

# Polynomial and rational solutions of linear differential/recurrence equations

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**M2, CR11**

**November 18, 2019**

## The exercise from last lecture

Let  $(a_n)_{n \geq 0}$  be a sequence with  $a_0 = a_1 = 1$  satisfying the recurrence

$$(n+3)a_{n+1} = (2n+3)a_n + 3na_{n-1}, \quad \text{for all } n > 0.$$

Show that  $a_n$  is an integer for all  $n$ .

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Show that  $a_n$  is an integer for all  $n$ .

Follow the next steps:

- 1 Compute the first 5 terms of the sequence,  $a_0, \dots, a_4$ ;
- 2 Determine a Hermite-Padé approximant of type  $(0, 1, 2)$  for  $(1, f, f^2)$ , where  $f = \sum_n a_n x^n$ ;
- 3 Deduce that  $P(x, f(x)) = 0 \pmod{x^5}$  for  $P(x, y) := 1 + (x-1)y + x^2y^2$ ;
- 4 Show that the equation  $P(x, y) = 0$  admits a root  $y = g(x) \in \mathbb{Q}[[x]]$  whose coefficients satisfy the same linear recurrence as  $(a_n)_{n \geq 0}$ ;
- 5 Deduce that  $a_{n+2} = a_{n+1} + \sum_{k=0}^n a_k \cdot a_{n-k}$  for all  $n$ , and conclude.

- ▷ 1. Let's compute the first 5 terms of the sequence:

```
> rec:=(n+3)*a(n+1)-(2*n+3)*a(n)-3*n*a(n-1): ini:=a(0)=1,a(1)=1:  
> pro:=gfun:-rectoproc({rec,ini}, a(n), list);  
> pro(4);
```

[1, 1, 2, 4, 9]

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- ▷ 2. Determine a Hermite-Padé approximant of type  $(0, 1, 2)$  for  $(1, f, f^2)$ :

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> f:=listtoseries(pro(4),x):  
> numapprox[hermite_pade]([1, f, f^2], x, [0, 1, 2]);
```

$[-1, -x + 1, -x^2]$

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- ▷ 3. We guessed  $P(x, f(x)) = 0 \pmod{x^5}$  for  $P(x, y) = 1 + (x - 1)y + x^2y^2$ .

▷ 4. The equation  $P(x, y) = 0$  admits a root  $y = g(x) = \sum_{n \geq 0} g_n x^n$  in  $\mathbb{Q}[[x]]$ :

$$g(x) = \frac{1 - x - \sqrt{1 - 2x - 3x^2}}{2x^2} = 1 + x + 2x^2 + 4x^3 + 9x^4 + \dots$$

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$$g_{n+2} = g_{n+1} + \sum_{k=0}^n g_k \cdot g_{n-k}.$$



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$$h(x) := \sqrt{1 - 2x - 3x^2} = 1 - x - 2x^2 g(x).$$

Clearly,  $h_0 = 1$ ,  $h_1 = -1$ ,  $h_{n+2} = -2g_n$  for  $n \geq 0$ , and

$$\frac{h'(x)}{h(x)} = \frac{3x + 1}{3x^2 + 2x - 1}$$

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- ▷ Conclusion:  $a_n = g_n$  is an integer for all  $n$ . □

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> deq:=gfun:-algeqtodiffeq(pol,y(x));
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$$\{(3n+3)g_n + (2n+5)g_{n+1} - (n+4)g_{n+2} = 0, g_0 = 1, g_1 = 1\}$$

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- ▷ Conclusion:  $a_n = g_n$  is an integer for all  $n$ . □

## Resolution of Linear Differential Equations (LDEs):

- homogeneous LDE

$$c_m(x)y^{(m)}(x) + \cdots + c_0(x)y(x) = 0, \quad c_i \in \mathbb{K}[x]$$

- inhomogeneous LDE

$$c_m(x)y^{(m)}(x) + \cdots + c_0(x)y(x) = b(x), \quad c_i, b \in \mathbb{K}[x]$$

- inhomogeneous parametrized LDE

$$c_m(x)y^{(m)}(x) + \cdots + c_0(x)y(x) = \sum_{j=0}^S \lambda_j b_j(x) \quad \lambda_j \in \mathbb{K}, c_i, b_j \in \mathbb{K}[x]$$

▷ Input in blue, output in red

## Resolution of Linear Recurrence Equations (LREs):

- homogeneous LRE

$$p_r(n)u_{n+r} + \cdots + p_0(n)u_n = 0, \quad p_i \in \mathbb{K}[n]$$

- inhomogeneous LRE

$$p_r(n)u_{n+r} + \cdots + p_0(n)u_n = b(n), \quad p_i, b \in \mathbb{K}[n]$$

- inhomogeneous parametrized LRE

$$p_r(n)u_{n+r} + \cdots + p_0(n)u_n = \sum_{j=0}^S \lambda_j b_j(n) \quad \lambda_j \in \mathbb{K}, p_i, b_j \in \mathbb{K}[n]$$

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- ▷ **Tasks:** solve all these equations for
  - **power series** solutions of LDEs, **finite-support** solutions of LREs
  - **polynomial** and **rational** solutions → central tool in  $\Sigma$  and  $\int$  algos
- ▷ **Sub-tasks:**
  - **structure** of space of solutions
  - **complexity** issues
  - inhomogeneous parametrized: **existence** and **values** of parameters  $\lambda_i$  for which there exist polynomial / rational solutions
- ▷ **Main focus:** **homogeneous case** (the other cases reduce to it)

▷ **Basic, but important, observation:** polynomial/rational solutions may have **exponentially large degree** with respect to the bit-size of the input equation

- E.g.,  $u_n = n(n+1) \cdots (n+10^{10}-1)$  is a degree- $10^{10}$  solution of LRE

$$nu_{n+1} - (n+10^{10})u_n = 0$$

- E.g.,  $y(x) = \sum_{k=0}^{10^{10}} \binom{10^{10}}{k} x^k$  is a degree- $10^{10}$  solution of LDE

$$(x+1)y'(x) - 10^{10}y(x) = 0$$

▷ In both cases, when written in the monomial basis of  $\mathbb{K}[n]$ , resp. of  $\mathbb{K}[x]$ , outputs have degree  $\tilde{O}(N)$  and bit-size  $\tilde{O}(N^2)$ , for inputs of bit-size  $\log N$ .

**Theorem.** The set of sequences  $(u_n)_{n \geq 0}$  solutions of the LRE

$$\text{(LRE):} \quad p_r(n)u_{n+r} + \cdots + p_0(n)u_n = 0$$

is a  $\mathbb{K}$ -vector space of finite dimension, lying between  $r$  and  $r + |H|$ , where

$$H := \{h \in \mathbb{N} \mid p_r(h) = 0\}.$$

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**Remarks:**

- ▷ When  $p_r \in \mathbb{K} \setminus \{0\}$  (e.g., for r.l.s.c.c.), this dimension is exactly  $r$ .
  - ▷ Same is true for all **regular** recurrences, that is when  $H = \emptyset$ .
  - ▷ A **singular** example:  $n(n-1)u_{n+3} - (n-1)u_{n+2} + (n+1)u_{n+1} + 2nu_n = 0$ 
    - for  $n = 0 \Rightarrow u_2 + u_1 = 0$  and  $u_3$  is free
    - for  $n = 1 \Rightarrow 2u_2 + 2u_1 = 0$  and  $u_4$  is free
    - for  $n \geq 2 \Rightarrow$  recurrence forces value of  $u_{n+3}$  in terms of previous ones
- So **dim = 4**:  $u_0, u_1, u_3, u_4$  are free, all other are linear combinations of them.

**Theorem.** The set of sequences  $(u_n)_{n \geq 0}$  solutions of the LRE

$$\text{(LRE):} \quad p_r(n)u_{n+r} + \cdots + p_0(n)u_n = 0$$

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    - for  $n \geq 2 \Rightarrow$  recurrence forces value of  $u_{n+3}$  in terms of previous ones
- So **dim = 3**:  $u_0, u_3, u_4$  are free, all other are linear combinations of them.

**Theorem.** The set of sequences  $(u_n)_{n \geq 0}$  solutions of the LRE

$$\text{(LRE):} \quad p_r(n)u_{n+r} + \cdots + p_0(n)u_n = 0$$

is a  $\mathbb{K}$ -vector space of finite dimension, lying between  $r$  and  $r + |H|$ , where

$$H := \{h \in \mathbb{N} \mid p_r(h) = 0\}.$$

**Proof:**

- Let  $J := \{0, \dots, r-1\} \cup \{h+r \mid h \in H\}$ . Note:  $|J| = r + |H|$ .
- For  $n \in \mathbb{N} \setminus J$ , value of  $u_{n+r}$  is uniquely determined from previous ones.
- $\dim \Lambda_\ell = \dim \Lambda_{\ell-1}$  for  $\ell \notin J$ , and  $\dim \Lambda_j \leq 1 + \dim \Lambda_{j-1}$  for  $j \in J$ , where

$$\Lambda_\ell := \left\{ (u_0, \dots, u_\ell) \in \mathbb{K}^{\ell+1} \mid \text{(LRE) true for } n \in \{0, \dots, \ell-r\} \right\}.$$

- $\dim \Lambda_0 = 1, \dots, \dim \Lambda_{r-1} = r$  and  $\dim \text{Sols}(\text{LRE}) = \dim \Lambda_{\max(J)}$  □

▷ Linear algebra in dimension  $\max(J)$   $\longrightarrow$  basis of solution space  $\text{Sols}(\text{LRE})$ .

**Theorem.** If  $y(x) = \sum_{n \geq 0} u_n x^n \in \mathbb{K}[[x]]$  is a solution of the LDE

$$\text{(LDE):} \quad c_m(x)y^{(m)}(x) + \cdots + c_0(x)y(x) = 0,$$

with  $c_i(x) = \sum_{j=0}^d c_{i,j}x^j$  in  $\mathbb{K}[x]$ , then the sequence  $(u_n)_{n \geq 0}$  satisfies the LRE

$$\sum_{i=0}^m \sum_{j=0}^d c_{i,j}(n+i-j) \cdots (n+1-j)u_{n+i-j} = 0,$$

for all  $n \in \mathbb{Z}$ , with  $u_\ell = 0$  for  $\ell < 0$ . This writes for some  $0 \leq r \leq m+d$ :

$$p_r(n)u_{n+r} + \cdots + p_0(n)u_n = 0,$$

together with the equations (translating the constraints  $u_\ell = 0$  for  $\ell < 0$ ):

$$p_r(-1)u_{r-1} + \cdots + p_1(-1)u_0 = 0, \quad p_r(-2)u_{r-2} + \cdots = 0, \dots, \quad p_r(-r)u_0 = 0.$$

**Proof:** Extract coefficient of  $x^n$  in  $\text{LDE}(y(x)) = 0$ .

Let, as before, the linear differential equation

$$\text{(LDE): } c_m(x)y^{(m)}(x) + \cdots + c_0(x)y(x) = 0,$$

with associated linear recurrence equation

$$\text{(LRE): } p_r(n)u_{n+r} + \cdots + p_0(n)u_n = 0.$$

**Def.** Polynomial  $p_r(n-r)$  in  $\mathbb{K}[x]$  is called **indicial polynomial at 0** of (LDE)

**Def.** Polynomial  $p_0(n)$  in  $\mathbb{K}[x]$  is called **indicial polynomial at infinity** of (LDE)

Let  $\text{(LDE}_\alpha\text{): } c_m(x+\alpha)y^{(m)}(x) + \cdots + c_0(x+\alpha)y(x) = 0$  for any  $\alpha \in \overline{\mathbb{K}}$ .

**Def.** **Indicial polynomial at  $x = \alpha$**  of (LDE) is indicial poly. **at  $x = 0$**  of  $\text{(LDE}_\alpha\text{)}$ .

▷ Indicial polynomials at infinity of LDE and of  $\text{(LDE}_\alpha\text{)}$  coincide.



$$\text{(LDE): } c_m(x)y^{(m)}(x) + \cdots + c_0(x)y(x) = 0, \quad \text{(LRE): } p_r(n)u_{n+r} + \cdots + p_0(n)u_n = 0.$$

### Theorem.

- (1) If  $F \in \mathbb{K}[[x]]$  is a solution of (LDE) then its valuation  $\text{val}(F)$  is a root of the indicial polynomial at 0 of (LDE).
- (2) (Cauchy) If  $c_m(0) \neq 0$  then (LDE) admits a basis of  $m$  solutions in  $\mathbb{K}[[x]]$ , with valuations  $0, 1, \dots, m - 1$ .

$$\text{(LDE): } c_m(x)y^{(m)}(x) + \cdots + c_0(x)y(x) = 0, \quad \text{(LRE): } p_r(n)u_{n+r} + \cdots + p_0(n)u_n = 0.$$

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### Proof:

- (1) Evaluate (LRE) at  $n = \text{val}(F) - r$ .
- (2) Assume  $c_m(0) \neq 0$ . Then (LRE) reads

$$0 = c_m(0) \cdot (n+m) \cdots (n+1)u_{n+m} + (\text{terms in } u_{<n+m})$$

- ▷ Thus, (LRE) has order at least  $m$ , and it is regular, so there exist  $y_i = x^i + O(x^m)$  solutions of (LDE) for  $i = 0, 1, \dots, m-1$ .
- ▷ The indicial polynomial at 0 equals  $p_r(n-r) = c_m(0) \cdot n \cdots (n-m+1)$ , so by (1) these  $y_i$  form a basis of solutions of (LDE).

$$(\text{LDE}): c_m(x)y^{(m)}(x) + \dots + c_0(x)y(x) = 0, \quad (\text{LRE}): p_r(n)y_{n+r} + \dots + p_0(n)y_n = 0.$$

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▷ Ex: The converse of (2) is also true: if (LDE) admits a basis of  $m$  solutions in  $\mathbb{K}[[x]]$  with valuations  $0, 1, \dots, m - 1$ , then  $c_m(0) \neq 0$ .

▷ When  $c_m(0) \neq 0$ , one says that  $x = 0$  is an **ordinary point** for (LDE).

**Corollary:** If  $x = 0$  is an ordinary point for (LDE), then one can compute a **full basis of power series solutions mod  $x^N$**  in  $O(N)$  operations in  $\mathbb{K}$ .

## Series solutions of LDEs: an example

(LDE) of order  $m = 3$

$$(x-1)(x+1)y'''(x) + y'(x) + (x^2-1)y(x) = 0$$

▷ Corresponding (LRE) of order  $r = 5$

$$-(n+3)(n+2)(n+1)u_{n+3} + (n+1)(n^2-n+1)u_{n+1} - u_n + u_{n-2} = 0$$

▷ Indicial equation at 0 equals  $-n(n-1)(n-2)$ .

▷ Basis of solutions of (LDE):

$$\begin{aligned} & 1 + 0 \cdot x + 0 \cdot x^2 - \frac{1}{6} \cdot x^3 + 0 \cdot x^4 - \frac{1}{120} \cdot x^5 + O(x^6), \\ & 0 + 1 \cdot x + 0 \cdot x^2 + \frac{1}{6} \cdot x^3 - \frac{1}{24} \cdot x^4 + \frac{1}{40} \cdot x^5 + O(x^6), \\ & 0 + 0 \cdot x + 1 \cdot x^2 + 0 \cdot x^3 + \frac{1}{12} \cdot x^4 - \frac{1}{60} \cdot x^5 + O(x^6). \end{aligned}$$

## Solutions with finite support of LRE

**Def:** A sequence  $(u_n)_{n \geq 0}$  **has finite support** if  $\exists N \geq 0$  s.t.  $u_n = 0$  for  $n > N$ . We call  $N$  the **degree** of  $(u_n)_{n \geq 0}$  if  $u_N \neq 0$ .

**Task:** Find solutions with finite support of

$$\text{(LRE):} \quad p_r(n)u_{n+r} + \cdots + p_0(n)u_n = 0.$$

- ▷ The degree of such a solution must be a root of  $p_0(n)$ .
- ▷ Solution space: vector space of  $\dim \leq$  number of roots in  $\mathbb{N}$  of  $p_0(n)$ .
- ▷ Algorithm:
  - Compute  $S := \{h \in \mathbb{N} \mid p_0(h) \cdot p_r(h) = 0\}$
  - If  $S = \emptyset$ , return no solution; if not, let  $N := \max S$ .
  - Solve a linear system:
    - equations: eval. (LRE) at  $n = 0, 1, \dots, N$ ; add  $u_{N+i} = 0$  for  $i = 1, \dots, r$
    - unknowns:  $u_0, u_1, \dots, u_{N+r}$
- ▷ Complexity:  $O(N^\theta)$
- ▷ One can do better:  $O(N + |K|^\theta)$

**Task:** Given  $c_0, \dots, c_m \in \mathbb{K}[x]$ , compute a basis of solutions in  $\mathbb{K}[x]$  of

$$\text{(LDE):} \quad c_m(x)y^{(m)}(x) + \dots + c_0(x)y(x) = 0.$$

- ▷ If a bound  $N$  on the degree of a solution  $y(x) \in \mathbb{K}[x]$  is known, then the coefficients of  $y(x)$  can be computed by linear algebra
- ▷ Such a bound on the degree is the largest root  $N$  in  $\mathbb{N}$  of the indicial polynomial of (LDE) at infinity
- ▷ One can compute a basis of polynomial solutions:
  - in  $O(N^\theta)$  by dense linear algebra
  - in  $O(N)$  by exploiting the structure (banded matrix)
- ▷ One can decide *existence* of nonzero polynomial solutions faster!  $\tilde{O}(\sqrt{N})$

## Existence of polynomial solutions of LDEs

**Task:** Given  $c_0, \dots, c_m \in \mathbb{K}[x]$ , decide if there exist solutions in  $\mathbb{K}[x] \setminus \{0\}$  of

$$\text{(LDE):} \quad c_m(x)y^{(m)}(x) + \dots + c_0(x)y(x) = 0.$$

- ▷ Up to a shift, one may assume  $c_m(0) \neq 0$ , i.e.  $x = 0$  is an ordinary point
- ▷ Let  $\{y_1, \dots, y_m\} \subset \mathbb{K}[[x]]$  be a basis of power series solutions of (LDE) (and  $N$  the largest root in  $\mathbb{N}$  of the indicial polynomial of (LDE) at infinity).
- ▷ Any solution in  $\mathbb{K}[x]_{\leq N}$  of (LDE) is a  $\mathbb{K}$ -linear combination of the  $y_i$ 's
- ▷ With  $d = \max(\deg c_i)$ , one can choose

$$y_i = x^i + u_{i,m}x^m + \dots + u_{i,N}x^N + \dots + u_{i,N+m+d}x^{N+m+d} + \dots$$

- ▷ One can compute coefficients in blue in  $\tilde{O}(\sqrt{N})$  by baby-steps/giant-steps
- ▷  $y = \sum_i \lambda_i y_i$  is in  $\mathbb{K}[x]_{\leq N} \iff \sum_{i=1}^m \lambda_i u_{i,N+j} = 0$  for all  $1 \leq j \leq m+d$
- ▷ Computing the kernel does not depend on  $N$

**Task:** Given  $p_0, \dots, p_r \in \mathbb{K}[n]$ , compute a basis of solutions in  $\mathbb{K}[n]$  of

$$\text{(LRE): } p_r(n)u_{n+r} + \dots + p_0(n)u_n = 0.$$

▷ Same principle as for LDEs: “degree bound + linear algebra”

**Main sub-task:** Degree bound

▷ With  $(\Delta \cdot u)(n) = u(n+1) - u(n)$ , rewrite recurrence equation (LRE) as

$$\text{(LRE): } L \cdot u = \sum_{k=0}^s b_k(n) \Delta^k u = 0.$$

▷ **Advantage:**  $\Delta$  decreases degrees by 1, thus  $\deg(\Delta^k P) \leq \max(\deg P - k, 0)$ .

▷ **Corollary:**  $\deg(L \cdot P) \leq B + \deg P$ , with  $B := \max_k \{\deg(b_k) - k\}$

▷ If  $P = n^D + \dots \in \mathbb{K}[n]$  is a solution of (LRE), then

- either  $D + B < 0$ , thus  $D$  is upper bounded by  $-(B + 1)$ ;
- or  $[n^{D+B}] (L \cdot P) = \sum_{\deg b_k - k = B} \text{lc}(b_k) D(D-1) \dots (D-k+1)$  is zero



## Bound on Degree for Polynomial solutions of LREs

**Task:** Given  $p_0, \dots, p_r \in \mathbb{K}[n]$ , bound the degree of polynomial solutions of

$$\text{(LRE):} \quad p_r(n)u_{n+r} + \dots + p_0(n)u_n = 0.$$

▷ With  $(\Delta \cdot u)(n) = u(n+1) - u(n)$ , rewrite recurrence equation (LRE) as

$$\text{(LRE):} \quad L \cdot u = \sum_{k=0}^s b_k(n) \Delta^k u = 0.$$

▷ **Theorem.** A bound on the degree of polynomial solutions of (LRE) is

$$\max \left( - (B + 1), \text{largest root } D \in \mathbb{N} \text{ of } \underbrace{\sum_{k \in E} \text{lc}(b_k) D(D-1) \cdots (D-k+1)}_{=: \text{indicial polynomial of (LRE)}} \right)$$

where

- $B := \max_k \{ \deg(b_k) - k \}$ , weighted  $(n, \Delta)$ -degree of  $L$ , for  $\deg(\Delta) = -1$ ;
- $E := \{ k \mid \deg(b_k) - k = B \}$ , indices where weighted degree is reached.

## Polynomial solutions of LREs: an example

**Question:** Find all polynomial solutions of

$$(n-3)u_{n+2} - (2n-3)u_{n+1} + nu_n = 0$$

▷ The recurrence can be rewritten:

$$(n-3)(\Delta^2 u_n + 2\Delta u_n + u_n) - (2n-3)(\Delta u_n + u_n) + nu_n = 0,$$

that is

$$L(u) = \left( (n-3)\Delta^2 - 3\Delta \right) (u_n) = 0$$

▷  $B$  is  $-1$  and the set  $E$  is  $\{1, 2\}$ .

▷ Indicial polynomial is  $D(D-1) - 3D = D(D-4)$

▷ Degree bound is  $4$ ; solution  $P(n) = n^4 + an^3 + bn^2 + cn + d$

▷  $L(P) = (-3a - 30)n^2 + (-21a - 70 - 4b)n - 21a - 45 - 9b - 3c = 0$ ,  
with solution  $a = -10$ ,  $b = 35$ ,  $c = -50$ , and  $d$  is free.

▷ Basis of solutions given by

$$\left\{ (n-1)(n-2)(n-3)(n-4), \quad n(n-5)(n^2 - 5n + 10) \right\}$$

**Remark:** Degree bound can be exponentially large in the bit size of (LRE)!

▷ E.g.,  $nu_{n+1} - (n + 100)u_n = 0$  has solution  $u_n = n(n + 1) \cdots (n + 99)$ .

▷ For LDEs, we exploited the fact that coefficients sequence of a series solution satisfies a LRE. **This is false for LREs in the monomial basis.** Still:

**Theorem.** If  $(u_n)$  is a **polynomial solution** of a LRE of order  $r$  with coefficients of degree at most  $d$ , then its coefficient sequence  $(c_k)$  in the **binomial basis**  $\left\{\binom{n}{k}, k \in \mathbb{N}\right\}$  is a solution with finite support of another LRE, of order at most  $d + r$ , degree at most  $d$ , with no additional singularities.

**Remark:** Degree bound can be exponentially large in the bit size of (LRE)!

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**Proof:**

$$u_n = \sum_{k=0}^d c_k \binom{n}{k} \implies nu_n = \sum_{k=0}^{d+1} k(c_k + c_{k-1}) \binom{n}{k}, \quad u_{n+1} = \sum_{k=0}^d (c_k + c_{k+1}) \binom{n}{k}.$$

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**Proof:**

$$u_n = \sum_{k=0}^d c_k \binom{n}{k} \implies nu_n = \sum_{k=0}^{d+1} k(c_k + c_{k-1}) \binom{n}{k}, \quad u_{n+1} = \sum_{k=0}^d (c_k + c_{k+1}) \binom{n}{k}.$$

▷ **Example:** for  $nu_{n+1} - (n + 100)u_n = 0$ , we have

$$k(c_k + c_{k+1} + c_{k-1} + c_k) - k(c_k + c_{k-1}) - 100c_k = kc_{k+1} + (k - 100)c_k = 0.$$

Find all polynomial solutions of the recurrence equation

$$3u_{n+2} - nu_{n+1} + (n - 1)u_n = 0.$$

**Task:** Given  $c_0, \dots, c_m \in \mathbb{K}[x]$ , compute a basis of solutions in  $\mathbb{K}(x)$  of

$$\text{(LDE):} \quad c_m(x)y^{(m)}(x) + \dots + c_0(x)y(x) = 0.$$

- ▷ Even assuming numerator/denominator bounds, a naive approach by undetermined coefficients would lead to a non-linear polynomial system.
- ▷ Better strategy [Liouville, 1833]:
  - ① find a **multiple of denominators** of all rational solutions  $\rightarrow$  **singularity analysis**
  - ② reduce to **polynomial solutions for numerators**  $\rightarrow$  **change of unknowns**
- ▷ Singularity analysis based on:
  - ① any solution  $y(x) \in \mathbb{K}(x)$  of (LDE) has an expansion, around any of its poles  $\alpha \in \overline{\mathbb{K}}$ , as **generalized series**  $(x - \alpha)^v \sum_i u_i (x - \alpha)^i$ , with  $v \in \mathbb{Z} \setminus \mathbb{N}$
  - ② order  $-v$  of the pole at  $x = \alpha$  is bounded by the opposite of the smallest root in  $\mathbb{Z}_-$  of the indicial polynomial at  $\alpha$

## Liouville(LDE)

**In:** A LDE  $c_m(x)y^{(m)}(x) + \dots + c_0(x)y(x) = 0$  with  $c_i \in \mathbb{K}[x]$

**Out:** A basis of its rational solutions in  $\mathbb{K}(x)$ .

- ① At each root  $\alpha$  of  $c_m$ :
  - compute the indicial polynomial  $p_\alpha(n)$ ;
  - compute the smallest root  $N_\alpha$  in  $\mathbb{Z}_-$  of  $p_\alpha$ ; if none, set  $N_\alpha := 0$ .
- ② Form the polynomial  $P = \prod_{c_m(\alpha)=0} (x - \alpha)^{-N_\alpha}$ .
- ③ Make the change of unknowns  $y = Y/P$  in (LDE)
- ④ Find a basis  $\mathcal{B}$  of polynomial solutions  $Y(x)$  of the new LDE.
- ⑤ Return  $\{b/P \mid b \in \mathcal{B}\}$ .



**Question:** Find (a bound on the degree of) all polynomial solutions of

$$(n-3)u_{n+2} - (2n-3)u_{n+1} + nu_n = 0$$

▷ If  $u_n$  polynomial solution of degree  $D$ , then  $y(x) = \sum_n u_n x^n$  is a **rational solution** of the form  $N(x)/(1-x)^{D+1}$  of the corresponding LDE:

$$\begin{aligned} & x(x-1)^2((5u_0 - 4u_1)x - 5u_0)y''(x) \\ & + 2(x-1)\left((5u_0 - 4u_1)x^2 + (5u_0 - 10u_1)x - 10u_0\right)y'(x) \\ & - 20u_1y(x) = 0 \end{aligned}$$

▷ Indicial equation at  $x = 1$  is  $(n+1)(n+5)$ , so  $D+1 \leq 5$

Consider the differential equation

$$(x - 1)(x^2 - 2)y''(x) + 2x(x^2 - x - 1)y'(x) + 4(x - 2)y(x) = 0.$$

- 1 Prove that this equation has no nonzero polynomial solution.
- 2 Show that any rational solution  $y(x) \in \mathbb{Q}(x)$  of this equation has at most a pole of order 1 at  $x = 1$ , and cannot have any other pole.
- 3 Find all rational solutions  $y(x) \in \mathbb{Q}(x)$  of this equation.

**Task:** Given  $p_0, \dots, p_r \in \mathbb{K}[n]$ , compute a basis of solutions in  $\mathbb{K}(n)$  of

$$\text{(LRE):} \quad p_r(n)u_{n+r} + \dots + p_0(n)u_n = 0.$$

▷ Even assuming numerator/denominator bounds, a naive approach by undetermined coefficients would lead to a non-linear polynomial system.

▷ Better strategy [Abramov, 1995]:

- ① find a **multiple of denominators** of all rational solutions  $\longrightarrow$  **singularity analysis**
- ② reduce to **polynomial solutions for numerators**  $\longrightarrow$  **change of unknowns**

## Rational solutions of LREs: singularity analysis

**Task:** Given  $p_0, \dots, p_r \in \mathbb{K}[n]$ , compute a basis of solutions in  $\mathbb{K}(n)$  of

$$\text{(LRE): } p_r(n)u_{n+r} + \dots + p_0(n)u_n = 0.$$

Assume  $u_n = P(n)/Q(n)$  solution of (LRE) with  $\gcd(P, Q) = 1$ . Multiplying

$$p_r(n) \frac{P(n+r)}{Q(n+r)} + \dots + p_0(n) \frac{P(n)}{Q(n)} = 0$$

by

$$M(n) = \text{lcm}(Q(n+1), \dots, Q(n+r))$$

yields

$$Q(n) \mid p_0(n)M(n).$$

▷ **Idea in simplest case:** Assume  $Q$  has no two roots with integer difference. Then  $\gcd(Q, M) = 1$ . Thus  $Q$  divides  $p_0(n)$ . Similarly,  $Q$  divides  $p_r(n-r)$ .

**Task:** Given  $p_0, \dots, p_r \in \mathbb{K}[n]$ , compute a basis of solutions in  $\mathbb{K}(n)$  of

$$\text{(LRE):} \quad p_r(n)u_{n+r} + \dots + p_0(n)u_n = 0.$$

**Theorem.** Assume  $u_n = P(n)/Q(n)$  solution of (LRE) with  $\gcd(P, Q) = 1$ . Let  $\alpha, \beta$  be roots in  $\overline{\mathbb{K}}$  of  $Q$  s.t.  $\alpha - \beta = H \in \mathbb{N}$  maximal ( $H = \text{dispersion}$  of  $Q$ ). Then  $p_0(\alpha) = 0$  and  $p_r(\beta - r) = 0$ .

**Proof:**  $\alpha + 1, \alpha + 2, \dots$ , are not roots of  $Q$ , i.e.,  $\alpha$  not pole of  $u_{n+1}, \dots, u_{n+r}$ . So  $\alpha$  is not a pole of  $p_0(n)u_n$ . Since  $\alpha$  is a pole of  $u_n$ , necessarily  $p_0(\alpha) = 0$ .  
 $\beta - 1, \beta - 2, \dots$ , are not roots of  $Q$ , i.e.,  $\beta$  not pole of  $u_{n-1}, \dots, u_{n-r}$ . So  $\beta$  is not a pole of  $p_r(n-r)u_n$ . Since  $\beta$  is a pole of  $u_n$ , necessarily  $p_r(\beta - r) = 0$ .  $\square$

**Corollary.** The dispersion  $H$  is a root of  $\text{Res}_n(p_0(n+h), p_r(n-r)) \in \mathbb{K}[h]$ .

**Proof:**  $(n, h) = (\beta, H)$  is solution of system  $p_0(n+h) = 0, p_r(n-r) = 0$ .  $\square$

**Task:** Given  $p_0, \dots, p_r \in \mathbb{K}[n]$ , compute a basis of solutions in  $\mathbb{K}(n)$  of

$$\text{(LRE): } p_r(n)u_{n+r} + \dots + p_0(n)u_n = 0.$$

**Theorem.** Assume  $u_n = P(n)/Q(n)$  solution of (LRE) with  $\gcd(P, Q) = 1$ . Let  $H$  be the dispersion of  $Q$ . Then

$$Q(n) \mid \gcd(p_0(n) \cdots p_0(n+H), p_r(n-r) \cdots p_r(n-r-H))$$

**Proof:** We've already proved that  $Q(n) \mid p_0(n) \operatorname{lcm}(Q(n+1), \dots, Q(n+r))$ . Substituting  $n \leftarrow n+1$  and repeating yields:

$$\begin{aligned} Q(n) &\mid p_0(n) \operatorname{lcm}(p_0(n+1) \operatorname{lcm}(Q(n+2), \dots, Q(n+r+1)), Q(n+2), \dots, Q(n+r)) \\ &\mid p_0(n)p_0(n+1) \operatorname{lcm}(Q(n+2), \dots, Q(n+r+1)) \\ &\cdots \mid p_0(n) \cdots p_0(n+j) \operatorname{lcm}(Q(n+j+1), \dots, Q(n+j+r)). \end{aligned}$$

As  $\gcd(Q(n), Q(n+j)) = 1$  for  $j > H$ , we get  $Q(n) \mid p_0(n) \cdots p_0(n+H)$ .  $\square$

**Task:** Given  $p_0, \dots, p_r \in \mathbb{K}[n]$ , compute a basis of solutions in  $\mathbb{K}(n)$  of

$$\text{(LRE): } p_r(n)u_{n+r} + \dots + p_0(n)u_n = 0.$$

**Theorem.** Assume  $u_n = P(n)/Q(n)$  solution of (LRE) with  $\gcd(P, Q) = 1$ . Let  $H$  be the dispersion of  $Q$ . Then

$$Q(n) \mid \gcd(p_0(n) \cdots p_0(n+H), p_r(n-r) \cdots p_r(n-r-H))$$

▷ **A refinement:**  $Q$  divides the (generally smaller) polynomial

$$\prod_{i=1}^m \prod_{j=0}^{h_i} \gcd(p_0(n-j+h_i), p_r(n-j-r)),$$

where  $h_1 > h_2 > \dots > h_m \geq 0$  are the roots in  $\mathbb{N}$  of

$$R(h) = \text{Res}_n(p_0(n+h), p_r(n-r)).$$

## Abramov(LRE)

**In:** A LRE  $p_r(n)u_{n+r} + \dots + p_0(n)u_n = 0$  with  $p_i \in \mathbb{K}[n]$

**Out:** A multiple of the denominator of all its rational solutions.

- 1 Compute the polynomial

$$R(h) = \text{Res}_n(p_0(n+h), p_r(n-r)).$$

- 2 If  $R$  has no roots in  $\mathbb{N}$ , then return 1; else, let  $h_1 > h_2 > \dots > h_m \geq 0$  be its roots in  $\mathbb{N}$ . Initialize  $Q$  to 1,  $A$  to  $p_0(n)$ ,  $B$  to  $p_r(n-r)$ .

- 3 For  $i = 1, \dots, m$  do

$$\begin{aligned} g(n) &:= \text{gcd}(A(n+h_i), B(n)); \\ Q(n) &:= g(n)g(n-1) \cdots g(n-h_i)Q(n); \\ A(n) &:= A(n)/g(n-h_i); \\ B(n) &:= B(n)/g(n). \end{aligned}$$

- 4 Return  $Q$ . //  $\prod_{i=1}^m \prod_{j=0}^{h_i} \text{gcd}(p_0(n-j+h_i), p_r(n-j-r))$



The aim of this exercise is to show that the sum

$$S_n := \sum_{k=1}^n \frac{1}{k!}$$

cannot be simplified under the form

$$S_n = r(n)/n!$$

where  $r(n)$  is a rational function in  $n$ .

- 1 Give a linear recurrence with polynomial coefficients satisfied by  $r(n)$ ;
- 2 Find all rational solutions of the recurrence equation:

$$u_{n+1} - (n+1)u_n = 1.$$

- 3 Conclude.