x)(1+2x+2x^3+x^4)u(x^2) \\ &+2(1+x)(1+x^2)(1+x^4)u(x^4)=0 \end{aligned}$$

$$\begin{aligned} x(1+x^2)u(x) \\ &- (1+x)(1+2x+2x^3+x^4)u(x^2) \\ &+ 2(1+x)(1+x^2)(1+x^4)u(x^4) = 0 \end{aligned}$$

$$u_n + u_{n-2} - \left(u_{\lfloor \frac{n+1}{2} \rfloor} + 2u_{\lfloor \frac{n}{2} \rfloor} + 2u_{\lfloor \frac{n-2}{2} \rfloor} + u_{\lfloor \frac{n-3}{2} \rfloor} \right) + 2\left(u_{\lfloor \frac{n+1}{4} \rfloor} + u_{\lfloor \frac{n-3}{4} \rfloor} \right) = 0$$

to summarize: from

$$u_{4n} = 2u_{2n}, u_0 = 0$$

$$u_{4n+1} = -4u_n + 3u_{2n} + u_{2n+1}$$

$$u_{4n+2} = -4u_n + u_{2n} + 3u_{2n+1}$$

$$u_{4n+3} = 2u_{2n+1}$$

t	1	Ъ
U	L	,

$$u_{n}+u_{n-2} \qquad u_{0} = 0, \ u_{1} = 1, \ u_{2} = 3$$
$$-\left(u_{\lfloor \frac{n+1}{2} \rfloor} + 2u_{\lfloor \frac{n}{2} \rfloor} + 2u_{\lfloor \frac{n-2}{2} \rfloor} + u_{\lfloor \frac{n-3}{2} \rfloor}\right)$$
$$+ 2\left(u_{\lfloor \frac{n+1}{4} \rfloor} + u_{\lfloor \frac{n-3}{4} \rfloor}\right) = 0$$

$$u_{n} + u_{n-2} - \left(u \lfloor \frac{n+1}{2} \rfloor + 2u \lfloor \frac{n}{2} \rfloor + 2u \lfloor \frac{n-2}{2} \rfloor + u \lfloor \frac{n-3}{2} \rfloor\right) + 2\left(u \lfloor \frac{n+1}{4} \rfloor + u \lfloor \frac{n-3}{4} \rfloor\right) = 0$$

$$u_{n-3} + u_{n-1} - u_{\frac{n-5}{2}} - 3u \frac{n-4}{2} - 2u \frac{n-3}{2} - 2u \frac{n-2}{2} - 3u \frac{n-4}{2} - 2u \frac{n-3}{2} - 2u \frac{n-2}{2} - 3u \frac{n-4}{2} - 2u \frac{n-4}{2} + 2u \frac{n-4}{2} + 2u \frac{n-4}{4} + 2u \frac{n-3}{4} + 2u \frac{n-3}{4} + 2u \frac{n-4}{4} + 2u \frac{n-3}{4} + 2u \frac{n-4}{4} + 2u \frac{n-4}$$

.

$$\begin{array}{ll} T_{2,0}u(x)=2T_{2,0}u(x) & u_{4n}=2u_{2n} \\ T_{4,1}u(x)=-4u(x)+3T_{2,0}u(x)+T_{2,1}u(x) & u_{4n+1}=-4u_n+3u_{2n}+u_{2n+1} \\ T_{4,2}u(x)=-4u(x)+T_{2,0}u(x)+3T_{2,1}u(x) & u_{4n+2}=-4u_n+u_{2n}+3u_{2n+1} \\ T_{4,3}u(x)=2T_{2,1}u(x) & u_{4n+3}=2u_{2n+1} \end{array}$$





Definition

A formal series (or a sequence) is rational wrt a numeration system with radix b, or is b-rational, if under the action of the section operators it generates a finite dimensionial \mathbb{K} -vector space.

Jean-Paul Allouche and Jeffrey Shallit. The ring of *k*-regular sequences. *Theoret. Comput. Sci.*, 98(2):163–197, 1992.

Proposition

A b-rational series satisfies a non trivial Mahler equation for the radix b.

Proposition

A formal series u(x) which satisfies a Mahler equation $(\omega \in \mathbb{N}_{\geq 0})$

$$x^{\omega}u(x) = c_1(x)u(x^b) + \dots + c_d(x)u(x^{b^d}),$$

with polynomial coefficients, is b-rational.

Proposition

A formal series u(x) which satisfies a Mahler equation $(\omega \in \mathbb{N}_{\geq 0})$

$$x^{\omega}u(x) = c_1(x)u(x^b) + \dots + c_d(x)u(x^{b^d}),$$

with polynomial coefficients, is b-rational.

Proposition

A sequence (u_n) which satisfies a fractional type recurrence

$$u_n = \sum_{k=1}^d \sum_{\ell=-s}^s c_{k,\ell} u_{\frac{n-\ell}{b^k}}$$

is b-rational.

Proposition

A formal series u(x) which satisfies a Mahler equation $(\omega \in \mathbb{N}_{\geq 0})$

$$x^{\omega}u(x) = c_1(x)u(x^b) + \dots + c_d(x)u(x^{b^d}),$$

with polynomial coefficients, is b-rational.

Proposition

A sequence (u_n) which satisfies a fractional type recurrence

$$u_n = \sum_{k=1}^d \sum_{\ell=-s}^s c_{k,\ell} u_{\frac{n-\ell}{b^k}}$$
 the true DAC recurrences!

is b-rational.
Theorem

The Nth coefficient of a b-rational series u(x) is expressed as

$$u_N = T_{b,r_\ell} \cdots T_{b,r_0} u(0)$$

if $N = (r_\ell \dots r_0)_b$.

Theorem

The Nth coefficient of a b-rational series u(x) is expressed as

$$u_N = T_{b,r_\ell} \cdots T_{b,r_0} u(0)$$

if $N = (r_\ell \dots r_0)_b$.

$$10 = (1010)_2 \qquad u(x) = u_0 + u_1 x + u_2 x^2 + \cdots$$
$$T_{2,0} u(x) = u_0 + u_2 x + u_4 x^2 + \cdots$$
$$T_{2,1} T_{2,0} u(x) = u_2 + u_6 x + u_{10} x^2 + \cdots$$
$$T_{2,0} T_{2,1} T_{2,0} u(x) = u_2 + u_{10} x + u_{18} x^2 + \cdots$$
$$T_{2,1} T_{2,0} T_{2,1} T_{2,0} u(x) = u_{10} + u_{26} x + u_{42} x^2 + \cdots$$
$$T_{2,1} T_{2,0} T_{2,1} T_{2,0} u(0) = u_{10}$$

Definition

A linear representation of a *b*-rational series u(x) or sequence (u_n) is a triplet (L, A, C) made from

- a row vector L (initial values);
- a family of square matrices $(A_r)_{0 \le r < b}$ (action);
- a column vector C (coordinates),

with the same size and coefficients in \mathbb{K} , such that

$$u_N = LA_{r_\ell} \cdots A_{r_0}C$$

when

$$N = (r_\ell \dots r_0)_b.$$

$$u_N = LA_{r_\ell} \cdots A_{r_0}C$$

when

$$N = (r_\ell \dots r_0)_b$$

for the Gray code:

$$L = \begin{bmatrix} 0 & 0 & 1 \end{bmatrix},$$

$$A_0 = \begin{bmatrix} 0 & 0 & -4 \\ 1 & 2 & 3 \\ 0 & 0 & 1 \end{bmatrix}, A_1 = \begin{bmatrix} 0 & -4 & 0 \\ 0 & 1 & 0 \\ 1 & 3 & 2 \end{bmatrix}, C = \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}$$

•

Bestiary



Divide and conquer algorithms



Divide and conquer algorithms



binary search extrema kth smallest element mergesort quicksort



binary powering Karatsuba algorithm Toom-Cook algorithm Schönhage-Strassen algorithm multipoint evaluation

sorting

algebraic computations



convex hull nearest pair Voronoï diagram maxima in dim ≥ 2



Strassen multiplication triangular matrix inversion fast Fourier transform singular values decomposition eigenvalues/vectors of symmetric tridiagonal matrices

matrix computations

79/151

algorithmic geometry



NETWORKS AND TELECOMMUNICATIONS SERIES



Formal Languages, Automata and Numeration Systems 1

> Introduction to Combinatorics on Words





WILEY

The goal is that, after reading this book (or at least parts of this book), the reader should be able to fruitfully attend a conference or a seminar in the field.

Michel Rigo

NETWORKS AND TELECOMMUNICATIONS SERIE



Formal Languages, Automata and Numeration Systems

> Introduction to Combinatorics on Words

> > Michel Rigo

SIE

WILE

1794

The goal is that, after reading this book (or at least parts of this book), the reader should be able to fruitfully attend a conference or a seminar in the field.

N. Pytheas Fogg

Substitutions in **Dynamics, Arithmetics** and Combinatorics

Editors: V. Rerthé S.Ferenczi C.Mauduit A. Siegel

Michel Rigo



Axel Thue

Marston Morse

 $e_1(n)$ number of 1's in binary expansion of n $e_{11}(n)$ Golay-Rudin-Shapiro sequence $(-1)^{e_1(n)}$ Thue-Morse sequence $(-1)^{e_1(3n)}$ Newman-Slater-Coquet overlapping free words

patterns counting

$$w_0 = \varepsilon$$
$$w_{k+1} = w_k 1 \overline{w}_k^R$$



substitutions

$$w_0 = \varepsilon$$
$$w_{k+1} = w_k 1 \overline{w}_k^R$$

$$\begin{array}{l} a \longrightarrow ab \\ b \longrightarrow cb \\ c \longrightarrow ad \\ d \longrightarrow cd \\ a \rightarrow ab \rightarrow abcb \rightarrow abcbadcb \rightarrow \dots \\ a := 1, \quad b := 1, \quad c := 0, \quad d := 0 \\ w_{\infty} = 0010011000110110\dots \end{array}$$



$$\begin{aligned} & u & \longrightarrow u_0 \\ w_0 &= \varepsilon \\ & w_{k+1} &= w_k 1 \overline{w}_k^R \\ & u_{4n} &= 0 \\ & u_{4n+2} &= 1 \\ & u_{2n+1} &= u_n \end{aligned} \qquad \begin{aligned} & u & \longrightarrow u_0 \\ & b & \longrightarrow cb \\ & c & \longrightarrow ad \\ & d & \longrightarrow cd \\ & a & \rightarrow ab & \rightarrow abcb & \rightarrow abcbadcb & \rightarrow \dots \\ & a &:= 1, \quad b &:= 1, \quad c &:= 0, \quad d &:= 0 \\ & w_\infty &= 0010011000110110\dots \end{aligned}$$

 α

\ ah



founded in 1964 by N. J. A. Sloane







Number theory



elementary number theory



Moritz Stern, Achille Brocot



Moritz Stern, Achille Brocot



$$u_0 = 0, u_1 = 1,$$
 $u_{2n} = u_n,$ $u_{2n+1} = u_n + u_{n+1}$
 $n \in \mathbb{N}_{>0} \longmapsto r_n = \frac{u_{n+2}}{u_{n+1}} \in \mathbb{Q}_{>0}$ one-t-one

Part II

Analysis

Overview of Part II

Slaves bound Goal Integers and words Extraction of classical rational sequences A mere idea Joint spectral radius Dilation equations Theorem A worked example Linear representation Joint spectral radius Jordan reduction Dilation equation Cascade algorithm

!

What I did not speak about



source

Jon Louis Bentley, Dorothea Haken, and James B. Saxe. A general method for solving divide-and-conquer recurrences. SIGACT News, 12(3):36–44, September 1980.

a good version:

 Alin Bostan, Frédéric Chyzak, Marc Giusti, Romain Lebreton, Grégoire Lecerf, Bruno Salvy, and Éric Schost.
 Algorithmes Efficaces en Calcul Formel.

Version provisoire disponible à l'url http://specfun.inria.fr/chyzak/mpri/poly.pdf, 2016.

Theorem

Let
$$(c_n)$$
 be s.t. $0 \le c_n \le \begin{cases} ac_{\lceil \frac{n}{b} \rceil} + t_n, & \text{if } n \ge n_0 \ge b, \\ \kappa & \text{otherwise,} \end{cases}$ with

- $b \ge 2$ is an integer;
- a > 0 is a real number;
- $\kappa \geq is \ a \ real \ number;$
- ► t a toll function
 - ▶ non decreasing,

• such that $a't_n \leq t_{bn} \leq a''t_n$ for some constants $a'' \geq a' > 1$,

then

$$c_n \mathop{=}\limits_{n \to \infty} \left\{ \begin{array}{ll} O(t_n) & a' > a, \\ O(t_n \log n) & if \ a' = a, \\ O(n^{\alpha - \alpha'} t_n) & if \ a' < a \end{array} \right.$$

with $\alpha = \log_b a$, $\alpha' = \log_b a'$







We want to catch the oscillations!

Goal

$$\begin{split} u(x) &= \sum_{n \ge 0} u_n x^n = \prod_{k \ge 0} \frac{1}{1 - \rho x^{2^k}} \\ \rho > 1 & u_n = u(1/\rho^2) \rho^n + O(\rho^{n/2}) \\ \rho &= 1 & \log u_{2n} = \log u_{2n+1} = \frac{1}{2\log 2} \log^2 \frac{n}{\log n} \\ &+ \left(\frac{1}{2} + \frac{1}{\log 2} + \frac{\log \log 2}{\log 2}\right) \log n \\ &+ O(\log \log n) \\ \rho < 1 & \sum_{n=1}^N u_n = \varphi(\log_2 n) N^\alpha + O(N^{\alpha - 1/2 + \varepsilon}) \\ &\alpha = \log_2 \frac{1}{1 - \rho} \end{split}$$

Goal

$$\begin{split} u(x) &= \sum_{n \ge 0} u_n x^n = \prod_{k \ge 0} \frac{1}{1 - \rho x^{2^k}} \\ \rho > 1 & u_n = u(1/\rho^2) \rho^n + O(\rho^{n/2}) \\ \rho &= 1 & \log u_{2n} = \log u_{2n+1} = \frac{1}{2\log 2} \log^2 \frac{n}{\log n} \\ &+ \left(\frac{1}{2} + \frac{1}{\log 2} + \frac{\log \log 2}{\log 2}\right) \log n \\ &+ O(\log \log n) \\ \rho < 1 & \sum_{n=1}^N u_n = \varphi(\log_2 n) N^\alpha + O(N^{\alpha - 1/2 + \varepsilon}) \\ &\alpha = \log_2 \frac{1}{1 - \rho} \end{split}$$

 $u_n = \rho u_{n-1} + u_{\frac{n}{2}}$

We want to study the asymptotic behavior of true divide and conquer sequences, that is b-rational sequences.

Some tools



Integers and words $b \ge 2, \mathcal{Z} = \{0, 1, \dots, b-1\}$

> generating series

formal series

$$u(x) = \sum_{n \ge 0} u_n x^n$$
$$T_{b,r} u(x) = \sum_{k \ge 0} u_{bk+r} x^k$$

$$s = \sum_{w \in \mathcal{Z}^*} s_w w$$
$$sr^{-1} = \sum_{w = w'r} s_w w'$$

Integers and words $b > 2, \mathcal{Z} = \{0, 1, \dots, b-1\}$ formal generating series series
$$\begin{split} u(x) &= \sum_{n \geq 0} u_n x^n \qquad \qquad s = \sum_{w \in \mathcal{Z}^*} s_w \, w \\ T_{b,r} u(x) &= \sum_{k \geq 0} u_{bk+r} x^k \qquad sr^{-1} = \sum_{w = w'r} s_w \, w' \end{split}$$
 $w = \overline{w'r}$ $n = (w)_b \in \mathbb{N} \longrightarrow w \in \mathcal{Z}^* \longrightarrow s_w = u_n$

maps composition

Integers and words

We do not use the words which begins with some zeroes.

Integers and words

We do not use the words which begins with some zeroes.

Definition

A linear representation (L, A, C) is insensitive to the leftmost zeroes, or zero-insensitive, if it satisfies $LA_0 = L$.

Integers and words

We do not use the words which begins with some zeroes.

Definition

A linear representation (L, A, C) is insensitive to the leftmost zeroes, or zero-insensitive, if it satisfies $LA_0 = L$.

Concretely, we always use zero-insensitive linear representations.

Extraction of classical rational sequences

sequence integers whose b-ary expansions have a regular expression e.g. $2^k = (10^k)_2, 2^k - 1 = (1^k)_2$

Stern-Brocot sequence

$$L = \begin{bmatrix} 0 & 1 \end{bmatrix}, \quad A_0 = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}, \quad A_1 = \begin{bmatrix} 0 & -1 \\ 1 & 2 \end{bmatrix}, \quad C = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$
$$u_{2^{k}-1} = LA_1^k C = \begin{bmatrix} 1 & 2 \end{bmatrix} \begin{bmatrix} 1-k & -k \\ k & 1+k \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = k+1$$
$$\begin{pmatrix} 1 \\ 2 \\ 1 \\ 1 \\ 3 \\ 1 \\ 5 \\ 4 \\ 7 \\ 1 \\ 6 \\ 5 \\ 9 \\ 4 \\ 11 \\ 7 \\ 10 \\ 3 \\ 11 \\ 8 \\ 13 \\ 5 \\ 12 \\ 7 \\ 9 \\ 2 \\ 9 \\ 2 \\ 9 \\ \dots$$
Extraction of classical rational sequences

Stern-Brocot sequence

$$L = \begin{bmatrix} 0 & 1 \end{bmatrix}, \quad A_0 = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}, \quad A_1 = \begin{bmatrix} 0 & -1 \\ 1 & 2 \end{bmatrix}, \quad C = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

$$\sum_{n=2^{k}}^{2^{k+1}-1} u_n = LA_1(A_0 + A_1)^k C = \begin{bmatrix} 1 & 2 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ (3^k - 1)/2 & 3^k \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = 3^k$$

u(x) b-rational series

u(x) *b*-rational series

 $\delta(x) = (1-x)u(x)$ b-rational

u(x) b-rational series $\delta(x) = (1-x)u(x)$ b-rational

(L, A, C)linear representation for $\delta(x)$, insensitive to the leftmost zeroes

u(x) b-rational series $\delta(x) = (1-x)u(x)$ b-rational

(L, A, C)

linear representation for $\delta(x)$, insensitive to the leftmost zeroes

$$u_N = \sum_{n=0}^N \delta_n = \sum_{n \le N} LA_w C, \qquad (w)_b = n$$

u(x) b-rational series $\delta(x) = (1-x)u(x)$ b-rational

(L, A, C)

linear representation for $\delta(x)$, insensitive to the leftmost zeroes

$$u_N = \sum_{n=0}^N \delta_n = \sum_{\substack{n \le N}} LA_w C, \qquad (w)_b = n$$
$$S_K(x) = \sum_{\substack{|w|=K\\(0.w)_b \le x}} A_w C \qquad 0 \le x \le 1$$

$$u_N = \sum_{n=0}^N \delta_n = \sum_{\substack{n \le N}} LA_w C, \qquad (w)_b = n$$
$$S_K(x) = \sum_{\substack{|w|=K\\(0.w)_b \le x}} A_w C \qquad 0 \le x \le 1$$

$$\begin{split} \delta_0 &= LA_0C\\ \delta_1 &= LA_1C\\ \delta_2 &= LA_1A_0C\\ \delta_3 &= LA_1A_1C\\ \delta_4 &= LA_1A_0A_0C\\ \delta_5 &= LA_1A_0A_1C \end{split}$$

$$u_N = \sum_{n=0}^N \delta_n = \sum_{\substack{n \le N}} LA_w C, \qquad (w)_b = n$$
$$S_K(x) = \sum_{\substack{|w| = K\\(0.w)_b \le x}} A_w C \qquad 0 \le x \le 1$$

$$\begin{split} \delta_{0} &= LA_{0}C & \delta_{0} &= LA_{0}A_{0}A_{0}C \\ \delta_{1} &= LA_{1}C & \delta_{1} &= LA_{0}A_{0}A_{1}C \\ \delta_{2} &= LA_{1}A_{0}C & \delta_{2} &= LA_{0}A_{1}A_{0}C \\ \delta_{3} &= LA_{1}A_{1}C & \delta_{3} &= LA_{0}A_{1}A_{1}C \\ \delta_{4} &= LA_{1}A_{0}A_{0}C & \delta_{4} &= LA_{1}A_{0}A_{0}C \\ \delta_{5} &= LA_{1}A_{0}A_{1}C & \delta_{5} &= LA_{1}A_{0}A_{1}C \end{split}$$

$$u_N = \sum_{n=0}^N \delta_n = \sum_{n \le N} LA_w C, \qquad (w)_b = n$$
$$S_K(x) = \sum_{\substack{|w| = K\\(0,w)_b \le x}} A_w C \qquad 0 \le x \le 1$$

$$\delta_0 = LA_0A_0A_0C$$

$$\delta_1 = LA_0A_0A_1C$$

$$\delta_2 = LA_0A_1A_0C$$

$$\delta_3 = LA_0A_1A_1C$$

$$\delta_4 = LA_1A_0A_0C$$

$$\delta_5 = LA_1A_0A_1C$$

$$u_N = \sum_{n=0}^N \delta_n = \sum_{n \le N} LA_w C, \qquad (w)_b = n$$
$$S_K(x) = \sum_{\substack{|w| = K\\(0,w)_b \le x}} A_w C \qquad 0 \le x \le 1$$

 $u_{5} = LA_{0}A_{0}A_{0}C$ $+ LA_{0}A_{0}A_{1}C$ $+ LA_{0}A_{1}A_{0}C$ $+ LA_{0}A_{1}A_{1}C$ $+ LA_{1}A_{0}A_{0}C$ $+ LA_{1}A_{0}A_{1}C$

$$u_N = \sum_{n=0}^N \delta_n = \sum_{n \le N} LA_w C, \qquad (w)_b = n$$
$$S_K(x) = \sum_{\substack{|w| = K\\(0.w)_b \le x}} A_w C \qquad 0 \le x \le 1$$
$$UA_0 A_0 A_0 C \qquad US_2(5/8) = UA_0 A$$

 $u_{5} = LA_{0}A_{0}A_{0}C \qquad LS_{3}(5/8) = \\ + LA_{0}A_{0}A_{1}C \qquad + \\ + LA_{0}A_{1}A_{0}C \qquad + \\ + LA_{1}A_{0}A_{0}C \qquad + \\ + LA_{1}A_{0}A_{1}C \qquad + \\ \end{pmatrix}$

 $LS_{3}(5/8) = LA_{0}A_{0}A_{0}C$ $+ LA_{0}A_{0}A_{1}C$ $+ LA_{0}A_{1}A_{0}C$ $+ LA_{0}A_{1}A_{1}C$ $+ LA_{1}A_{0}A_{0}C$ $+ LA_{1}A_{0}A_{1}C$

$$u_N = \sum_{n=0}^N \delta_n = \sum_{\substack{n \le N}} LA_w C, \qquad (w)_b = n$$
$$S_K(x) = \sum_{\substack{|w| = K\\(0.w)_b \le x}} A_w C \qquad 0 \le x \le 1$$

Proposition

Let (L, A, C) be a insensitive to the leftmost zeroes linear representation for the sequence (δ_n) of backward differences of a *b*-rational sequence (u_n) . Then

$$u_N = LS_{K+1}(b^{\{\log_b N\}-1}),$$

with $K = \lfloor \log_b N \rfloor$ and $\{t\} = t - \lfloor t \rfloor$.

$$u_N = \sum_{n=0}^N \delta_n = \sum_{n \le N} LA_w C, \qquad (w)_b = n$$
$$S_K(x) = \sum_{\substack{|w|=K\\(0.w)_b \le x}} A_w C \qquad 0 \le x \le 1$$

Proposition

The sequence $S_K(x)$ satisfies

$$S_{K+1}(x) = \sum_{r_1 < x_1} A_{r_1} Q^K C + A_{x_1} S_K(bx - x_1),$$

for $x = (0.x_1x_2...)_b$ in [0, 1[.

$$u_N = LA_{r_\ell} \cdots A_{r_0}C \qquad \text{for } N = (r_\ell \dots r_0)_b$$

$|u_N| \le ||L|| ||A_{r_\ell}|| \cdots ||A_{r_0}|| ||C||$

$$u_N = LA_{r_\ell} \cdots A_{r_0}C \qquad \text{for } N = (r_\ell \dots r_0)_b$$

$$|u_N| \le ||L|| ||A_{r_{\ell}}|| \cdots ||A_{r_0}|| ||C||$$

$$\le ||L|| ||C||a^{\ell+1} = ||L|| ||C||a^{\lfloor \log_b N \rfloor} \le K N^{\log_b a}$$

$$u_N = LA_{r_\ell} \cdots A_{r_0}C$$
 for $N = (r_\ell \dots r_0)_b$

$$\begin{aligned} |u_N| &\leq \|L\| \|A_{r_\ell}\| \cdots \|A_{r_0}\| \|C\| \\ &\leq \|L\| \|C\| a^{\ell+1} = \|L\| \|C\| a^{\lfloor \log_b N \rfloor} \leq K N^{\log_b a} \end{aligned}$$

Proposition

A b-rational sequence has a growth order at most polynomial.

Proposition

Let $A = (A_z)_{z \in \mathbb{Z}}$ be a finite family of square matrices. The sequence

$$\hat{\rho}_{\ell}(A) = \max_{w \in \mathcal{Z}^{\ell}} \|A_w\|^{1/\ell},$$

converges towards

$$\hat{\rho}(A) = \lim_{\ell \to +\infty} \hat{\rho}_{\ell}(A) = \inf_{\ell} \hat{\rho}_{\ell}(A).$$

Moreover the limit is independent of the used multiplicative norm. It is the joint spectral radius of A.

Proposition

If (L, A, C) is a linear representation for a b-rational sequence (u_n) , then for all $\varepsilon > 0$

$$u_N \underset{N \to +\infty}{=} O(N^{\log_b \hat{\rho}(A) + \varepsilon})$$

Karatsuba

$$A_{0} = \begin{bmatrix} 3 & 1 & 0 & 0 & 0 \\ 0 & 2 & 0 & 0 & 0 \\ 0 & -1 & 1 & 0 & 0 \\ 4 & 10 & 1 & 2 & 1 \\ 4 & -1 & 0 & 0 & 1 \end{bmatrix}, \qquad A_{1} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 2 & 3 & 0 & 0 & 0 \\ -1 & 4 & 2 & 1 & 0 \\ 10 & 4 & 0 & 1 & 2 \\ -1 & 0 & 0 & 0 & 0 \end{bmatrix}$$

$$\|M\|_{1} = \max_{j} \sum_{i} |M_{i,j}|, \quad \|M\|_{\infty} = \max_{i} \sum_{j} |M_{i,j}|,$$
$$\|M\|_{F} = \left(\sum_{i,j} |M_{i,j}|^{2}\right)^{1/2}$$



128 / 151

Proposition

If the matrices of $A = (A_z)_{z \in \mathbb{Z}}$ can be simultaneously block-triangulated,

$$P^{-1}A_z P = \begin{pmatrix} B_z & C_z \\ 0 & D_z \end{pmatrix}, \qquad z \in \mathcal{Z},$$

then the joint spectral radius of A is

 $\hat{\rho}(A) = \max(\hat{\rho}(B), \hat{\rho}(D)).$

$$A_{0} = \begin{bmatrix} 3 & 1 & 0 & 0 & 0 \\ 0 & 2 & 0 & 0 & 0 \\ 0 & -1 & 1 & 0 & 0 \\ 4 & 10 & 1 & 2 & 1 \\ 4 & -1 & 0 & 0 & 1 \end{bmatrix} = \begin{bmatrix} B_{0} & 0 \\ C_{0} & D_{0} \end{bmatrix}, A_{1} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 2 & 3 & 0 & 0 & 0 \\ -1 & 4 & 2 & 1 & 0 \\ 10 & 4 & 0 & 1 & 2 \\ -1 & 0 & 0 & 0 & 0 \end{bmatrix} = \begin{bmatrix} B_{1} & 0 \\ C_{1} & D_{1} \end{bmatrix}$$
$$B_{0} = \begin{bmatrix} 3 & 1 \\ 0 & 2 \end{bmatrix}, \quad B_{1} = \begin{bmatrix} 1 & 0 \\ 2 & 3 \end{bmatrix},$$

$$P = \begin{bmatrix} 2/3 & 1/3 \\ -2/3 & 2/3 \end{bmatrix}, \quad P^{-1}B_0P = \begin{bmatrix} 2 & 1 \\ 0 & 3 \end{bmatrix}, \quad P^{-1}B_1P = \begin{bmatrix} 1 & -1 \\ 0 & 3 \end{bmatrix}.$$
$$\hat{\rho}(B) = 3, \ \hat{\rho}(D) = 2 \qquad \hat{\rho}(A) = \max(3, 2) = 3$$

Consequence:

$$S_{K}(x) = \sum_{\substack{|w|=K\\(0.w)_{b} \le x}} A_{w}C \qquad 0 \le x \le 1$$
$$S_{K+1}(x) = \sum_{r_{1} \le x_{1}} A_{r_{1}}Q^{K}C + A_{x_{1}}S_{K}(bx - x_{1}),$$

Proposition

Let V be an eigenvector of $Q = A_0 + \cdots + A_{b-1}$ for an eigenvalue $\rho \omega$ with $|\omega| = 1$ and $\rho \leq \hat{\rho}(A)$. Then

$$S_K(x) = \sum_{\substack{|w|=K\ (0.w)_b \leq x}} A_w V$$

is $O(r^K)$ uniformly wrt x for $r > \hat{\rho}(A) \ge \rho$.

coin tossing $(T_n)_{n \ge 1}$ i.i.d. with $\mathbf{P}(T=0) = p_0$, $\mathbf{P}(T=1) = p_1$ $p_0 + p_1 = 1, \ 0 < p_0, \ p_1 < 1$

$$X = \sum_{n \ge 1} \frac{T_n}{2^n}$$

distribution function F(x)

coin tossing $(T_n)_{n\geq 1}$ i.i.d. with $\mathbf{P}(T=0) = p_0$, $\mathbf{P}(T=1) = p_1$ $p_0 + p_1 = 1, \ 0 < p_0, \ p_1 < 1$

$$X = \sum_{n \ge 1} \frac{T_n}{2^n}$$

distribution function F(x) $0 \le x < 1/2$

$$F(x) = \mathbf{P}(X \le x) = \mathbf{P}(T_1 = 0, \sum_{n \ge 2} \frac{T_n}{2^{n-1}} \le 2x) = p_0 F(2x)$$

 $1/2 \leq x \leq 1$

$$F(x) = \mathbf{P}(X \le x) = \mathbf{P}(T_1 = 0) + \mathbf{P}(T_1 = 1, \sum_{n \ge 2} \frac{T_n}{2^{n-1}} \le 2x - 1)$$
$$= p_0 + p_1 F(2x - 1),$$

$$0 \le x < 1/2 \qquad F(x) = p_0 F(2x)$$

1/2 \le x \le 1 \qquad F(x) = p_0 + p_1 F(2x - 1)

$$0 \le x < 1/2 \qquad F(x) = p_0 F(2x)$$

1/2 \le x \le 1 \qquad F(x) = p_0 + p_1 F(2x - 1)

$$F(x) = p_0 F(2x) + p_1 F(2x - 1)$$

 $F(x) = 0 \text{ for } x \le 0 \qquad F(x) = 1 \text{ for } x \ge 1$

dilation equation two-scale difference equation





cascade algorithm

Hölder with exponent $\log_2 1 / \max(p_0, p_1)$

in wavelet theory (Daubechies)



in interpolation scheme (Dubuc, Deslauriers)

multidimensional version

Proposition

Under the hypothesis $\rho > \hat{\rho}(A)$, the dilation equation

$$\rho\omega F(x) = \sum_{0 \le r < b} A_r F(bx - r)$$

with boundary conditions

$$F(x) = 0$$
 for $x \le 0$, $F(x) = V$ for $x \ge 1$,

where V is an eigenvector for $Q = A_0 + \cdots + A_{b-1}$ and the eigenvalue $\rho\omega$, $|\omega| = 1$, has a unique continuous solution from \mathbb{R} into \mathbb{C}^d . Moreover this solution is Hölder with exponent $\log_b(\rho/r)$ for $r > \hat{\rho}(A)$.

Consequence:

Proposition

Let V be an eigenvector for an eigenvalue $\rho\omega$, $|\omega| = 1$, $\rho > \hat{\rho}(A)$, of $Q = A_0 + \cdots + A_{b-1}$. Then

$$S_K(x) = \sum_{\substack{|w|=K\ (0.w)_b \leq x}} A_w V$$

satisfies

$$S_K(x) \underset{K \to \infty}{=} (\rho \omega)^K F(x) + O(r^K)$$

for $\rho > r > \hat{\rho}(A)$ uniformly wrt x.

Theorem

Theorem

Let (u_n) be a b-rational sequence and (L, A, C) a linear representation for the sequence of its backward differences. Then the sequence (u_n) has an asymptotic expansion which is a sum of terms

$$N^{\log_b \rho} \binom{\log_b N}{m} \times e^{i\vartheta \log_b N} \times \varphi(\log_b N).$$

In this writing, $\rho e^{i\vartheta}$ is an eigenvalue of $Q = A_0 + A_1 + \dots + A_{b-1}$ with a modulus $\rho > \hat{\rho}(A)$. The integer m is bounded by the maxima size of the Jordan blocks related to $\rho e^{i\vartheta}$. The function $\varphi(t)$ is 1-periodic and Hölder with exponent $\log_b(\rho/r)$ for $\rho > r > \hat{\rho}(A)$. The error term is $O(N^{\log_b r})$ for $r > \hat{\rho}(A)$.

A worked example

Karatsuba!

$$xu(x) - (1+x)(2+x)u(x^2) = -x^2 + 4\frac{x^3}{(1-x)^2}$$
$$\delta(x) = \frac{(2+x)}{x}\delta(x^2) - \frac{x-6x^2+x^3}{1-x}.$$

basis

$$\delta(x), \frac{\delta(x)}{x}, \frac{1}{1-x}, \frac{x}{1-x}, \frac{x^2}{1-x}, \frac{x^3}{1-x}$$

A worked example: Linear representation

A worked example: Joint spectral radius

 $\hat{\rho}(A) = 2$

A worked example: Jordan reduction

$$P = \begin{bmatrix} 0 & 0 & 0 & 24 & 0 & 0 \\ 16 & 0 & 96 & -96 & -24 & 0 & 0 \\ 16 & 0 & 0 & 0 & -24 & 0 & 0 \\ 16 & 0 & 96 & -96 & -48 & -334 \\ 0 & 0 & 0 & 0 & 0 & -24 \end{bmatrix}, \quad Q' = \begin{bmatrix} 3 & 0 & 0 & 0 & 0 & 0 \\ 0 & 2 & 0 & 0 & 0 & 0 & 0 \\ 3 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 24 \end{bmatrix}$$
A worked example: Jordan reduction

 $\rho = 3, \quad \omega = 1, \quad V = E_2 + 6E_3 + 2E_4$

A worked example: Dilation equation

$$3F(x) = A_0F(2x) + A_1F(2x-1).$$

F(x) = 0 for $x \le 0$, $F(x) = E_2 + 6E_3 + 2E_4$ for $x \ge 1$.

A worked example: Dilation equation

$$\begin{split} f_1\left(x\right) &= \frac{2}{3}f_1\left(2x\right) - f_4\left(2x\right), \\ f_2\left(x\right) &= \frac{1}{9}f_1\left(2x\right) + \frac{1}{3}f_2\left(2x\right) - \frac{2}{3}f_4\left(2x\right) + \frac{2}{3}f_1\left(2x-1\right) + \frac{1}{3}f_2\left(2x-1\right), \\ f_3\left(x\right) &= \frac{1}{12}f_1\left(2x\right) + \frac{1}{3}f_3\left(2x\right) + \frac{1}{12}f_4\left(2x\right) - \frac{1}{12}f_5\left(2x\right) - \frac{155}{288}f_6\left(2x\right) \\ &\quad - \frac{1}{3}f_1\left(2x-1\right) + \frac{5}{3}f_2\left(2x-1\right) + \frac{1}{3}f_3\left(2x-1\right) + \frac{1}{3}f_4\left(2x-1\right), \\ f_4\left(x\right) &= \frac{1}{3}f_4\left(2x\right) + \frac{5}{3}f_1\left(2x-1\right) - \frac{1}{3}f_2\left(2x-1\right) + \frac{1}{3}f_5\left(2x-1\right) + \frac{1}{3}f_6\left(2x-1\right), \\ f_5\left(x\right) &= \frac{1}{3}f_6\left(2x\right), \\ f_6\left(x\right) &= 0. \end{split}$$

$$f_{j}(x) = 0 \quad \text{for } x \le 0 \qquad \qquad \begin{array}{l} f_{1}(x) = 0 \\ f_{2}(x) = 1 \\ f_{3}(x) = 6 \\ f_{4}(x) = 2 \\ f_{5}(x) = 0 \\ f_{6}(x) = 0 \end{array} \qquad \qquad \begin{array}{l} \text{for } x \ge 1 \\ f_{5}(x) = 0 \\ f_{6}(x) = 0 \end{array}$$

 $f_{1}(m) = 0$

A worked example: Cascade algorithm



 $f(x) = LF(x) = f_2(x) + f_3(x)$

A worked example: !



What I did not speak about

▶ analytic number theory

- Michael Drmota and Peter J. Grabner.
 Analysis of digital functions and applications.
 In Combinatorics, automata and number theory, volume 135 of Encyclopedia Math. Appl., pages 452–504. Cambridge Univ.
 Press, Cambridge, 2010.
- probability theory
 - Louis H.Y. Chen, Hsien-Kuei Hwang, and Vytas Zacharovas. Distribution of the sum-of-digits function of random integers: a survey.

Probababilty Surveys, 11:177–236, 2014.

Thanks for your attention!



Philippe Dumas SpecFun INRIA Saclay