

# Creative Telescoping

—From 1G to 4G algorithms—

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## Several generations of Creative Telescoping algorithms

- 1G, *brutal elimination*: [Fasenmyer, 1947], [Lipshitz, 1988], [Zeilberger, 1990], [Takayama, 1990], [Wilf, Zeilberger, 1990], [Chyzak, Salvy, 2000]
  - 2G, *linear diff/rec rational solving*: [Zeilberger, 1990], [Zeilberger, 1991], [Almkvist, Zeilberger, 1990], [Chyzak, 2000], [Koutschan, 2010]
  - 3G, combines  $1G + 2G + \text{linear algebra}$ : [Apagodu, Zeilberger, 2005], [Koutschan 2010], [Chen, Kauers 2012], [Chen, Kauers, Koutschan 2014]
- ▷ Advantages:
- 1G–3G: very general algorithms;
  - 2G/3G algorithms are able to solve non-trivial problems.
- ▷ Drawbacks:
- 1G: slow;
  - 2G: bad or unknown complexity;
  - 1G and 3G: non-minimality of telescopers;
  - 1G–3G: all compute (big) certificates.

# Several generations of Creative Telescoping algorithms

4G: roots in [Ostrogradsky, 1845], [Hermite, 1872] and [Picard, 1902].

- univariate:

- rational  $f$ : [B., Chen, Chyzak, Li, 2010];
- hyperexponential  $f$ : [B., Chen, Chyzak, Li, Xin, 2013]
- hypergeometric  $\Sigma$ : [Chen, Huang, Kauers, Li, 2015], [Huang, 2016]
- mixed  $f + \Sigma$ : [B., Dumont, Salvy, 2016]
- algebraic  $f$ : [Chen, Kauers, Koutschan, 2016]
- D-finite Fuchsian  $f$ : [Chen, van Hoeij, Kauers, Koutschan, 2018]
- D-finite  $f$ : [B., Chyzak, Lairez, Salvy, 2018], [van der Hoeven, 2018]

- multiple:

- rational bivariate  $\mathcal{F}\mathcal{F}$ : [Chen, Kauers, Singer, 2012]
- rational: [B., Lairez, Salvy, 2013], [Lairez 2016]
- binomial sums: [B., Lairez, Salvy, 2017]

▷ Advantages:

- complexity;
- minimality of telescopers;
- does not need to compute certificates;
- fast in practice,

▷ Drawback: not (yet) as general as 1G–3G algorithms.

**Input:** a **hyperexponential** function  $H(t, x)$ , i.e.,  $\frac{\partial H}{\partial t}$  and  $\frac{\partial H}{\partial x}$  are in  $\mathbb{Q}(t, x)$

**Output:**

- a linear differential operator  $P(t, \partial_t)$  satisfied by  $I(t) = \oint_{\gamma} H(t, x) dx$
- a  $G(t, x) \in \mathbb{Q}(t, x)$  such that  $P(t, \partial_t) \cdot H(t, x) = \frac{\partial}{\partial x} (G(t, x) \cdot H(t, x))$ .

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**Algorithm:** Write  $\mathbb{L} = \mathbb{Q}(t)$ . For  $r = 0, 1, 2, \dots$  do

- ① compute  $a(x) := \frac{\partial H}{\partial x} \in \mathbb{L}(x)$  and  $b_k(x) := \frac{\partial^k H}{\partial t^k} \in \mathbb{L}(x)$  for  $k = 0, \dots, r$
- ② rational solving: decide whether the inhomogeneous parametrized LDE

$$\frac{\partial G}{\partial x} + a(x) \cdot G = \sum_{k=0}^r c_k \cdot b_k(x)$$

admits a rational solution  $G \in \mathbb{L}(x)$ , for some  $c_0, \dots, c_r \in \mathbb{L}$  not all zero

- ③ if so, then return  $P := \sum_{k=0}^r c_k \partial_t^k$  and  $G$ ; else increase  $r$  by 1 and repeat

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$$H = \partial_x(g) + \frac{a}{Q^*},$$

where  $Q^*$  is the squarefree part of  $Q$  and  $\deg_x(a) < d^* := \deg_x(Q^*)$ .

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**Algorithm** [B., Chen, Chyzak, Li, 2010]

(1) For  $i = 0, 1, \dots, d^*$  compute Hermite reduction of  $\partial_t^i(H)$ :

$$\partial_t^i(H) = \partial_x(g_i) + \frac{a_i}{Q^*}, \quad \deg_x(a_i) < d^*$$



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▷  $L = \sum_{k=0}^r c_k \partial_t^k$  is a **telescoper** (and  $\sum_{k=0}^r c_k g_k$  the corresponding certificate).

Compute a telescoper for the diagonal of the rational power series

$$\frac{1}{1-x-y} = \sum_{i,j \geq 0} \binom{i+j}{i} x^i y^j$$

in three different ways:

- ① using the 2G (Almkvist-Zeilberger) creative telescoping algorithm;
- ② using the 4G (Hermite reduction-based) creative telescoping algorithm;
- ③ using the 4G (Griffiths-Dwork reduction-based) creative telescoping algorithm.

## An exercise from a past lecture

Let  $(a_n)_{n \geq 0}$  be a sequence with  $a_0 = a_1 = 1$  satisfying the recurrence

$$(n+3)a_{n+1} = (2n+3)a_n + 3na_{n-1}, \quad \text{for all } n > 0.$$

Show that  $a_n$  is an integer for all  $n$ .

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Follow the next steps:

- 1 Compute the first 5 terms of the sequence,  $a_0, \dots, a_4$ ;
- 2 Determine a Hermite-Padé approximant of type  $(0, 1, 2)$  for  $(1, f, f^2)$ , where  $f = \sum_n a_n x^n$ ;
- 3 Deduce that  $P(x, f(x)) = 0 \pmod{x^5}$  for  $P(x, y) := 1 + (x-1)y + x^2y^2$ ;
- 4 Show that the equation  $P(x, y) = 0$  admits a root  $y = g(x) \in \mathbb{Q}[[x]]$  whose coefficients satisfy the same linear recurrence as  $(a_n)_{n \geq 0}$ ;
- 5 Deduce that  $a_{n+2} = a_{n+1} + \sum_{k=0}^n a_k \cdot a_{n-k}$  for all  $n$ , and conclude.

- ▷ 1. Let's compute the first 5 terms of the sequence:

```
> rec:=(n+3)*a(n+1)-(2*n+3)*a(n)-3*n*a(n-1): ini:=a(0)=1,a(1)=1:  
> pro:=gfun:-rectoproc({rec,ini}, a(n), list);  
> pro(4);
```

[1, 1, 2, 4, 9]

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- ▷ 2. Determine a Hermite-Padé approximant of type  $(0, 1, 2)$  for  $(1, f, f^2)$ :

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> f:=listtoseries(pro(4),x):  
> numapprox[hermite_pade]([1, f, f^2], x, [0, 1, 2]);
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- ▷ 3. We guessed  $P(x, f(x)) = 0 \pmod{x^5}$  for  $P(x, y) = 1 + (x - 1)y + x^2y^2$ .



▷ 4. The equation  $P(x, y) = 0$  admits a root  $y = g(x) = \sum_{n \geq 0} g_n x^n$  in  $\mathbb{Q}[[x]]$ :

$$g(x) = \frac{1 - x - \sqrt{1 - 2x - 3x^2}}{2x^2} = 1 + x + 2x^2 + 4x^3 + 9x^4 + \dots$$

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$$h(x) := \sqrt{1 - 2x - 3x^2} = 1 - x - 2x^2 g(x).$$

Clearly,  $h_0 = 1$ ,  $h_1 = -1$ ,  $h_{n+2} = -2g_n$  for  $n \geq 0$ , and

$$\frac{h'(x)}{h(x)} = \frac{3x + 1}{3x^2 + 2x - 1}$$

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- ▷ Conclusion:  $a_n = g_n$  is an integer for all  $n$ . □

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