Computer Algebra in the Service of Enumerative Combinatorics

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ABSTRACT
Classifying lattice walks in restricted lattices is an important problem in enumerative combinatorics. Recently, computer algebra has been used to explore and to solve a number of difficult questions related to lattice walks. We give an overview of recent results on structural properties (e.g., algebraicity versus transcendence) and on explicit formulas for generating functions of walks with small steps in the quarter plane. In doing so, we emphasize the algorithmic nature of the methodology, especially two important paradigms: “guess-and-prove” and “creative telescoping”.

CCS CONCEPTS
• Computing methodologies → Algebraic algorithms.

KEYWORDS
Computer algebra; Experimental mathematics; Guess-and-Prove; Creative telescoping; Enumerative combinatorics; Lattice paths; Generating functions; D-finite functions; Algebraic functions.

1 GENERAL PRESENTATION
1.1 Prelude
Consider the following innocent-looking problem.

A tandem-walk is a path in $\mathbb{Z}^2$ taking steps from $\{↑, ←, \downarrow\}$ only. Show that, for any integer $n ≥ 0$, the following quantities are equal:

(i) the number $a_n$ of tandem-walks of length $n$ (i.e., using $n$ steps), confined to the upper half-plane $\mathbb{Z} \times \mathbb{N}$, that start and end at $(0, 0)$;

(ii) the number $b_n$ of tandem-walks of length $n$ confined to the quarter plane $\mathbb{N}^2$, that start at $(0, 0)$ and finish on the diagonal $x = y$.

For instance, when $n = 3$, this common value is $a_3 = b_3 = 3$, as shown in the next picture, obtained by exhaustive enumeration.

It appears that this problem is far from being trivial. Several solutions exist, but none of them is elementary. One of the main aims of the present text is to convince the reader that this problem (and many others with a similar flavor) can be solved with the help of a computer. More precisely, computer algebra tools can be used to discover and to prove the following equalities

\[ a_{3n} = b_{3n} = \frac{(3n)!}{n!^2 \cdot (n+1)!} \quad \text{and} \quad a_m = b_m = 0 \quad \text{if} \ 3 \nmid m. \quad (1) \]

It goes without saying that such a simple and beautiful expression cannot be the result of mere chance. It turns out that closed forms are quite rare for this kind of enumeration problem. Nevertheless, even in the absence of nice formulas, the structural properties of the corresponding enumeration sequences reflect the symmetries of the step set and of the evolution domain. Equation (1) shows that the sequences $(a_n)$ and $(b_n)$ are P-recursive, that is, they satisfy a linear recurrence with polynomial coefficients (in the index $n$). Equivalently, their generating functions $\sum_{n≥0} a_n x^n$ and $\sum_{n≥0} b_n x^n$ are D-finite, that is, they satisfy linear differential equations with polynomial coefficients (in the variable $t$). On the methodological (i.e., computer-algebraic) side, one of the main messages that will emerge from the text is that, in the absence of closed formulas, the (recurrence/differential) equations themselves constitute the appropriate data structure to represent and manipulate P-recursive sequences and D-finite functions [65]. On the application (i.e., combinatorial) side, the main message is that these important properties of the enumeration sequences are intimately related to the finiteness of a certain group, naturally attached to the step set $\{↑, ←, \downarrow\}$.

1.2 General context: walks confined to cones
Let us put the previous problem into a more general framework. Let $d ≥ 1$ be an integer (dimension), let $\mathcal{S}$ be a finite subset (called step set, or model) of vectors in $\mathbb{Z}^d$, and $p_0 ∈ \mathbb{Z}^d$ (starting point). An $\mathcal{S}$-path (or, $\mathcal{S}$-walk) of length $n$ starting at $p_0$, is a sequence $(p_0, p_1, \ldots, p_n)$ of elements in the lattice $\mathbb{Z}^d$ such that $p_{i+1} - p_i ∈ \mathcal{S}$ for all $0 ≤ i < n$. Let $C$ be a cone of $\mathbb{R}^d$, that is, a subset of $\mathbb{R}^d$ such that $r \cdot v ∈ C$ for any $v ∈ C$ and $r ≥ 0$, assumed to contain $p_0$. In combinatorics and in probability theory, one is

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interested in the (exact or asymptotic) enumeration of \(\mathcal{S}\)-walks confined to the cone \(\mathcal{C}\), potentially subject to additional constraints.

**Example 1.1.** Consider the model \(\mathcal{S} = \{-\rightarrow, \leftrightarrow, \rightarrow, \sqrt{1}\}\), that is \(\mathcal{S} = \{(1, 0), (-1, 0), (1, 1), (-1, -1)\}\) (called the Gessel model) in dimension \(d = 2\), with starting point \(p_0 = (0, 0)\) and cone \(\mathcal{C} = \mathbb{R}_{\geq 0}^2\) (the quarter plane). Figure 1 displays the model \(\mathcal{S}\) (on the left), and an \(\mathcal{S}\)-walk of length \(n = 2178\) confined to \(\mathcal{C}\) (on the right).

The main typical questions in this context are the following:
- What is the number \(a_n\) of \(n\)-step \(\mathcal{S}\)-walks contained in \(\mathcal{C}\)?
- For \(i \in \mathcal{C}\), what is the number \(a_{n,i}\) of such walks ending at \(i\)?
- What is the nature of their generating functions

\[
A(t) = \sum_n a_n t^n \quad \text{and} \quad A(t; x) = \sum_{n,i} a_{n,i} t^n x^i.
\]

The answers to these questions are not simple, and heavily depend on the various parameters. The aim of this text is to provide a survey of recent results, notably classification results and closed form expressions, obtained using Computer Algebra.

### 1.3 Why count walks in cones?

Lattice paths are fundamental objects in combinatorics. They have been studied at least since the second half of the 19th century, in connection with the so-called ballot problem [64]. Various approaches have progressively been involved in the study of lattice walks, separately or in interaction. These methods arise from various fields of classical mathematics (algebra, combinatorics, complex analysis, probability theory), and more recently from computer science. There are several reasons for the ubiquity of lattice walks, but the main one is that they encode several important classes of mathematical objects, in discrete mathematics (permutations, trees, words, urns), in statistical physics (magnetism, polymers), in probability theory (branching processes, games of chance) and in operations research (birth-death processes, queueing theory). Therefore, many questions from all these various fields can be reduced to solving lattice path problems. Several books are entirely devoted to lattice paths and their applications (e.g., [35, 36, 59, 60]). For more details, the reader is referred to the introduction of [3].

### 1.4 Blending experimental mathematics and computer algebra

The enumeration of lattice walks is still a topical issue, with a lot of recent activity, new and exciting results, and many open questions. Even if we only count articles published since 2000, there are hundreds of references in this area; significant progress has been obtained by many authors, see [9] and [48] for an extensive bibliography. The dominating viewpoint in these works is to develop uniform approaches, rather than ad-hoc solutions to specific questions. My personal contribution is to combine an experimental mathematics approach, as promoted in the beautiful and inspiring books by Borwein and collaborators [2, 8], with modern tools from the computer algebra arsenal [10, 41], in order to conjecture and prove enumerative and asymptotic results for lattice paths.

### 1.5 Classification of power series

Before stating the main results, let us recall a few definitions on (univariate and multivariate) power series.

**Definition 1.2.** A series \(S(t) = \sum_{n=0}^{\infty} s_n t^n \in \mathbb{Q}\{t\}\) is called:
- **algebraic** if it is a root of a non-trivial polynomial \(P \in \mathbb{Q}\{t, T\}\), i.e., \(P(t, S(t)) = 0\);
- **D-finite** (differentially finite, or holonomic) if it satisfies a non-trivial linear differential equation with coefficients in \(\mathbb{Q}\{t\}\);
- **hypergeometric** if the sequence \((s_n)_n\) satisfies a non-trivial linear homogeneous recurrence of order 1 with coefficients in \(\mathbb{Q}\{n\}\).

An important class of hypergeometric series is that of **Gauss hypergeometric functions** \(2F_1\) with parameters \(a, b, c \in \mathbb{Q}, c \notin \mathbb{N}\):

\[
2F_1 \left( \begin{array}{c} a \ b \\ c \end{array}; t \right) = \sum_{n=0}^{\infty} \frac{(a)_n (b)_n}{(c)_n} \frac{t^n}{n!},
\]

where \((x)_n = x(x+1) \cdots (x+n-1)\) is the rising factorial. This class contains important elementary power series such as

\[
\ln(1-t) = -\sum_{n=1}^{\infty} \frac{t^n}{n}, \quad (1 + t)^\alpha = \sum_{n=0}^{\infty} \binom{\alpha}{n} t^n, \quad (\alpha \in \mathbb{Q}),
\]

and less elementary ones, such as the complete elliptic integral

\[
2F_1 \left( \frac{1}{2}, \frac{1}{2}; 1 \middle| t \right) = \frac{\pi}{2} \int_0^1 \frac{dx}{\sqrt{(1-x^2)(1-tx)}} = \sum_{n=0}^{\infty} \left( \frac{2}{16} \right)^n n!. \quad (2)
\]

The notion of \(2F_1\) admits an obvious extension to the so-called **generalized hypergeometric function** \(pF_q\), involving \(p\) rising factorials in the numerator and \(q\) rising factorials in the denominator.

These three important classes of power series (algebraic, D-finite, hypergeometric) are connected in the following ways. **Hypergeometric series are D-finite:** this is an immediate consequence of the fact that coefficient sequences of D-finite series are exactly P-recursive sequences [68]. **Algebraic series are D-finite:** this fact has been observed by Abel in 1827. Algorithms for the computation of differential equations satisfied by algebraic functions were studied from the complexity perspective by Chudnovsky and Chudnovsky [30], and more recently by Bostan et al. [12]. Finally, understanding which power series are simultaneously algebraic and hypergeometric is an old and difficult question. It was solved for \(2F_1\) by Schwarz [66] in 1873 and in general only one century later, by Beukers and Heckman [7]; the decision procedure is completely algorithmic.
Similar definitions for algebraicity and D-finiteness apply to multivariate power series. For instance, $S \in \mathbb{Q}[x, y]$ is algebraic if it is the root of a non-trivial polynomial $P \in \mathbb{Q}[x, y, t]$, and it is D-finite if the set of all partial derivatives of $S$ spans a finite-dimensional vector space over $\mathbb{Q}(x, y, t)$, in other words if $S$ satisfies a system of linear PDEs with polynomial coefficients of the form
\[
\sum_i a_i(t, x, y) \frac{\partial^i S}{\partial t^i} = 0, \quad \sum_i b_i(t, x, y) \frac{\partial^i S}{\partial y^i} = 0, \quad \sum_i c_i(t, x, y) \frac{\partial^i S}{\partial t^i} = 0.
\]
As in the univariate case, multivariate algebraic series are D-finite [54].

1.6 Walks in the quarter plane

Let us now turn back to the general problem stated in §1.2. The simplest possible cone is the full space $\mathcal{C} = \mathbb{R}^d$. In that case, the full generating function has the simplest possible structure: it is rational. Indeed, a simple reasoning shows that
\[
a_n = |\mathcal{A}|^n, \quad \text{i.e.} \quad A(t) = \sum_{n \geq 0} a_n t^n = \frac{1}{1 - |\mathcal{A}|t},
\]
and, more precisely,
\[
A(t; x) = \sum_{n \geq 0} a_n t^n = \frac{1}{1 - \sum_n x^n t^n}.
\]

The next case by increasing order of difficulty is when the cone is a half-space, typically $\mathbb{R}^{d-1} \times \mathbb{R}_{\geq 0}$. The full generating function $A(t; x)$ is generally not rational anymore, but is nevertheless still algebraic. This result is due to Bousquet-Mélou and Petkovšek, see [25, Theorem 13]. The main ingredient in the proof is the so-called kernel method. A powerful generalization of this method is presented in [23], itself a special case of Popescu’s theorem on Artin approximation with nested conditions [63, Thm. 1.4].

Next comes the case of a cone obtained as the intersection of two half-spaces. Up to modifying the step set by a linear transformation, one may assume that the cone is the basic orthant $\mathcal{C} = \mathbb{R}^d_{\geq 0}$. Quite surprisingly at first sight, generating functions for walks constrained to evolve in an orthant need not be algebraic, and not even D-finite. The first model of walks for which the generating function was proved to be non-D-finite (by Bousquet-Mélou and Petkovšek, in [26, §3]) is the so-called knight walks model: these are walks confined to $\mathbb{Z}^2$ that start from $p_0 = (1, 1)$ and take their steps in $\mathcal{C} = \{(2, -1), (-1, 2)\}$. This surprising result was the starting point of a massive classification effort, initiated by Mishna [57], intensified in a seminal work by Bousquet-Mélou and Mishna [24], and continued by many researchers.

From now on, we focus on small-step walks (or nearest-neighbor walks) in the quarter plane. These are walks in the lattice $\mathbb{Z}^2$, confined to the cone $\mathcal{C} = \mathbb{R}^2_\geq$ (we will often say confined to $\mathbb{R}^2_\geq$), that start at $p_0 = (0, 0)$ and use steps in a model $\mathcal{C}$ which is a fixed subset of $\{(\uparrow, \downarrow, \leftarrow, \rightarrow, \nearrow, \searrow)\}$. Let us denote by $f_{n,i,j}$ the number of walks of length $n$ ending at $(i, j)$. The full counting sequence $(f_{n,i,j})_{n,i,j}$ admits several interesting specializations:

- $f_{n,0,0}$: the number of walks of length $n$ returning to $(0, 0)$;
- $f_n := \sum_{i,j \geq 0} f_{n,i,j}$, the number of walks with prescribed length $n$.

To these enumeration sequences one attaches various power series, namely the full generating function
\[
F_\mathcal{C}(t; x, y) = \sum_{n=0}^{\infty} \left( \sum_{i,j=0}^{\infty} f_{n,i,j} x^i y^j \right) t^n \in \mathbb{Q}[x, y][[t]],
\]
and its corresponding univariate specializations:

- $F_\mathcal{C}(t; 0, 0) := \sum_{n \geq 0} f_{n,0,0} t^n$, the “excursion generating function”;
- $F_\mathcal{C}(t; 1, 1) := \sum_{n \geq 0} f_{n,1,1} t^n$, the “length generating function”;
- $F_\mathcal{C}(t; 0, 1)$, resp. $F_\mathcal{C}(t; 0, 0)$, the generating function of walks ending on the horizontal, resp. vertical, axis (“boundary returns”);
- $\hat{F}_\mathcal{C}(t; 0, \infty) := \left[x^0 \right] F_\mathcal{C}(t; x, 1/x)$, the generating function of walks ending on the diagonal $x = y$ of $\mathbb{R}^2$ (“diagonal returns”).

(Here, and in all the text, $[x^n]$ denotes coefficient extraction of $x^n$.)

The general questions formulated in §1.2 specialize to the quarter-plane setting as follows: given the model $\mathcal{C}$, what can be said about the generating function $F_\mathcal{C}(t; x, y)$, resp. about the counting sequence $(f_{n,i,j})_{i,j,n}$, and their specializations? More precise sub-questions concern structures, explicit forms and asymptotics:

- Structures: is $F_\mathcal{C}$ algebraic? Is it D-finite?
- Explicit forms: do $F_\mathcal{C}(t; x, y)$ and $(f_{n,i,j})_{i,j,n}$ (or some of their specializations) admit closed-form expressions?
- Asymptotics: how do $(e_n)_{n}$ and $(f_n)_{n}$ behave when $n \to \infty$?

1.7 Small-step models of interest

Among the $2^8$ models $\mathcal{C} \subseteq \{-1, 0, 1\}^2 \setminus \{(0, 0)\}$, some are trivial (e.g., if $\mathcal{C} \subseteq \{\downarrow, \leftarrow, \nearrow, \searrow\}$, then $F_\mathcal{C}(t; x, y) \equiv 1$), others are intrinsic to the half-plane (therefore $F_\mathcal{C}(t; x, y) \equiv 0$), and others come in pairs by diagonal symmetry (if $\mathcal{C}$ and $\mathcal{C}'$ are symmetric with respect to the diagonal of $\mathbb{R}^2$, then $F_{\mathcal{C}'}(t; x, y) \equiv F_\mathcal{C}(t; y, x)$).

After discarding these cases, Bousquet-Mélou and Mishna [24] found that there are 79 interesting distinct models of small-step walks in the quarter plane. They are represented in Fig. 2 and are grouped in two classes: 74 non-singular models (or genus-1 models) and 5 singular models (or genus-0 models). Singular models are the ones for which walks never return to the origin, that is for which the excursions generating function is trivial $F(t; 0, 0) \equiv 1$.

Among the 79 models, a few “special” ones are considered interesting enough and were sufficiently studied in order to deserve names (as “special functions” do): Tandem: $\mathcal{C}$ (the one in §1.1); Pólya: $\mathcal{C}$; Kreweras: $\mathcal{C}$; Gessel: $\mathcal{C}$; Gouyou-Beauchamps: $\mathcal{C}$;
1.8 Two nice models: Kreweras’ and Gessel’s

An interesting model in the world of quarter-plane walks is Kreweras’ model \( \mathcal{K} \). It is related to a version of the three-candidate ballot problem [61] and also to an applied probabilistic context [39] (double queue that arises when arriving customers simultaneously place two demands handled independently by two servers). In a long paper, Kreweras [49] obtained the following result, which was reproved using various methods in [5, 16, 20, 21, 42, 45, 50, 61]. References [16, 45] provide two different computer-aided proofs. In what follows, we denote by \( K(t; x, y) = F_\mathcal{K}(t; x, y) \) the full generating function for Kreweras walks in the quarter plane, and by \( K(t; 0, 0) \) the generating function for Kreweras excursions.

**Theorem 1.3** ([49]). \( K(t; 0, 0) = 1 + 2t^3 + 16t^5 + \cdots \) is equal to

\[
3F_2\left(\begin{array}{c} 1 \frac{3}{2} \frac{3}{2} 1 \\ 0 \frac{3}{2} \frac{3}{2} 0 \\ \sqrt{t} \end{array} \right) = \sum_{n=0}^{\infty} \frac{4^n (3n)}{(n+1)(2n+1)} t^{3n}.
\] (3)

Theorem 1.3 implies, together with the results in §1.5 (e.g. [7]) that \( K(t; 0, 0) \) is an algebraic power series. In fact, much more is true:

**Theorem 1.4** ([21, 39, 42]). \( K(t; x, y) \) is algebraic.

The most difficult model of small-step walks in the quarter plane is Gessel’s model \( \mathcal{G} \). In 2001, Ira Gessel formulated two conjectures equivalent to the following statements:

**Conjecture 1.** The generating function \( G(t; 0, 0) = 1 + 2t^2 + 11t^4 + 85t^6 + 782t^8 + \cdots \) of Gessel excursions is equal to

\[
3F_2\left(\begin{array}{c} \frac{1}{2} \frac{1}{2} \frac{1}{2} 2 \\ 0 \frac{1}{2} \frac{1}{2} 0 \\ \sqrt{t} \end{array} \right) = \sum_{n=0}^{\infty} \frac{2^n n!}{\left(\frac{3}{2} n\right) n!} (4t)^{2n}.
\]

**Conjecture 2.** The generating function \( G(t; x, y) \) is not D-finite.

Here, \( G(t; x, y) = F_\mathcal{G}(t; x, y) \) for \( \mathcal{G} = \{\nearrow, \searrow, \leftarrow, \rightarrow\} \) denotes the full generating function for Gessel walks in the quarter plane, and \( G(t; 0, 0) \) denotes the generating function for Gessel excursions.

1.9 Algebraic translation: functional equations

Gessel’s problem admits the following purely algebraic reformulation. If \( G(t; x, y) \in \mathbb{Q}[x, y][[t]] \) denotes the full generating function for Gessel walks in the quarter plane, then a simple inclusion-exclusion reasoning represented pictorially in Fig. 3 implies that \( G(t; x, y) \) satisfies a functional equation called the **kernel equation**

\[
G(t; x, y) = 1 + t \left( xy + x + \frac{1}{xy} + \frac{1}{x} \right) G(t; x, y) - t \frac{1}{x} G(t; 0, y) - t \frac{1}{xy} (G(t; x, 0) - G(t; 0, 0)).
\] (4)

Moreover, \( G(t; x, y) \) is completely characterized by the functional equation (4): it is its unique solution in \( \mathbb{Q}[x, y][[t]] \), and even in the ring \( \mathbb{Q}[x, y, t] \). Therefore, the task is simply to solve equation (4).

To any of the 79 models introduced in §1.7 is attached a very similar functional equation. Again, this equation merely reflects a step-by-step construction of quarter-plane walks, and is based on the most elementary decomposition: a walk is either the empty walk, or a shorter walk followed by a permissible step. This observation is naturally translated into a generating function equation using the **inventory** \( \chi_{\mathcal{G}}(x, y) \), \( x(t) = \sum_{(i, j) \in \mathcal{G}} x^i y^j \), and the kernel \( R_{\mathcal{G}}(t; x, y) := xy(1 - t \chi_{\mathcal{G}}(x, y)) \). The decomposition is translated into the kernel equation (we omit the subscript \( \mathcal{G} \)):

\[
R(t; x, y)F(t; x, y) = xy + R(t; x, 0)F(t; x, 0) + R(t; 0, y)F(t; 0, y) - R(t; 0, 0)F(t; 0, 0).
\] (5)

Observe that the last term of the right-hand side occurs only if the step \( \uparrow \searrow \) belongs to the model \( \mathcal{G} \).

2 MAIN RESULTS

In §1.9 we saw that classifying lattice walks in the quarter plane amounts to solving 79 equations of the form (5). In this section we describe several classes of results that have been obtained, using computer algebra tools, on the solutions of these 79 equations.

2.1 Algebraicity of Gessel walks

After a first attempt in [45], Gessel’s first conjecture was finally solved in 2009 by Kauers, Koutschan and Zeilberger in [44] using an extension of the guess-and-prove approach described in [45].

**Theorem 2.1** ([44]). \( G(t; 0, 0) = 3F_2\left(\begin{array}{c} 5/6 1/2 1 \\ 3/5 2/3 0 \end{array} \right) \).

This result implies in particular that \( G(t; 0, 0) \) is D-finite, but has no immediate implications concerning the D-finiteness of \( G(t; x, y) \). It came as a total surprise when Bostan and Kauers [16] proved that Gessel’s second conjecture was false.

**Theorem 2.2** ([16]). The generating function \( G(t; x, y) \) for Gessel walks is algebraic.

Prior to this result, even the algebraicity of \( G(t; 0, 0) \) had been overlooked, even though the classical results recalled in §1.5 apply. The original discovery and proof of Theorem 2.2 was computer-driven, and used a guess-and-prove approach, based on **Hermite–Padé approximants** [16]. Note that as a byproduct of this proof, an estimate on the size of the minimal polynomial of \( G(t; x, y) \) has been given: it has more than \( 10^{11} \) terms when written in dense (expanded) form, for a total size of \( \approx 30 \) Gb (!). The guess-and-prove method is a 3-step process:

1. **(S1)** compute \( G(t; x, y) \) to precision \( t^{1200} \) (= 1.5 billion coefficients);
2. **(S2)** conjecture polynomial equations for \( G(x; 0, t) \) and \( G(0, y; t) \) (degree 24 each, coeffs. of degree (46, 56), with 80-bit digits coeffs.), relying on the Beckermann–Labahn algorithm [4];
3. **(S3)** conclude the proof by computing multivariate results of (very big) polynomials (30 pages each), using fast algorithms for manipulating algebraic series [14].
Several human proofs of Theorems 2.1 and 2.2 have been discovered since the publication of [16]; the first one used complex analysis [17], the second one was purely algebraic [22], the third one is both combinatorial and analytic [28], while the more recent one is probably the most elementary [6].

### 2.2 Explicit form for \( G(t; x, y) \)

An interesting refinement of Theorem 2.2 is the next result, which contains a closed-form formula for the generating function \( G(t; x, y) \).

**Theorem 2.3** ([16]). Let \( V = 1 + 4t^2 + \cdots \) be the root in \( Q[t] \) of \((V-1)(1+V)^3 = (16t^2)\), let \( U = 1 + 2t^2 + 16t^4 + 2t^6 + \cdots \) be the root in \( Q[[t]] \) of

\[
x(V-1)(V+1)U^3 - 2V(3x + 5xV - 8Vt)U^2
\]

and let \( W = t^2 + (y + 8)t^4 + 2(y^2 + 8y + 44)t^6 + \cdots \) be the root in \( Q[y][[t]] \) of \( y(V-1)W^4 + 5y(V + 3)W^2 - (V + 3)W + V - 1 = 0 \).

Then \( G(t; x, y) \) is equal to

\[
\frac{64(U/V+1-2V)^{1/2}}{(y(V-1)(1-Wy)W^{3/2}+(y+V)W)} - \frac{1}{xy(y+1)}.
\]

Again, the original discovery and proof of this result was computer-driven, based on effective Galois theory. During the proof, a few other remarkable facts have been noticed, namely that \( G(t; x, y) \) can be expressed using nested radicals; for instance the length generating function \( G(t; 1, 1) = 1 + 2t + 7t^2 + 21t^3 + 78t^4 + \cdots \) reads

\[
G(t; 1, 1) = -\frac{1}{2t} + \frac{\sqrt{5}}{6t} \sqrt{H(t) + \frac{16t(2t + 3) + 2}{(1 - 4t)^2H(t)} - H(t)^2 + 3},
\]

where \( H(t) = \sqrt{1 + 4t^{1/3}(1 + 4t)^{2/3}/(1 - 4t)^{4/3}} \).

The proof of Theorem 2.3 uses the minimal polynomials for \( G(t; x, 0) \) and \( G(t; 0, y) \) that were guessed and proved during the algebraicity proof. A striking feature of Theorem 2.3 is the relative simplicity of the closed-form expression, especially when compared to the size of the minimal polynomial of \( G(t; x, y) \). Similarly to Theorem 2.2, the result in Theorem 2.3 admits several recent human proofs [6, 17, 22].

### 2.3 Models with D-finite generating functions

The computer-driven approach that allowed Bostan and Kauers [16] to discover and prove the properties of the puzzling generating function for Gessel walks was used as early as 2008 by the same authors to provide a (conjecturally) exhaustive list of models having (conjecturally) D-finite and algebraic generating functions [15]. That resulted in an experimental classification synthesized in Fig. 4, which displays 23 models of walks in the quarter plane for which the minimal generating function \( F(t; 1, 1) \) was conjectured to be D-finite. The computerized discovery used again (differential) guessing, based on Hermite–Padé approximation. Note that the guess-and-prove method used to prove Theorem 2.2 is robust enough to also prove that all models in Fig. 4 are indeed D-finite (with resultant computations in [53] replaced by differential closure algorithms). The labels used in column "OEIS" are taken from Sloane’s On-Line Encyclopedia of Integer Sequences [67]. The "alg?" column refers to the algebraicity or transcendence of \( F(t; 1, 1) \): cases marked "Y" were conjectured algebraic (using algebraic Hermite–Padé approximation), the other cases were conjectured transcendental. For cases 1–22, these results on D-finiteness, resp. algebraicity, were confirmed by human proofs obtained almost simultaneously with [15] by Bousquet-Mélou and Mishna [24], using a uniform approach. We discussed the difficult case 23 (Gessel’s model) in §2.1 and §2.2. Concerning the conjectural transcendence results, the first unified proof was given in [13] and it is computer-driven; this will be discussed in §2.6. The reference [13] also contains the first proof, again computer-driven, that the (differential / recurrence / algebraic) equations conjectured in [15] are indeed correct.

In addition, Bostan and Kauers demonstrated that computer algebra tools are also able to produce conjectural expressions for the asymptotics of \( f_{n} = [t^{n}] F(t; 1, 1) \). Their results are displayed in column "asymptotics" of Fig. 4 and have been obtained using a combination of algorithmic tools, including Hermite–Padé approximation, constant recognition algorithms built on integer relation detection algorithms like LLL [53] and PSLQ [37], and convergence acceleration techniques [27]. These results have been confirmed a few years later by human proofs by Melczer and Wilson [56], using the theory of analytic combinatorics in several variables [62] that extends the univariate theory from [38].

### 2.4 The group of a model

In order to formulate more results on the classification of lattice walks in the quarter plane, we need to introduce an important concept, the group of a small-step walk model \( \mathcal{F} \). Let us decompose the inventory \( \chi_{\mathcal{F}} \) of \( \mathcal{F} \) along powers of \( x \), resp. of \( y \), as follows:

\[
\chi_{\mathcal{F}}(x, y) = \sum_{(i,j) \in \mathcal{F}} x^{i} y^{j} = \sum_{i=1}^{B_{1}} B_{i}(y)x^{i} = \sum_{j=1}^{A_{1}} A_{j}(x)y^{j}.
\]
The basic, yet fundamental, observation is that $\chi_\mathcal{S}(x, y)$ is left invariant under two rational transformations

$$
\psi(x, y) = \left( \frac{A_{-1}(x)}{A_{+1}(x)} \right), \quad \phi(x, y) = \left( \frac{B_{-1}(y)}{B_{+1}(y)} \right),
$$

and thus under any element of the group $\mathcal{G}_\mathcal{S} := \{\psi, \phi\}$ of birational transformations generated by $\psi$ and $\phi$. When it is finite, $\mathcal{G}_\mathcal{S}$ is isomorphic to a dihedral group, since $\psi$ and $\phi$ are involutions.

Mishna [57] showed that for the 23 models in Fig. 4, the group is finite, of cardinality either 4 (for models 1–16 with an axial symmetry), 6 (for the models 17, 18, 20, 21, 22, with a diagonal or an anti-diagonal symmetry), or 8 (for the remaining models 19 and 23). Boussquet-Mélou and Mishna [24] proved in addition that for all the other 56 models, the group is infinite.

At this point, we know that the finiteness of the group for some model implies the D-finiteness of the generating function for that model. One important remaining question is: is the converse true? Another important question is: in the D-finite cases, are there any closed-form expressions for the generating functions? The next two subsections will bring answers and completely clarify the situation.

### 2.5 Explicit expressions

Models 20–23 in Fig. 4 admit full generating functions that are algebraic. Moreover, closed formulas exist for them. For the three models 20–22 related to the Kreweras model, such formulas are displayed in [24, 46]. The most difficult case among these four is model 23 (Gessel’s), for which Theorem 2.3 provides a closed-form expression. We now focus on models 1–19.

As a consequence, $F_\mathcal{S}(t; x, y)$ is expressible using iterated integrals of hypergeometric $\mathcal{G}_t$ expressions. More precisely, the expressions of the generating functions $F_\mathcal{S}(t; x, y)$ and thus under any element of the group $\mathcal{G}_\mathcal{S} := \{\psi, \phi\}$ of birational transformations generated by $\psi$ and $\phi$. When it is finite, $\mathcal{G}_\mathcal{S}$ is isomorphic to a dihedral group, since $\psi$ and $\phi$ are involutions.

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**Theorem 2.4 ([13]).** Let $\mathcal{S}$ be one of the models 1–19 in Fig. 4. Then $F_\mathcal{S}(t; x, y)$ is expressible as a finite sum of iterated integrals of products of algebraic functions in $x$, $y$, $t$ and of expressions of the form $\mathcal{G}_t\left(\begin{array}{c}a \\ b \\ c \end{array}\right)\omega(t)$, where $c \in \mathbb{N}$ and $\omega(t) \in \mathbb{Q}(t)$.

Once again, the discovery and proof of this result are computer-driven; no human proof is available yet. The proof is based, among other tools, on creative telescoping, an efficient algorithmic technique for the symbolic integration of multivariate functions with parameters. This powerful computer algebra tool was introduced in the early 1990s by Zeilberger in the “hyper” setting [70, 71], vastly generalized by Chyzak in the 2000s to the framework of holonomic functions [31], and further enhanced by Koutschan in the 2010s [46]. Since its birth, 30 years ago, the methodology of creative telescoping gained more and more popularity. To date, there are no less than four generations of creative telescoping algorithms; examples of 4G integration algorithms are in [18] and [11]. For more details, we refer the reader to the surveys [29, 32, 47].

Using the kernel method, Boussquet-Mélou and Mishna showed in [24, Prop. 8] that the generating function $F_\mathcal{S}(t; x, y)$ for models 1–19 can be written in terms of the D-part of a rational function,

$$
F_\mathcal{S}(t; x, y) = \frac{1}{xy} \left[ x^r \right] \left[ y^s \right] \frac{N(x, y)}{1 - tS(x, y)},
$$

where $N(x, y)$ and $S(x, y)$ are certain Laurent polynomials in $y$ with coefficients that are rational functions in $x$. By a result of Lipshitz [54, Thm. 2.7], Equation (6) already implies the D-finiteness of $F_\mathcal{S}(t; x, y)$. To obtain explicit differential equation for $F_\mathcal{S}(t; x, y)$ and its specializations, the key is creative telescoping. Once a linear differential equation is obtained, the last steps in the proof of Theorem 2.4 rely on factoring differential operators [69] and on solving second-order differential equations in terms of $\mathcal{G}_tF_\mathcal{S}$ [43, 51].

The expressions of the generating functions $F(t; 0, 0), F(t; 0, 1), F(t; t, 1), F(t; t, 0), F(t; x, 0), F(t; y)$ and $F(t; x, y)$ are too large to be displayed here, and are available online. It turns out that the involved hypergeometric functions have a very particular form: they are intimately related to elliptic integrals such as (2). For instance, for King walks (model 4), the length generating function is

$$
F_{\mathcal{S}}(t; 1, 1) = \int_0^t \frac{1}{(1 + 4x)^3} \cdot 2F_1\left(\frac{3}{2}, 2; 3; \frac{16(x + 1)}{(1 + 4x)^2}\right) \, dx.
$$

No human proof is currently available for the structure Theorem 2.4, and even less for the explicit expressions implied by it. However, the book [36] probably contains the premises of the key explanation: the kernel for models 1–19 is an elliptic curve, therefore the solution of the kernel equation (5) should be related to elliptic integrals.

### 2.6 Transcendence results

As previously mentioned, models 20–23 in Fig. 4 admit full generating functions that are algebraic. What about the full generating function $F_\mathcal{S}(t; x, y)$, and its combinatorially meaningful specializations $F_\mathcal{S}(t; 0, 0), F_\mathcal{S}(t; 0, 1), F_\mathcal{S}(t; 1, 0)$ and $F_\mathcal{S}(t; 1, 1)$ for the models 1–19? Computer algebra is able to answer this question.

**Theorem 2.5 ([13]).** Let $\mathcal{S}$ be one of the models 1–19 in Fig. 4. Then for any $(\alpha, \beta) \in \{(0, 0), (1, 0), (0, 1), (1, 1)\}$, the power series $F_\mathcal{S}(t; \alpha, \beta)$ is transcendental, except in the following four cases:

- $\mathcal{S} = \text{(model 17, ”Tandem”)}$ and $(\alpha, \beta) = (1, 1), \quad \mathcal{S} = \text{(model 18)}$ and $(\alpha, \beta) = (0, 0), (0, 1), (1, 1)$.

As a consequence, $F_\mathcal{S}(t; x, y), F_\mathcal{S}(t; x, 0)$, and $F_\mathcal{S}(t; y, 0)$ are transcendental for all the 19 models. Additionally, the generating functions of the four algebraic cases are equal to:

- $F_\mathcal{S}(t; 1, 1) = \frac{1}{24t^3} \left(1 - t - \sqrt{1 + (1 + t)(1 - 3t)}\right),
- F_\mathcal{S}(t; 1, 1) = 2F_1\left(2; 1, 1\right),
- F_\mathcal{S}(t; 0, 1) = 2F_1\left(0; 1\right) = \frac{(1 - 6t)^{1/2}(1 + 2t)^{1/2} - 4t^2 + 8t - 1}{32t^2}.

Again, the proof of Theorem 2.5 is computer-driven and heavily relies on the use of several modern computer algebra algorithms. Transcendence proofs had been considered in some isolated cases only. The computer-aided proof of Theorem 2.5 was the first complete and unified transcendence proof applying to all 19 models.

### 2.7 Non-D-finiteness results

The last question in view of the complete classification of small step walks in the quarter plane concerns the 56 models with an infinite group. Among them, 5 models are singular; for them, a variant of the kernel method, called the iterated kernel method was used by Mishna and Rechnitzer [58] (for two models) and by Melczer and Mishna [55] (for all five models), who showed...
that the length generating function \( F(t; 1, 1) \), and thus also the full generating function \( F(t; x, y) \), are non-D-finite. The remaining question concerns the 51 non-singular models with an infinite group: is the full generating function again non-D-finite?

Computer algebra is able to help prove the following result.

**Theorem 2.6** ([19]). Let \( \mathcal{S} \subseteq \{0, \pm 1\}^2 \) be any of the 51 non-singular step sets in \( \mathbb{Z}^2 \) with infinite group \( \mathcal{G} \). Then the generating function \( F_{\mathcal{S}}(t; 0, 0) \) of \( \mathcal{S} \)-excursions is not D-finite. Equivalently, the excursion sequence \((e_n)_{n \geq 0}\) is not P-recursive.

In particular, the full generating function \( F_{\mathcal{S}}(t; x, y) \) is not D-finite in the 51 cases, since D-finiteness is preserved by specialization [54]. This corollary had been already obtained by Kurkova and Raschel [52], but the approach in [19] is simpler, and at the same time it delivers a more accurate information. This new proof combines asymptotic information about the coefficients \( e_n \) of \( F_{\mathcal{S}}(t; 0, 0) \), and arithmetic information about the constrained behavior of the asymptotics of \( e_n \) when \( F_{\mathcal{S}}(t; 0, 0) \) is D-finite. More precisely, [19] first makes explicit consequences of the general results by Denisov and Wachtel [34] in the case of walks in the quarter plane.

Moreover, under (1)–(5), \( |\mathcal{G}| \) is equal to 2 min \( \{ t \in \mathbb{N}_g \mid \frac{1}{2\pi t} \in \mathbb{Z} \} \). Still under (1)–(5), \( F_{\mathcal{S}}(t; x, y) \) is algebraic if and only if \( \text{OS}_{\mathcal{S}} = 0 \).

We now state the main result of this text, obtained by combining all previous results. It provides a complete characterization of the non-singular small-step sets with D-finite full generating function.

**Theorem 2.2.** Let \( \mathcal{S} \subseteq \{0, \pm 1\}^2 \) be any of the 74 nonsingular quarter-plane models in Fig. 2. The following assertions are equivalent:

1. the full generating function \( F_{\mathcal{S}}(t; x, y) \) is D-finite;
2. the excursions generating function \( F_{\mathcal{S}}(t; 0, 0) \) is D-finite;
3. the sequence \([z^n] F_{\mathcal{S}}(t; 0, 0) \) is \( \sim k \cdot \rho^n \cdot n^z \), with \( \alpha \in \mathbb{Q} \);
4. the group \( \mathcal{G} \) is finite;
5. \( \mathcal{S} \) has either an axial symmetry, or zero drift and \( \text{OS}_{\mathcal{S}} \neq 5 \).

Moreover, under (1)–(5), \( |\mathcal{G}| \) is equal to 2 min \( \{ t \in \mathbb{N}_g \mid \frac{1}{2\pi t} \in \mathbb{Z} \} \).

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**3 BIOGRAPHY**

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